Sufficiency

Ontline

- 1) Review

 - 2) Sufficiency
 3) Factorization Theorem

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Sufficiency
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Motivation: Coin flipping Suppose X,,..., X ind Bernoulli (0) $\Rightarrow \chi \sim \pi \Theta^{\chi_i}_{(1-\Theta)}^{-\chi_i} \qquad \text{on} \qquad [0,13]^n$ Then $T(X) = \sum X_i \sim Binom(n, 0)$ $= O^{t}(1-\Theta)^{n-t} \binom{n}{t} \qquad \text{on} \qquad \{0,-,n\}$ (X,,-,X) -> T(x) is throwing away data. How do we justify this? In exp. fam. lingo, T(x) is the "sufficient etathstic" for X. Today we'll see why we call it that. Definition Let P= [Po: 0 = a] be a statistical

model for data X. T(x) is sufficient for P if P(XIT) does not depend on O

Example (Cont'd) $P_{o}(X=x, T=t)$ $P_{\theta}(X=x|T=\epsilon)=$ $P_{\theta}(T = \epsilon)$ $= \frac{0^{\xi_{x_i}} (1-0)^{n-\xi_{x_i}}}{0^{t} (1-0)^{n-t} (1)}$ $= 1\{\xi \times_i = t\} / \binom{n}{t}$

So given T(x) = t, X is uniform on all seq.s with $\sum x_i = t$

Interpretations of Sufficiency

Recall we only rare about X in the first place because it is (indirectly) informative about 0 Sufficiency means only T(x) is informative

We can think of the data as being generated in two stages:

silce of x'1) Generate T: distribution dep. on O "Generate XIT: does <u>not</u> dep on O within the liver principle within cach slice.

Sufficiency Principle within cach slice.

If T(x) is sufficient for P then any statistical procedure should depend on X only through T(x)

In fact, we could throw away X and generate a new $\hat{X} \sim P(X|T)$ and it would be just as good as X

In graphical model form:

Step 2

Step 2

Step 2

No reason to pay

Any attention

Step 2

Tust as good as X!

Factorization Theorem

There is a very convenient way to verify sufficiency of a statistic based only on the density:

Theorem (Factorization Theorem)

Let $P = \{P_0 : O \in \Theta\}$ be a family of distributions dominated by M ($P << M + \Theta$) densities P_0 .

T is sufficient for P iff there exist non-neg.

functions go, h such that

 $\rho_{\theta}(x) = g_{\theta}(T(x)) h(x)$ for a.e. x under M $[n(x \cdot p_{\theta} \neq g_{\theta}(T(x)) \cdot h(x)] = 0$ Rigorous proof in Keener 6.4

"Physics proof": (rigorous for discrete X)

 $(E) p(x|T=t) = 1\{T(x)=t\} \cdot \frac{g_0(t) h(x)}{\sum_{T(z)=t}^{g_0(t) h(z)} d_{\mu(z)}}$

(=)) Take $g_0(t) = \int_{T(x)=t}^{p_0(x)} dn(x) = P_0(T(x)=t)$

 $h(x) = \frac{\rho_0(x)}{\int_{\mathbb{T}(z)=1}^{\rho_0(z)} d\mu(z)}$

 $= \mathcal{H}_{\mathcal{A}}(X = X \mid T(X) = T(X))$

 $g_{o}(T(x))h(x) = P_{o}(T = T(x))P(X = x | T = T(x))$

Examples

Ex. Exponential Families

$$P_{\theta}(x) = e^{\gamma(\theta)^{T}(x)} - B(\theta) h(x)$$
 $P_{\theta}(x) = e^{\gamma(\theta)^{T}(x)} - B(\theta) h(x)$

Ex. $X_{1},..., X_{n} \stackrel{iid}{\sim} P_{\theta}^{(1)}$ for any model

 $P_{\theta}^{(1)} = \{P_{\theta}^{(1)} : \theta \in \Theta\}$ on $X \leq \mathbb{R}$

Position is invariant to perm s of $X = (X_{1},..., X_{n})$
 $\Rightarrow \frac{\text{order statistics}}{\text{are sufficient}} (X_{(i)})_{i=1}^{n} (X_{(i)} = e^{ixt} \text{ smellest})$
 $\Rightarrow \frac{\text{order statistics}}{\text{loses information}} (X_{(i)})_{i=1}^{n} (X_{(i)})_{i=1}^{n}$
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Not important that it's a measure in this context; just keeps trade of which values came up how many times

$$E_{X}. \quad X_{1}, ..., X_{n} \stackrel{iid}{\sim} U[\Theta, \Theta+1]$$

$$= 1\{\Theta \leq x \leq \Theta+1\}$$

$$P_{\Theta}(x) = \frac{1}{i=1} 1\{\Theta \leq x_{i} \leq \Theta+1\}$$

$$= 1\{\Theta \leq X_{(i)}\} 1\{X_{(n)} \leq \Theta+1\}$$

$$= (X_{(i)}, X_{(n)}) \quad \text{is sufficient.}$$

Minimal Sufficiency Consider $X_1, ..., X_n \stackrel{iid}{\sim} N(\theta, i)$: $T(x) = \sum X_i \quad \text{sufficient}$ $\overline{X} = \frac{1}{n} \sum X_i \quad \text{also}$ $S(x) = (X_{(1)}, ..., X_{(n)}) \quad \text{too}$ $X = (X_1, ..., X_n) \quad \text{too}$

Which can be recovered from which others?

These can be compressed

further

These are the most

as possible?

compressed as possible?

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Prop If T(X) is sufficient and T(x) = f(s(x))
    then S(x) is sufficient
\frac{P_{root}}{P_{oot}} : \rho_{o}(x) = g_{o}(T(x)) h(x)
                  = (go of) (S(x)) h(x)
Definition: T(X) is minimal sufficient ; f
      1) T(x) is sufficient
      a) For any other sufficient S(X)
          T(x) = f(s(x)) for some f

(a.s. in P)
 So, no neither how many more suff. stats we add
 to our diagram, they will all have arrows
 pointing to EXi
How to check minimal sufficiency? Basically, "equivalent to knowing likelihood ratios"
Definition Assume J= {Po:0 = (A) has densiting
   Po(x) wrt common M, data X. The (log)likelihood function is the (log) density, reframed as a random function of \Theta:
    Lik(\Theta;X) = p_{\Theta}(x)
                                    (0; x) = log Lik(0; x)
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Note if T(X) is sufficient then $Lik(\theta;x) = g_{\theta}(T(x)) h(x)$ T determines the scaling "shape" Than 3.11 Assume 9 = (Po: OE a)?, densities Po T(X) sufficient for X. If $Lik(\theta; x) \propto_{\theta} Lik(\theta; y) \Rightarrow T(x) = T(y)$ then T(X) is minimal sufficient

"I determines the likelihood shape in a
one-to-one fashion"

I < M Proof Suppose S is sufficient and Af s.t. f(s(x)) = T(x)Then $\exists_{x,y}$ with S(x) = S(y), $T(x) \neq T(y)$ $Lik(\theta;x) = g_{\theta}(S(x))h(x)$ = $Lik(\theta; y)$ that implies T(x) = T(y) by assumption.

Ex.
$$\rho_{\theta}(x) = e^{\gamma(\theta)'T(x) - \beta(\theta)} h(x)$$

Is $T(x)$ minimal? doesn't const. in θ change with x

Assume Lik(θ ; x) σ_{θ} Lik(θ ; y). WTS $T(x) = T(y)$

Lik(θ ; x) σ_{θ} Lik(θ ; y). WTS $T(x) = T(y)$
 $\Leftrightarrow \gamma(\theta)'T(x) = \gamma(\theta)'T(y) + \alpha(x,y)$
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 $\Leftrightarrow \gamma(\theta)'T(x) = \gamma(\theta)'T(y) + \alpha(x,y)$
 $\Leftrightarrow \gamma(\theta) - \gamma(\theta_2) T(x) = (\gamma(\theta) - \gamma(\theta_1))T(y) + \gamma(\theta_2)$
 $\Leftrightarrow \gamma(\theta_1) - \gamma(\theta_2) T(x) = (\gamma(\theta_1) - \gamma(\theta_2) + \gamma(\theta_2$

[could T(x) still be minimal?]

Ex.
$$X \sim N_2(\Lambda(\theta), I_2)$$
 $\Theta \in \mathbb{R}$

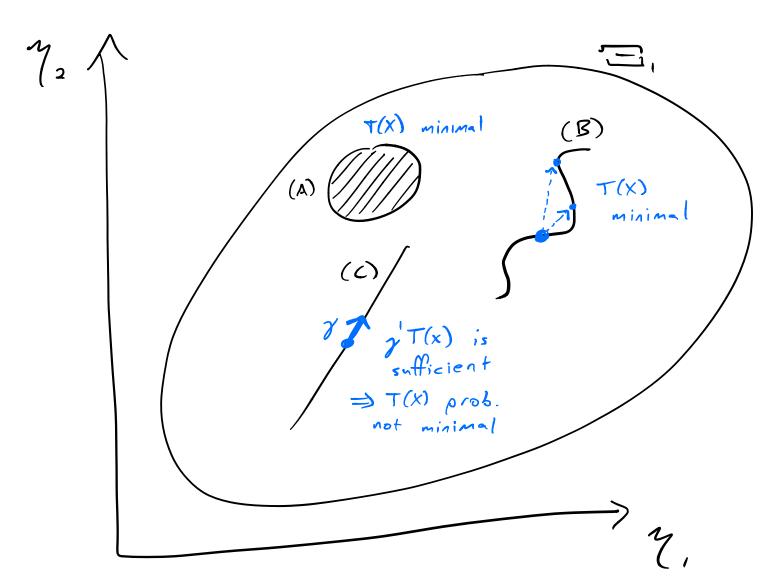
$$= e^{M(\theta)' \times -B(\theta)} e^{-\frac{1}{2} \times ' \times}$$

If $\Omega = \mathbb{R}$, $\Lambda(\theta) = a + \theta b$ for a, $b \in \mathbb{R}^2$

$$P_{\theta}(x) = e^{(a + \theta b)' \times -B(\theta)} e^{-\frac{1}{2} \times ' \times}$$

$$= e^{\theta(b' \times) -B(\theta)} e^{-\frac{1}{2} (x - 2a)' \times}$$

$$b' \times is sufficient \Rightarrow \times not minimal$$



Ex Laplace location family

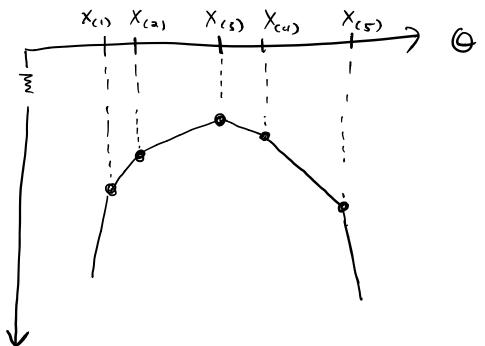
$$X_{1,1}, X_{n} \stackrel{\text{iid}}{\sim} \rho_{\theta}^{(i)}(x) = \frac{1}{2} e^{-|x-\theta|}$$

$$P_{\theta}(x) = \frac{1}{2} e^{-|x-\theta|} e^{-|x-\theta|}$$

$$I(\theta; x) = \log p_{\theta}(x)$$

$$= -\frac{2}{1} |x_{i}-\theta| - \log 2$$
Piecewise linear in θ , knots at

Piecewise linear in O, knots at X(i)



$$l(0; X) = l(0; Y) + const \Leftrightarrow X, Y same order statistics$$

Thm 3.11 => order state are minimal suff.