## 1. Local weak convergence of graphs/networks

- Stuff that's obvious when you think about it
- 4 non-obvious examples/results

2. The core idea in our probabilistic reformulation of special cases of the cavity method is: do exact calculation on some infinite random graph (tree-like, in practice). LWC provides link with the finite-n problem. Illustrate with

- mean-field TSP
- flow through a disordered network.

Aldous-Steele survey "The objective method ..." on my home page.

## Some math infrastructure

Consider an abstract space $S$ (complete separable metric space) with a notion of convergence $x_{n} \rightarrow x$. There is automatically a notion of convergence of probability measures on $S$ (all reasonable definitions are equivalent).
$\mu_{n} \rightarrow \mu_{\infty}$ iff there exist $S$-valued random variables $X_{n}$ such that

$$
\operatorname{dist}\left(X_{n}\right)=\mu_{n} ; \quad P\left(X_{n} \rightarrow X_{\infty}\right)=1
$$

This is called weak convergence.

Conceptual point: When you consider some new abstract space $S$, you don't need to think about what convergence of distributions means.

Most concrete case is $S=R^{1}$, where we have e.g. the central limit theorem

$$
n^{-1 / 2} \sum_{i=1}^{n} \xi_{i} \xrightarrow{d} \operatorname{Normal}(0,1)
$$

for i.i.d. $\left(\xi_{i}\right)$ with $E \xi=0$ and $\operatorname{var} \xi=1$.
Best known abstract case is

$$
S=\{\text { continuous functions }[0,1] \rightarrow R\}
$$

which allows one to formalize "rescaled random walk converges to Brownian motion".


Another abstract case is

$$
S=\left\{\text { locally finite point sets in } R^{2}\right\}
$$

which allows one to formalize " $n$ uniform random points in square of area $n$ converges to the Poisson point process on $R^{2 \prime \prime}$.


A graph has vertices and edges
A network is a graph whose edges have postive real lengths (default length $=1$ ) and maybe extra structure indicated by numbers/labels on vertices/edges. Write $G$ for a network.

Consider the abstract space

$$
S=\{\text { locally finite rooted networks }\} .
$$

What should convergence $G_{n} \rightarrow G_{\infty}$ mean? Note: interesting case is where $G_{n}$ is finite and $G_{\infty}$ is infinite.

Window of radius $r$ in $G$ defines subgraph $G[r]$ of vertices within distance $r$ from root, with edges both of whose endpoints are within window (a convention which turns out convenient).

Definition: $G_{n} \rightarrow G_{\infty}$ means that for each fixed generic $0<r<\infty$, for large $n$ there is graph-isomorphism between $G_{n}[r]$ and $G_{\infty}[r]$ such that edge-lengths of isomorphic edges converge as $n \rightarrow \infty$ (and also other labels converge).

Given $n$-vertex network (deterministic or random) let $U_{n}$ be uniform random vertex. Write $G_{n}\left[U_{n}\right]$ for $G_{n}$ rooted at $U_{n}$.

Definition. If $G_{n}\left[U_{n}\right] \xrightarrow{d}$ some $G_{\infty}$, call this local weak convergence (LWC) of $G_{n}$ to $G_{\infty}$.

Formalizes the idea: for large $n$ the local structure of $G_{n}$ near a typical vertex is approximately the local structure of $G_{\infty}$ near the root.

Intuition: in models where degree distribution is bounded in probability as $n \rightarrow \infty$ we expect LWC to some limit infinite network.

## Obvious examples

1a $G_{n}$ : geometric graph (all edges of length $\leq c$ ) on $n$ random points in square of area $n$ $G_{\infty}$ : geometric graph on Poisson point process (rate 1) on $R^{2}$ with point at origin.

1b As above with complete graphs.
$2 G_{n}$ : discrete cube $C_{m}^{d} \subset Z^{d}$ with i.i.d. (independent random) edge-lengths.
$G_{\infty}$; all $Z^{d}$ with i.i.d. edge-lengths.
3a $G_{n}$ : Erdos-Renyi random graph $\mathcal{G}(n, c / n)$
$G_{\infty}$ : tree of Galton-Watson branching process with Poisson(c) offspring.

3b $G_{n}$ : random r-regular graph $G_{\infty}$ : infinite degree- $r$ tree.

3c $G_{n}$ : random graph model designed as "random subject to degree distribution approximately a prescribed distribution ( $p(i), i \geq 0$ )
$G_{\infty}$ : Galton-Watson tree with offspring distribution $\tilde{p}(i) \propto(i+1) p(i+1)$
$4 G_{n}$ : de Bruijn graph on $n=2^{b}$ binary strings $G_{\infty}$ : infinite tree with in-degree 2 and outdegree 2.

| 001100 |  | 110010 | 100100 |
| :--- | :--- | :--- | :--- |
|  | 011001 |  | 100101 |
| 101100 |  | 110011 | 100110 |
|  |  |  | 100111 |

$5 G_{n}$ : Simple random walk with $n$ steps $G_{\infty}$ : 2-sided infinite simple RW.
(represent as linear graph with edge-marks $\pm 1$ ).

6a-z: For many models of random $n$-vertex trees one can explicitly describe $G_{\infty}$. For instance
$G_{n}$ : uniform random tree on $n$ labeled vertices $G_{\infty}$ : infinite path from root; i.i.d. "bushes" are Galton-Watson trees with Poisson(1) offspring.


Remark: qualitative behavior similar in most models: semi-infinite path with i.i.d. finite bushes, whose mean size is infinite. Bush at root gives limit of subtree defined by random vertex in original rooted tree.
$7 G_{n}$ : complete graph on $n$ vertices; edge-lengths random, independent Exponential(mean $n$ ) distribution. $G_{\infty}$ : the PWIT (Poisson weighted infinite tree)


Distances $0<\xi_{1}<\xi_{2}<\xi_{3}<\ldots$ from a vertex to its near neighbors (indicated by lines) are successive points of a Poisson (rate 1) process on ( $0, \infty$ ). Continue recursively.

## Q: So what's the use of knowing LWC . . . ?

A: not much, but it's a start .......

Let's mention 4 results/examples not related to cavity method.

Result A: According to the graph-theoretic definition of planar graph, the infinite binary tree is a planar graph. but this seems silly to a probabilist, because probabilistic models (random walk, percolation, interacting particles) behave quite differently on trees than on $Z^{2}$. The class of random networks defined as
(*)LWC limits of finite random planar graphs provides a more natural formalization of "random infinite planar graphs". Benjamini-Schramm (2001) show that on graphs (*) with bounded degree, RW is recurrent. Suggests many other questions........

Result B: Particular models of random planar $n$-vertex graphs include

- uniform random triangulations (Angel-Schramm 2003)
- uniform random quadrangulations (ChassaingSchaeffer 2004).

In each case there is a LWC limit which may be called the uniform infinite planar triangulation/quadrangulation.
xxx pictures

Result C: Because of the 'uniform random rooting" in the definition

Definition. If $G_{n}\left[U_{n}\right] \xrightarrow{d}$ some $G_{\infty}$, call this local weak convergence (LWC) of $G_{n}$ to $G_{\infty}$
a random infinite network $G_{\infty}$ which is a LWC limit is not entirely arbitrary, but has a property interpretable as "each vertex is equally likely to be the root" (stationary or involution invariant or unimodular).

Not obvious (but true: Aldous-Lyons, in preparation) that any random infinite network with this property really is some LWC limit.
(This is technically useful in extending obvious results in finite setting to the infinite setting)

Result D: A tractable complex network model (Aldous 2003/4) designed to have a LWC limit within which explicit formulas can be calculated (giving $n \rightarrow \infty$ asymptotics for finite- $n$ models).


Q: So what's the use of knowing LWC ...?
One goal is to prove that solution of CO problem on $G_{n}$ converges to solution of CO problem on $G_{\infty}$. Not always true, of course!

Example E: Suppose edge-lengths are distinct. Then $G_{n}$ has a unique MST (minimum spanning tree). Also we can define the (wired) minimum spanning forest (MSF) on an infinite network $G_{\infty}$.

Lemma: If $G_{n} \rightarrow G_{\infty}$ (LWC), if (technical condition on $G_{\infty}$ ), then

$$
\left(G_{n}, \operatorname{MST}\left(G_{n}\right)\right) \rightarrow\left(G_{\infty}, \operatorname{MSF}\left(G_{\infty}\right)\right)(\operatorname{LWC})
$$

In particular, on the PWIT one can calculate

$$
\begin{gathered}
E(\text { length of MSF per vertex }) \\
=\frac{1}{2} E(\text { length of MSF-edges at root) } \\
=\zeta(3)=\sum_{j=1}^{\infty} j^{-3}
\end{gathered}
$$

and re-derive result of Frieze (1985) that in complete graph with random edge lengths model

$$
n^{-1} E(\text { length of MST }) \rightarrow \zeta(3)
$$

## xXx java picture

Example F: (in Aldous-Steele survey). Uniform random tree on $n$ vertices. Put i.i.d. positive weights on edges.
$M_{n}:=$ weight of max-weight partial matching Then $n^{-1} E M_{n} \rightarrow E M_{\infty}$ where

$$
M_{\infty}=\frac{1}{2} \text { (weight of edge at root) }
$$

in max-weight matching on limit infinite tree


Explicitly, for Exponential(1) edge-weights we get $E M_{\infty} \approx 0.2396$ where the limit equals $\int_{0}^{\infty} s e^{-s} d s \int_{0}^{s} c\left(e^{-y}-b e^{-s}\right) \exp \left(-c e^{-y}-c e^{-(s-y)}\right) d y$ where $c \approx 0.7146$ is the strictly positive solution of $c^{2}+e^{-c}=1$ and $b=\frac{c^{2}}{c^{2}+2 c-1} \approx 0.5433$.

## LWC and the cavity method

## Recall model

$G_{n}$ : complete graph on $n$ vertices; edge-lengths random, independent Exponential(mean $n$ ) distribution.

Write $L_{n}$ for length of TSP tour under this model. Mezard-Parisi (1980s) used replica/cavity methods to argue

$$
n^{-1} E L_{n} \rightarrow c \approx 2.04
$$

We have explicit program to make rigorous but can't carry through two of the technical steps.

Here are two relaxations of TSP for $n$-vertex network.
(M2F): minimum 2-factor. Minimize total length of a 2 -factor, that is an edge-set in which each vertex has degree 2 . That is, a union of cycles which spans all vertices.
(MA2F): minimum almost 2-factor. Minimize total length over edge-sets $\mathcal{E}_{n}$ such that

$$
n^{-1} \mid\{v: \text { degree }(v) \neq 2\} \mid \rightarrow 0 .
$$

Proposition (Frieze 2004) In this model,

$$
n^{-1}\left(E L_{n}-E L_{n}^{\prime}\right) \rightarrow 0
$$

where $L_{n}$ is TSP length and $L_{n}^{\prime}$ is M2F length.
Missing Proposition Want to know

$$
n^{-1}\left(E L_{n}-E L_{n}^{\prime \text { prime }}\right) \rightarrow 0
$$

where $L_{n}$ is TSP length and $L_{n}^{\text {'prime }}$ is MA2F length.

## xxx java slide; hand write Zs

Central part of method - which I'll explain only superficially - is to do analysis of TSP on the (infinite) PWIT. each edge $e$ of PWIT splits it into two subtrees. There are random variables $Z^{1}(e), Z^{2}(e)$, measurable functions of the subtrees, such that
$e \in$ TSP-path iff length $(e)<Z^{1}(e)+Z^{2}(e)$
(another "missing proposition" in proving this) from which one can calculate mean length of TSP-path edges.

Q: How do we go back from the PWIT to the finite-n model?

More math infrastructure

A measurable function $f\left(Y_{1}, Y_{2}, \ldots\right)$ of some infinite collection of r.v.'s can be approximated arbitrary closely by continuous functions $f_{k}\left(Y_{1}, \ldots, Y_{k}\right)$ of finitely many of the r.v.'s.

So on the PWIT we can define an edge-set $\mathcal{E}_{r}$ such that
(i) the edges of $\mathcal{E}_{r}$ at a vertex $v$ are determined by the restriction of the PWIT to the window of radius $r$ around $v$; (ii) $\delta_{r}:=P$ ( some edge at $v \in \mathcal{E}_{r} \triangle\{ ) \rightarrow 0$ as $r \rightarrow \infty$.

Using LWC of the finite- $n$ model to the PWIT, we can apply the same rule to a window of radius $r$ around a vertex $v$, and define edgesets $\mathcal{E}_{r, n}$ such that
$\limsup _{n} P\left(\right.$ degree $(v)$ in $\left.\mathcal{E}_{r, n} \neq 2\right) \leq 2 \delta_{r}$ xxx and similarly the edge-lengths xxx .

This constructs an almost-2-factor of $G_{n}$ whose cost-per-vertex converges to the $c$ given in the PWIT analysis.

