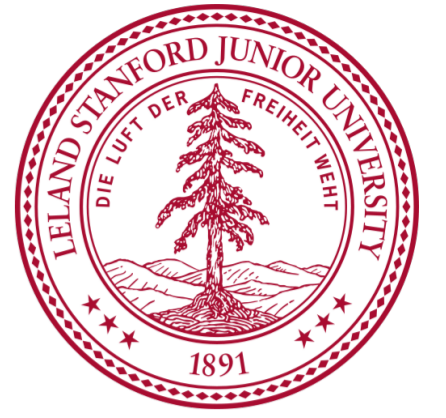


# Minimax Rates for Memory-Constrained Sparse Linear Regression



Jacob Steinhardt    John Duchi

{jsteinha, jduchi}@stanford.edu

## Resource-Constrained Learning

How do we solve statistical problems with limited resources?

- communication / memory constraints (Zhang et al., 2013; Garg et al., 2014; Shamir, 2014)
- privacy, computation constraints (Kasiviswanathan et al., 2011; Duchi et al., 2013; Berthet and Rigollet, 2013)
- NP-hardness of sparse regression (Zhang et al., 2014; Natarajan, 1995)

This work: sparse linear regression under memory constraints.

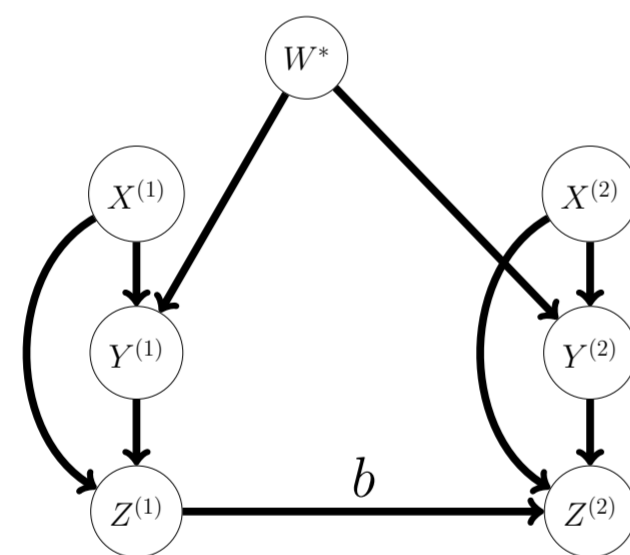
## Setting

Sparse linear regression in  $\mathbb{R}^d$ :

- $Y^{(i)} = \langle w^*, X^{(i)} \rangle + \epsilon^{(i)}$
- $\|w^*\|_0 = k, k \ll d$

Memory constraint:

- $(Y^{(i)}, X^{(i)})$  observed as read-only stream
- Only keep  $b$  bits of state  $Z^{(i)}$  between successive observations



## Problem Statement

How much data  $n$  is needed to obtain estimator  $\hat{w}$  with

$$\mathbb{E}[\|\hat{w} - w^*\|_2^2] \leq \epsilon?$$

Classical case (no memory constraint):

**Theorem** (Wainwright, 2009).

$$\frac{k}{\epsilon} \log(d) \lesssim n \lesssim \frac{k}{\epsilon} \log(d)$$

With memory constraints  $b$ :

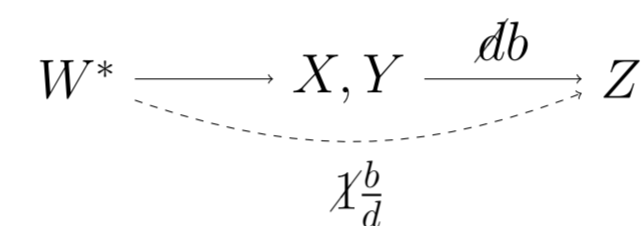
**Theorem** (S. & Duchi, 2015).

$$\frac{k d}{\epsilon b} \lesssim n \lesssim \frac{k d}{\epsilon^2 b}$$

Exponential increase if  $b \ll d$ !

## Proof Overview

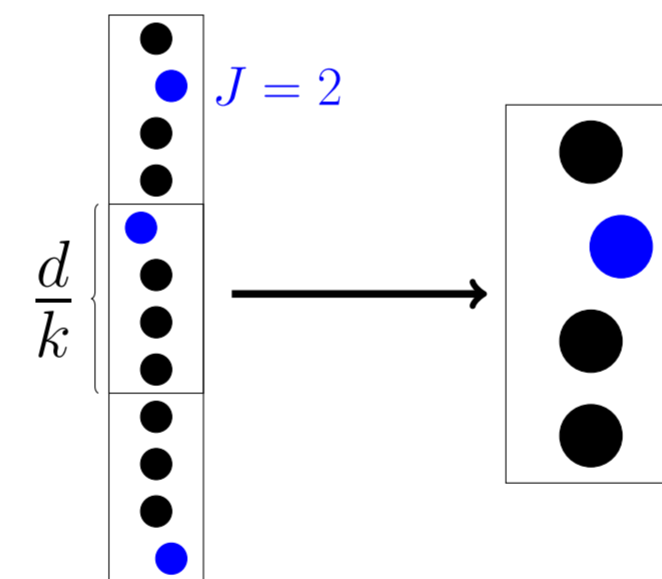
- Lower bound:
  - information-theoretic
  - strong data-processing inequality



- Upper bound:
  - count-min sketch +  $\ell^1$ -regularized dual averaging
  - more regularization  $\rightarrow$  easier sketching problem

## Lower Bound Construction

- Split coordinates into  $k$  blocks of size  $d/k$
- $w^*$  in each block: single non-zero coordinate  $J$ ,  $\pm\delta$  with equal probability
- Direct sum argument: reduce to  $k = 1$



- Estimation to testing:

$$\mathbb{E}[\|w^* - \hat{w}\|_2^2] \geq \frac{\delta^2}{2} \mathbb{P}[J \neq \hat{J}]$$

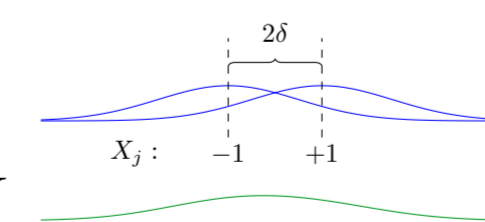
Looking ahead: bound KL between  $P_j$  and base distribution  $P_0$

## Some Information Theory

- Let  $X \sim \text{Uniform}(\{\pm 1\}^d)$
- Let  $P_j(Z^{(1:n)})$  be distribution conditioned on  $J = j$
- Let  $P_0(Z^{(1:n)})$  be distribution with  $Y$  independent of  $X$
- Assouad's method:

$$\mathbb{P}[J \neq \hat{J}] \geq \frac{1}{2} - \sqrt{\frac{1}{d} \sum_{j=1}^d \text{KL}(P_0(Z^{(1:n)}) \| P_j(Z^{(1:n)}))}$$

- Intuition:  $\text{KL}(P_0 \| P_j)$  small unless  $Z$  stores info about  $X_j$



## Strong Data-Processing Inequality

Focus on a single index  $Z = Z^{(i)}$ , with  $\hat{z} = z^{(1:i-1)}$  fixed.

**Proposition.** For any  $\hat{z}$ ,

$$\begin{aligned} \text{KL}(P_0(Z | \hat{z}) \| P_j(Z | \hat{z})) &\leq 4\delta^2 I(X_j; Z | Y, \hat{Z} = \hat{z}) \\ &\leq 4\delta^2 \underbrace{I(X_j; Z, Y | \hat{Z} = \hat{z})}_{\text{mutual information}} \end{aligned}$$

Plug into Assouad:

$$\begin{aligned} \frac{1}{d} \sum_{j=1}^d \text{KL}(P_0 \| P_j) &\leq \frac{4\delta^2}{d} \sum_{j=1}^d I(X_j; Z, Y | \hat{Z}) \\ &\leq \frac{4\delta^2}{d} \underbrace{I(X; Z, Y | \hat{Z})}_{b+O(1)} \end{aligned}$$

Only get  $\frac{4\delta^2 b}{d}$  bits per round!

## Upper Bound

Solve  $\ell^1$ -regularized dual averaging problem (Xiao, 2010),  $\lambda \gg 1$ :

$$\begin{aligned} w^{(i)} &= \underset{w}{\text{argmin}} \left\{ \langle \theta^{(i)}, w \rangle + \lambda \sqrt{n} \|w\|_1 \right\}, \\ \theta^{(i)} &= \sum_{i'=1}^{i-1} x^{(i')} (y^{(i')} - \langle w^{(i')}, x^{(i')} \rangle). \end{aligned}$$

Hard part: determine support of  $w^{(i)}$ .

- Need to distinguish  $|\theta_j| \geq \lambda \sqrt{n}$  (signal) from  $|\theta_j| \approx \sqrt{n}$  (noise)
- Can use count-min sketch, memory usage  $\approx \frac{d \log(d)}{\epsilon^2}$   
 $\implies$  computational-statistical tradeoff; seen before in  $\ell^2$  case (Shalev-Shwartz & Zhang, 2013; Bruer et al., 2014)

## Discussion

Summary:

- Upper and lower bounds on memory-constrained regression
- Lower bound: extend data processing inequality to handle covariates
- Upper bound: use  $\ell^1$ -regularizer to reduce to sketching

Future work:

- Close the gap ( $kd/b\epsilon$  vs  $kd/b\epsilon^2$ )
- Weaken upper bound assumptions

The first author was supported by the Hertz foundation and by the NSF.