

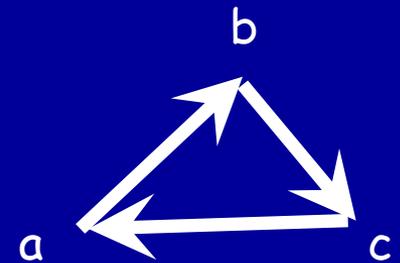
Arrow Theorem

Elchanan Mossel
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Condorcet Paradox

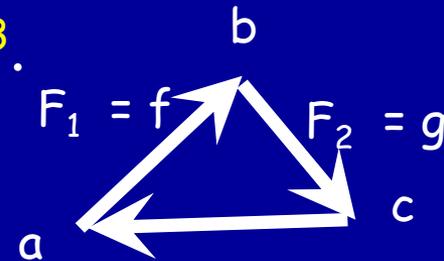
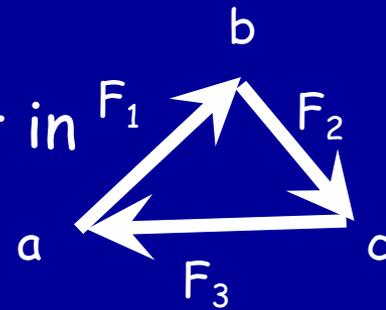
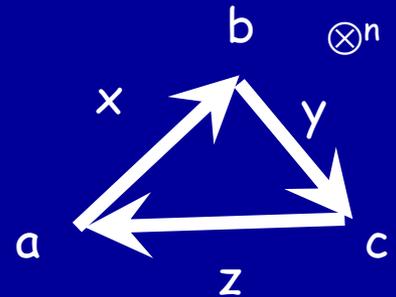


- n voters are to choose between 3 alternatives.
- Condorcet: Is there a rational way to do it?
- More specifically, for majority vote:
- Could it be that all of the following hold:
 - Majority of voters rank **a** above **b**?
 - Majority of voters rank **b** above **c**?
 - Majority of voters rank **c** above **a**?
- Condorcet(1785): Could be.
- Defined by Marquis de Condorcet as part of a discussion of the best way to elect candidates to the French academy of Science.



Properties of Constitutions

- n voters are to choose between 3 alternatives
- Voter i ranking $:= \sigma_i \in S(3)$. Let:
 - $x_i = +1$ if $\sigma_i(a) > \sigma_i(b)$, $x_i = -1$ if $\sigma_i(a) < \sigma_i(b)$,
 - $y_i = +1$ if $\sigma_i(b) > \sigma_i(c)$, $y_i = -1$ if $\sigma_i(b) < \sigma_i(c)$,
 - $z_i = +1$ if $\sigma_i(c) > \sigma_i(a)$, $z_i = -1$ if $\sigma_i(c) < \sigma_i(a)$.
- Note: (x_i, y_i, z_i) correspond to a σ_i iff $(x_i, y_i, z_i) \notin \{(1,1,1), (-1,-1,-1)\}$



- Def: A constitution is a map $F : S(3)^n \rightarrow \{-1,1\}^3$.
- Def: A constitution is transitive if for all σ :
 - $F(\sigma) \in \{-1,1\}^3 \setminus \{(1,1,1), (-1,-1,-1)\}$
- Def: Independence of Irrelevant Alternatives (IIA) is satisfied by F if: $F(\sigma) = (f(x), g(y), h(z))$ for all σ and some f, g and h .

Arrow's Impossibility Thm



- Def: A constitution F satisfies Unanimity if

$$\sigma_1 = \dots = \sigma_n \Rightarrow F(\sigma_1, \dots, \sigma_n) = \sigma_1$$

- Thm (Arrow's "Impossibility", 61): Any constitution F on 3 (or more) alternatives which satisfies

- **IIA**,
- **Transitivity** and
- **Unanimity:**

Arrow received a Nobel Prize in Economics in 1972

Is a dictator: There exists an i such that:

$$F(\sigma) = F(\sigma_1, \dots, \sigma_n) = \sigma_i \text{ for all } \sigma$$

Arrow's Impossibility Thm



The Royal Swedish Academy of Sciences has decided to award the Bank of Sweden Prize in Economic Sciences in Memory of Alfred Nobel, 1972 , to

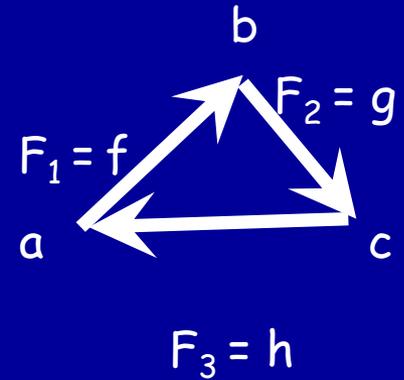
John R Hicks, Oxford University, U K
and

Kenneth Arrow, Harvard University, USA

for their pioneering contributions to general economic equilibrium theory and **welfare theory**.

A Short Proof of Arrow Thm

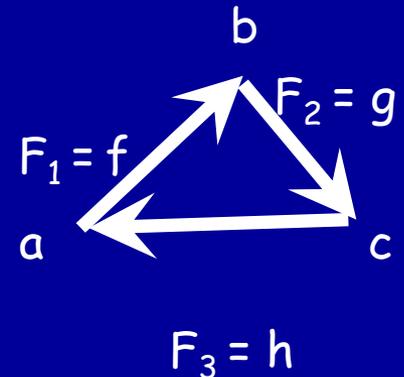
- Def: Voter 1 is pivotal for f (denoted $I_1(f) > 0$) if: $f(-, x_2, \dots, x_n) \neq f(+, x_2, \dots, x_n)$ for some x_2, \dots, x_n (similarly for other voters).
- Lemma (Barbera 82): Any constitution $F=(f, g, h)$ on 3 alternatives which satisfies **IIA** and has
 - $I_1(f) > 0$ and $I_2(g) > 0$
 - has a non-transitive outcome.
- Pf: $\exists x_2, \dots, x_n$ and y_1, y_3, \dots, y_n s.t:
 - $f(+1, +x_2, +x_3, \dots, +x_n) \neq f(-1, +x_2, +x_3, \dots, +x_n)$
 - $g(+y_1, +1, +y_3, \dots, +y_n) \neq g(+y_1, -1, +y_3, \dots, +y_n)$
 - $h(-y_1, -x_2, -x_3, \dots, -x_n) := v$ and choose x_1, y_2 s.t.: $f(x) = g(y) = v$
 \Rightarrow outcome is not transitive.
- Note: $(x_1, y_1, -y_1), (x_2, y_2, -x_2), (x_i, y_i, -x_i)$ not in $\{(1, 1, 1), (-1, -1, -1)\}$



A Short Proof of Arrow Thm

- Pf of Arrow Thm:
- Let $F = (f, g, h)$.
- Let $I(f) = \{\text{pivotal voters for } f\}$.
- Unanimity $\Rightarrow f, g, h$ are not constant
 $\Rightarrow I(f), I(g), I(h)$ are non-empty.
- By Transitivity + lemma $\Rightarrow I(f) = I(g) = I(h) = \{i\}$ for some i .
- $\Rightarrow F(\sigma) = G(\sigma_i)$
- By unanimity $\Rightarrow F(\sigma) = \sigma_i$.

- Q: How to prove for $k > 3$ alternatives?
- Q: Can we do without unanimity?



A Short Proof of Arrow Thm

- Q: How to prove for $k > 3$ alternatives?
- A: For each 3 alternatives there is a dictator so we only need to show it is the same dictator for all pairs of alternatives. If $\{a,b\}, \{c,d\}$ are two such pairs look at (a,b,c) and (b,c,d) .
- Q: Can we do without unanimity?
- A: Except the last step the same proof works if instead of unanimity we have that: for each pair of alternatives a, b in some outcome a beats b and in another b beats a .
- Then in the last step we get $F = G(\sigma_i)$
- Only such F that satisfy IIA is $F(\sigma) = \sigma$ and $F(\sigma) = -\sigma$.

A more general Arrow Theorem

- Def: Write $A \succ_F B$ if for all σ and all $a \in A$ and $b \in B$ it holds that $F(\sigma)$ ranks a above b .
- Thm (Wilson 72 as stated in M'10): A constitution F on k alternatives satisfies **IIA** and **Transitivity** iff
- F satisfies that there exists a partition of the k alternatives into sets A_1, \dots, A_s s.t:
- $A_1 \succ_F \dots \succ_F A_s$ and
- If $|A_r| > 2$ then F restricted to A_r is a dictator on some voter j .
- Note: "Dictator" now is also $F(\sigma) = -\sigma$.
- Def: Let $F_k(n) :=$ The set of constitutions on n voters and k alternatives satisfying IIA and Transitivity.

Pf of Wilson's Theorem

- Def: Write $A \succ_F B$ if for all σ and all $a \in A$ and $b \in B$ it holds that $F(\sigma)$ ranks a above b .
- Thm (Wilson 72 as stated in M'10): A constitution F on k alternatives satisfies **IIA** and **Transitivity** iff
- F satisfies that there exists a partition of the k alternatives into sets A_1, \dots, A_s s.t:
- $A_1 \succ_F \dots \succ_F A_s$ and
- If $|A_r| > 2$ then F restricted to A_r is a dictator on some voter j .
- Note: every function as above is IIA and transitive, so need to show that if f is IIA and transitive then satisfies the conditions above.

Pf of Wilson's Theorem

- Assume F is transitive and IIA.
- For two alternatives a, b write $a \succ_F b$ if a is always ranked above b . Write $a \sim_F b$ if there are outcome where $a \succ b$ and outcome where $b \succ a$.
- Claim: \succ_F is transitive.
- Claim: If there exists a profile σ where $a \succ b$ and a profile τ where $b \succ c$ then there exists an outcome where $a \succ c$.
- Pf: As in Barbera pf look at the configuration with a, b preferences taken from σ and b, c preferences taken from τ .
- Claim: \sim_F is transitive moreover if $a \succ_F b$ and $a \sim_F c$ and $b \sim_F d$ then $c \succ_F d$.

Pf of Wilson's Theorem

- Claim: \succ_F is transitive.
- Claim: \sim_F is transitive moreover if $a \succ_F b$ and $a \sim_F c$ and $b \sim_F d$ then $c \succ_F d$
- Claim: There exists a partition of the alternatives $A_1 \succ_F A_2 \succ_F \dots \succ_F A_s$
- Pf of Wilson's theorem: Apply Arrow thm to each of the A_i 's.

Ties

- Note: So far we assumed that each voters provides a strict ranking.
- Arrow and other work considered the more general case where voters are allowed to have a ranking with ties such as:
- $a > b \sim c$ or $a \sim b > c$ etc.
- Under this condition one can state Arrow's and Wilson's theorems but only one sided versions:
- Arrow theorem with ties:
- If F satisfies unanimity, IIA and transitivity then it is a dictator or null where
- Def: Dictator is a voter whose strict preferences are followed.

Some Examples of dictators

- Example 1: $F(\sigma) = \sigma_1$.
- Example 2: All the strict inequalities of σ_1 are followed and:
 - for every pair of alternatives $a \sim b$ in σ_1 run a majority vote on the pairwise preferences between a and b .
- Note:
 - Example 1 satisfies IIA while example 2 doesn't.
 - If and only if characterization in M-Tamuz-11.

Random Ranking:



- Assume uniform voting



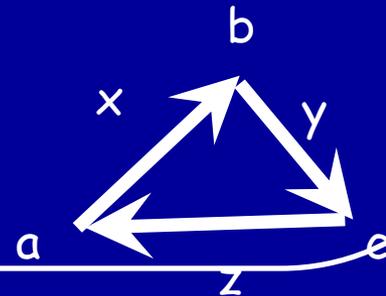
- Note: Rankings are chosen **uniformly** in S_3^n
- Assume IIA: $F(\sigma) = (f(x), g(y), h(z))$
- Q: What is the probability of a **paradox**:
- Def: $PDX(F) = P[f(x) = g(y) = h(z)]?$
- Arrow Theorem implies: If $F \neq$ dictator and f, g, h are non-constant then: $PDX(f) \geq 6^{-n}$.
- Notation: Write $D(F, G) = P(F(\sigma) \neq G(\sigma))$.
- Q: Suppose F is low influence or transitive and fair - what is the lowest possible probability of paradox?



Paradoxes and Stability

- Lemma 1 (Kalai 02):
- $PDX(F) = \frac{1}{4} (1 + E[f(x)g(y)] + E[f(x)h(z)] + E[g(y)h(z)])$
- Pf: Look at $s : \{-1,1\}^3 \rightarrow \{0,1\}$ which is 1 on $(1,1,1)$ and $(-1,-1,-1)$ and 0 elsewhere. Then
- $s(a,b,c) = \frac{1}{4} (1+ab+ac+bc)$. ■

- Note that (X,Y) is distributed as:
- $E[X_i] = E[Y_i] = 0$ and $E[X_i Y_i] = -1/3$.
- If F is fair then f,g,h are fair and we can write:
- $PDX(F) = \frac{1}{4} (1 - E[f(x)g(y)] - E[f(x)h(z)] - E[g(y)h(z)])$
- Where now (X,Y) is distributed as:
- $E[X_i] = E[Y_i] = 0$ and $E[X_i Y_i] = +1/3$



Paradoxes and Stability

- $PDX(F) = \frac{1}{4} (1 - E[f(x)g(y)] - E[f(x)h(z)] - E[g(y)h(z)])$
- Where now (X, Y) is distributed as:
- $E[X_i] = E[Y_i] = 0$ and $E[X_i Y_i] = +1/3$
- Fairness implies $E[f] = E[g] = E[h] = 0$.
- By majority is stablest
 $E[f(x)g(y)] < E[m_n(x) m_n(y)] + \epsilon$.
- Thm(Kalai 02): If F is fair and of max influence at most δ or transitive then:
- $PDX(F) > \lim PDX(Maj_n) - \epsilon$ where $\epsilon \rightarrow 0$ as $(\delta \rightarrow 0 / n \rightarrow \infty)$

Probability of a Paradox

- We already know that we cannot avoid paradoxes for low influence functions.
- Q: Can we avoid paradoxes with good probability for any non-dictatorial function?

Probability of a Paradox

- We already know that we cannot avoid paradoxes for low influence functions.
- Q: Can we avoid paradoxes with good probability with any non-dictatorial function?
- Let $f=g=h$ where $f(x) = x_1$ unless $x_2 = \dots = x_n$ in which case $f(x) = x_2$.
- Non-dictatorial system.
- Paradox probability is exponentially small.
- Q (more reasonable): Is it true that the only functions with small paradox probability are close to dictator?

Probability of a Paradox

- Kalai-02: If IIA holds with $F = (f, g, h)$ and
- $E[f] = E[g] = E[h] = 0$ then
- $PDX(F) < \varepsilon \Rightarrow \exists$ a dictator i s.t.:
- $D(F, \sigma_i) < K \varepsilon$ or $D(F, -\sigma_i) < K \varepsilon$
- Where K is some absolute constant.
- Keller-08: Same result for symmetric distributions.

Probability of a Paradox

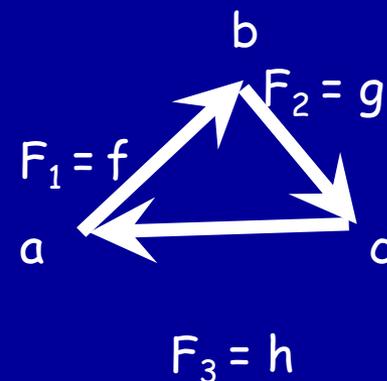
- Thm M-10: $\forall \varepsilon, \exists \delta$ s.t.:
- If IIA holds with $F = (f, g, h)$ and
- $\max \{|E[f]|, |E[g]|, |E[h]|\} < 1 - \varepsilon$ and
- $\min_i \min \{D(F, \sigma_i), D(F, -\sigma_i)\} > \varepsilon$
- Then $P(F) > \delta$.

- General Thm M-10: $\forall k, \varepsilon \exists \delta$ s.t.:
- If IIA holds for F on k alternatives and
- $\min \{D(F, G) : G \in F_k(n)\} > \varepsilon$
- Then: $P(F) > \delta$.

- Comment: Can take $\delta = k^{-2} \exp(-C/\varepsilon^{21})$

A Quantitative Lemma

- Def: The influence of voter 1 on f (denoted $I_1(f)$) is:
- $I_1(f) := P[f(-, x_2, \dots, x_n) \neq f(+, x_2, \dots, x_n)]$
- Lemma (M-09): Any constitution $F=(f, g, h)$ on 3 alternatives which satisfies **IIA** and has
- $I_1(f) > \varepsilon$ and $I_2(g) > \varepsilon$
- Satisfies $PDX(F) > \varepsilon^3/36$.
- Pf:
- Let $A_f = \{x_3, \dots, x_n : 1 \text{ is pivotal for } f(*, *, x_3, \dots, x_n)\}$
- Let $B_g = \{y_3, \dots, y_n : 2 \text{ is pivotal for } g(*, *, y_3, \dots, y_n)\}$
- Then $P[A_f] > \varepsilon$ and $P[B_g] > \varepsilon$
- By "Inverse Hyper-Contraction": $P[A_f \cap B_g] > \varepsilon^3$.
- By Lemma: $PDX[F] \geq 1/36 P[A_f \cap B_g] > \varepsilon^3/36$.



Inverse Hyper Contraction

- Note: (x_i, y_i) are i.i.d. with $E(x_i, y_i) = (0, 0)$ and $E[x_i y_i] = -1/3$
- Results of C. Borell 82: \Rightarrow
- Let $f, g : \{-1, 1\}^n \rightarrow \mathbb{R}_+$ then
- $E[f(x) g(y)] \geq |f|_p |g|_q$ if $1/9 \leq (1-q)(1-p)$ and $p, q < 1$.
- In particular: taking f and g indicators obtain:
- $E[f] > \varepsilon$ and $E[g] > \varepsilon \Rightarrow E[fg] > \varepsilon^3$.
- Implications in: M-O'Donnell-Regev-Steif-Sudakov-06.

- Note: "usual" hyper-contraction gives:
- $E[f(x) g(y)] \leq |f|_p |g|_q$ for all functions if
- $(p-1)(q-1) \geq 1/9$ and $p, q > 1$.

Inverse Hyper Contraction

The Use of Swedish Technology



IKEA Store Falls Apart! Experts Blame Cheap Parts, Confusing Blueprint
From SD Headliner, Mar 25, 09.

Quantitative Arrow - 1st attempt

- Thm M-10: $\forall \varepsilon, \exists \delta$ s.t if IIA holds with $F = (f, g, h)$ &
- $\max \{|E[f]|, |E[g]|, |E[h]|\} < 1 - \varepsilon$ &
- $\min \{D(F, G) : G \in F_3(n)\} > 3\varepsilon$
- Then $PDX(F) > (\varepsilon/96n)^3$.

- Pf Sketch: Let $P_f = \{i : I_i(f) > \varepsilon n^{-1/4}\}$
- Since $\sum I_i(f) > \text{Var}[f] > \varepsilon/2$, P_f is not empty.
- If there exists $i \neq j$ with $i \in P_f$ and $j \in P_g$ then $PDX(F) > (\varepsilon/96n)^3$ by quantitative lemma.
- So assume $P_f = P_g = P_h = \{1\}$ and $P(F) < (\varepsilon/96n)^3$
- $\Rightarrow D(f, \pm x_i) \leq \varepsilon$ or $D(f, \pm 1) \leq \varepsilon$ (same for g and h)
- $\Rightarrow D(F, G) \leq 3\varepsilon$ where $G(\sigma) = G(\sigma_1)$.
- $PDX(G) \leq 3\varepsilon + (\varepsilon/96n)^3 < 1/6 \Rightarrow G \in F_3(n)$.

Quantitative Arrow - Real Proof

- Pf High Level Sketch:
- Let $P_f = \{i : I_i(f) > \varepsilon\}$.
- If there exists $i \neq j$ with $i \in P_f$ and $j \in P_g$ then $PDX(F) > \varepsilon^3 / 36$ by quantitative lemma.
- Two other cases to consider:
- I. $P_f \cap P_g = P_f \cap P_h = P_g \cap P_h$ is empty
- In this case: use Invariance + Gaussian Arrow Thm.
- II. $P_f \cup P_g \cup P_h = \{1\}$.
- In this case we condition on voter **1** so we are back in case I.

Quantitative Arrow - Real Proof

- The Low Influence Case:
- We want to prove the theorem under the condition that $P_f \cap P_g = P_f \cap P_h = P_g \cap P_h$ is empty.
- Let's first assume that $P_f = P_g = P_h$ is empty - all functions are of low influence.
- Recall:
- $PDX(F) = \frac{1}{4} (1 + E[f(x)g(y)] + E[f(x)h(z)] + E[g(y)h(z)])$
- Where now (X, Y) is distributed as:
- $E[X_i] = E[Y_i] = 0$ and $E[X_i Y_i] = -1/3$
- By a version of Maj-Stablest Majority is Stablest:
- $PFX(F) > PDX(u, v, w) + \text{error}(I)$ where
- $u(x) = \text{sgn}(\sum x_j + u_0)$ and $E[u] = E[f]$ etc.

Quantitative Arrow - Real Proof

- By Majority is Stablest:
- $PFX(F) > PDX(u,v,w) + \text{error}(I)$ where
- $u(x) = \text{sgn}(\sum x_j + u_0)$ and $E[u] = E[f]$ etc.
- Remains to bound $PDX(u,v,w)$
- By CLT this is approximately:
- $P[U>0, V>0, W>0] + P[U<0, V<0, W<0]$ where $U \sim N(E(u), 1)$, $V \sim N(E(v), 1)$ and $W \sim N(E(w), 1)$ &
- $\text{Cov}[U, V] = \text{Cov}[V, W] = \text{Cov}[W, U] = -1/3$.
- For Gaussians possible to bound.

Quantitative Arrow - Real Proof

- In fact the proof works under the weaker condition that $P_f \cap P_g = P_f \cap P_h = P_g \cap P_h$ is empty.
- The reason is that the strong version of majority is stablest (M-10) says:
- If $\min(I_i(f), I_i(g)) < \delta$ for all i and u and v are majority functions with $E[f]=u$, $E[g]=v$ then:
- $E[f(X)g(Y)] < \lim_n E[u_n(X)v_n(Y)] + \varepsilon(\delta)$ where
- $\varepsilon(\delta) \rightarrow 0$ as $\delta \rightarrow 0$.

Probability of a Paradox for Low Inf Functions

- Thm: (Follows from MOO-05): $\forall \epsilon > 0 \exists \delta > 0$ s.t. If
- $\max_i \max\{I_i(f), I_i(g), I_i(h)\} < \delta$
then $PDX(F) > \lim_{n \rightarrow \infty} PDX(f_n, g_n, h_n) - \epsilon$
- where $f_n = \text{sgn}(\sum_{i=1}^n x_i - a_n)$, $g_n = \text{sgn}(\sum_{i=1}^n y_i - b_n)$, $h_n = \text{sgn}(\sum_{i=1}^n z_i - c_n)$ and a_n, b_n and c_n are chosen so that $E[f_n] \sim E[f]$ etc.
- Thm (Follows from M-08): The same theorem holds with $\max_i 2^{\text{nd}}(I_i(f), I_i(g), I_i(h)) < \delta$.
- So case I. of quantitative Arrow follows if we can prove Arrow theorem for threshold functions.
- (Recall case I.: $P_f \cap P_g = P_f \cap P_h = P_g \cap P_h$ is empty)
- Pf for "threshold functions" using Gaussian analysis.

Pf of Majority is Stablest

- Majority is Stablest Conj: If $E[f] = E[g] = 0$ and f, g have all influences less than δ then $E[f(x)g(y)] > E[m_n(x) m_n(y)] - \epsilon$.
- Ingredients:
 - I. Thm (Borell 85): (N_i, M_i) are i.i.d. Gaussians with $E[N_i] = E[M_i] = 0$ and $E[N_i M_i] = -1/3$, $E[N_i^2] = E[M_i^2] = 1$ and f and g are two functions from \mathbb{R}^n to $\{-1, 1\}$ with $E[f] = E[g] = 0$ then:
 - $E[f(X) g(Y)] \geq E[\text{sgn}(X_1) \text{sgn}(Y_1)]$.
 - By the CLT: $E[\text{sgn}(X_1) \text{sgn}(Y_1)] = \lim_{n \rightarrow \infty} E[m_n(x) m_n(y)]$
 - II. Invariance Principle [M+O'Donnell+Oleszkiewicz(05)]:
 - Gaussian case \Rightarrow Discrete case.

The Geometry Behind Borell's Result

- I. Thm (Borell 85): (N_i, M_i) are i.i.d. Gaussians with
- $E[N_i] = E[M_i] = 0$ and $E[N_i M_i] = -1/3$, $E[N_i^2] = E[M_i^2] = 1$ and f and g are two functions from \mathbb{R}^n to $\{-1, 1\}$ with $E[f] = E[g] = 0$ then:
- $E[f(X) g(Y)] \geq E[\text{sgn}(X_1) \text{sgn}(Y_1)]$.

- Spherical Version: Consider $X \in S^n$ uniform and $Y \in S^n$ chosen uniformly conditioned on $\langle X, Y \rangle \leq -1/3$.
- Among functions f, g with $E[f] = E[g] = 0$ what is the minimum of $E[f(X) g(Y)]$?
- Answer: $f = g =$ same half-space.

The Geometry Behind Borell's Result

- More general Thm (Isaksson-M 09): (N^1, \dots, N^k) are k n -dim Gaussian vectors $N^i \sim N(0, I)$.
- $\text{Cov}(N^i, N^j) = \rho I$ for $i \neq j$, where $\rho > 0$.
- Then if f_1, \dots, f_k are functions from \mathbb{R}^n to $\{0, 1\}$ with $E[f] = 0$ then:
- $E[f_1(N^1) \dots f_k(N^k)] \leq E[\text{sgn}(N^1_1) \dots \text{sgn}(N^k_1)]$
- Proof is based on re-arrangements inequalities on the sphere.
- Gives that majority maximizes probability of unique winner in Condorcet voting for low influence functions.

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- More general Thm (Isaksson-M 09): (N^1, \dots, N^k) are k n -dim Gaussian vectors $N^i \sim N(0, I)$.
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- Then if f_1, \dots, f_k are functions from \mathbb{R}^n to $\{0, 1\}$ with $E[f] = 0$ then:
- $E[f_1(N^1) \dots f_k(N^k)] \leq E[\text{sgn}(N^1_1) \dots \text{sgn}(N^k_1)]$
- Proof is based on re-arrangements inequalities on the sphere.
- Gives that majority maximizes probability of unique winner in Condorcet voting for low influence functions.

HW 1

- Let $f=g=h$ be the $m \times m$ electoral college and consider IIA vote with $F=(f,g,h)$.
- Given a uniform vote x and y obtained from x by a single uniformly chosen voting error, what is $\lim P[f(x) \neq f(y)] \times m$ as $m \rightarrow \infty$.
- Assume x is obtained from y by flipping each coordinate with probability ε independently.
What is $\lim P[f(x) \neq f(y)]$ as $m \rightarrow \infty$
- What is the limiting probability of an Arrow paradox assuming uniform voting and $m \rightarrow \infty$?

HW 2

- Consider the function $\Psi(f, i)$ which given
- a function $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$ and a voter i returns an x s.t. f is pivotal on x and voter i . The function returns **Null** if no such x exist.
- Given access to $\Psi(f, ?)$ $\Psi(g, ?)$ and $\Psi(h, ?)$ Design an efficient algorithm that decides if (f, g, h) has a non-transitive outcome and if such an outcome exist it produces it. The running time of the algorithm should be linear in n .

HW 2 - continued

- Assume that the functions f, g and h are monotone and submodular so that for all x, y :
- $f(\min(x, y)) + f(\max(x, y)) \leq f(x) + f(y)$

where the maximum is taken coordinate-wise.

Show that the problems of deciding if all outcome of (f, g, h) are transitive and finding a non-transitive outcomes if such exist can both be solve in linear time (assuming access to f, g and h takes one unit of time)

HW 3

- Consider the 3-reursive majority functions f_n :
- $f_1(x(1),x(2),x(3)) = \text{maj}(x(1),x(2),x(3))$
- $f_{k+1}(x(1),\dots,x(3^{k+1})) = \text{maj}(f_1(x),f_2(y),f_3(z))$ where
- $x = (x(1),\dots,x(3^k)), y = (x(3^{k+1}+1),\dots,x(2 \cdot 3^k)), z = \dots$
- Let (x,y) be uniform with y different from x in one coordinate. What is $P[f_k(x) = f_k(y)]$?
- Assuming x is uniform and y is obtained from x by flipping each coordinate with probability ε , show:
- $P[f_k(x) = f_k(y)] = \frac{1}{2} + 3^k \alpha + o(1)$ for some α . Find $\alpha(\varepsilon)$
- Consider ranking using $F=(f(x),f(y),f(z))$. What is the limit of $P[F(\sigma) \text{ is non-transitive}]$? What is the next order term (both as $k \rightarrow \infty$)

HW 4

- Consider the Plurality coordination problem on a social network where initially each player receives one of 3 colors.
- Design a protocol using the color and one extra bit of memory that reaches coordination.

HW 5

- Consider the voter model on $G=(V,E)$.
- Assume that the model is run for k different topics and that further
- Assume that for each topic k , time t and all $v \in V$ the opinion of v at topic k denoted $v(k,t)$ is known but:
 - The graph E is not known.
 - Design an algorithm that finds the edges of the graph G from the record of the votes.
 - How large should k and t be for the algorithm to have a high probability of recovering G ?