### 1. Local weak convergence of graphs/networks

- Stuff that's obvious when you think about it
- 4 non-obvious examples/results

2. The core idea in our probabilistic reformulation of special cases of the cavity method is: do exact calculation on some infinite random graph (tree-like, in practice). LWC provides link with the finite-n problem. Illustrate with

- mean-field TSP
- flow through a disordered network.

Aldous-Steele survey "The objective method ...." on my home page.

# Some math infrastructure

Consider an abstract space S (complete separable metric space) with a notion of convergence  $x_n \rightarrow x$ . There is automatically a notion of convergence of probability measures on S (all reasonable definitions are equivalent).

 $\mu_n \to \mu_\infty$  iff there exist S-valued random variables  $X_n$  such that

dist
$$(X_n) = \mu_n$$
;  $P(X_n \to X_\infty) = 1$ .

This is called weak convergence.

**Conceptual point:** When you consider some new abstract space S, you don't need to think about what convergence of distributions means.

Most concrete case is  $S = R^1$ , where we have e.g. the central limit theorem

$$n^{-1/2} \sum_{i=1}^{n} \xi_i \stackrel{d}{\rightarrow} \operatorname{Normal}(0,1)$$

for i.i.d.  $(\xi_i)$  with  $E\xi = 0$  and var  $\xi = 1$ .

Best known abstract case is

 $S = \{$ continuous functions  $[0, 1] \rightarrow R \}$ which allows one to formalize "rescaled random walk converges to Brownian motion".



Another abstract case is

 $S = \{ \text{locally finite point sets in } R^2 \}$ which allows one to formalize "*n* uniform random points in square of area *n* converges to the Poisson point process on  $R^2$ ".



A graph has vertices and edges A network is a graph whose edges have postive real lengths (default length = 1) and maybe extra structure indicated by numbers/labels on vertices/edges. Write G for a network.

Consider the abstract space

 $S = \{ \text{locally finite rooted networks} \}.$ 

What should convergence  $G_n \rightarrow G_\infty$  mean? <u>Note</u>: interesting case is where  $G_n$  is finite and  $G_\infty$  is infinite.

Window of radius r in G defines subgraph G[r] of vertices within distance r from root, with edges **both** of whose endpoints are within window (a convention which turns out convenient).

**Definition:**  $G_n \to G_\infty$  means that for each fixed generic  $0 < r < \infty$ , for large n there is graph-isomorphism between  $G_n[r]$  and  $G_\infty[r]$  such that edge-lengths of isomorphic edges converge as  $n \to \infty$  (and also other labels converge).

Given *n*-vertex network (deterministic or random) let  $U_n$  be uniform random vertex. Write  $G_n[U_n]$  for  $G_n$  rooted at  $U_n$ .

**Definition.** If  $G_n[U_n] \xrightarrow{d}$  some  $G_\infty$ , call this **local weak convergence (LWC)** of  $G_n$  to  $G_\infty$ .

Formalizes the idea: for large n the local structure of  $G_n$  near a typical vertex is approximately the local structure of  $G_\infty$  near the root.

**Intuition:** in models where degree distribution is bounded in probability as  $n \to \infty$  we expect LWC to some limit infinite network.

### **Obvious examples**

**1a**  $G_n$ : geometric graph (all edges of length  $\leq c$ ) on n random points in square of area n

 $G_{\infty}$ : geometric graph on Poisson point process (rate 1) on  $R^2$  with point at origin.

**1b** As above with complete graphs.

**2**  $G_n$ : discrete cube  $C_m^d \subset Z^d$  with i.i.d. (independent random) edge-lengths.

 $G_{\infty}$ ; all  $Z^d$  with i.i.d. edge-lengths.

**3a**  $G_n$ : Erdos-Renyi random graph  $\mathcal{G}(n, c/n)$  $G_\infty$ : tree of Galton-Watson branching process with Poisson(c) offspring.

**3b**  $G_n$ : random r-regular graph  $G_\infty$ : infinite degree-r tree.

**3c**  $G_n$ : random graph model designed as "random subject to degree distribution approximately a prescribed distribution  $(p(i), i \ge 0)$ 

 $G_{\infty}$ : Galton-Watson tree with offspring distribution  $\tilde{p}(i) \propto (i+1)p(i+1)$ 

**4**  $G_n$ : de Bruijn graph on  $n = 2^b$  binary strings  $G_\infty$ : infinite tree with in-degree 2 and out-degree 2.

001100 101100	011001	110010	100100
			100101
		110011	100110
			100111

**5**  $G_n$ : Simple random walk with n steps  $G_\infty$ : 2-sided infinite simple RW. (represent as linear graph with edge-marks  $\pm 1$ ).

**6a-z**: For many models of random *n*-vertex trees one can explicitly describe  $G_{\infty}$ . For instance

 $G_n$ : uniform random tree on *n* labeled vertices  $G_\infty$ : infinite path from root; i.i.d. "bushes" are Galton-Watson trees with Poisson(1) off-spring.



<u>Remark</u>: qualitative behavior similar in most models: semi-infinite path with i.i.d. finite bushes, whose mean size is infinite. Bush at root gives limit of subtree defined by random vertex in original rooted tree. **7**  $G_n$ : complete graph on n vertices; edge-lengths random, independent Exponential(mean n) distribution.

 $G_{\infty}$ : the PWIT (Poisson weighted infinite tree)



Distances  $0 < \xi_1 < \xi_2 < \xi_3 < \dots$  from a vertex to its **near neighbors** (indicated by lines) are successive points of a Poisson (rate 1) process on  $(0, \infty)$ . Continue recursively.

## Q: So what's the use of knowing LWC ...?

A: not much, but it's a start .....

Let's mention 4 results/examples <u>not</u> related to cavity method.

**Result A:** According to the graph-theoretic definition of **planar graph**, the infinite binary tree is a planar graph. but this seems silly to a probabilist, because probabilistic models (random walk, percolation, interacting particles) behave quite differently on trees than on  $Z^2$ . The class of random networks defined as

(\*)LWC limits of finite random planar graphs

provides a more natural formalization of "random infinite planar graphs". Benjamini-Schramm (2001) show that on graphs (\*) with bounded degree, RW is recurrent. Suggests many other questions ..... **Result B:** Particular models of random planar *n*-vertex graphs include

- uniform random triangulations (Angel-Schramm 2003)
- uniform random quadrangulations (Chassaing-Schaeffer 2004).

In each case there is a LWC limit which may be called the **uniform infinite planar trian**gulation/quadrangulation.

xxx pictures

**Result C:** Because of the 'uniform random rooting' in the definition

**Definition.** If  $G_n[U_n] \xrightarrow{d}$  some  $G_\infty$ , call this **local weak convergence (LWC)** of  $G_n$  to  $G_\infty$ 

a random infinite network  $G_{\infty}$  which is a LWC limit is not entirely arbitrary, but has a property interpretable as "each vertex is equally likely to be the root" (stationary or involution invariant or unimodular).

Not obvious (but true: Aldous-Lyons, in preparation) that any random infinite network with this property really is some LWC limit.

(This is technically useful in extending obvious results in finite setting to the infinite setting)

**Result D:** A tractable complex network model (Aldous 2003/4) designed to have a LWC limit within which explicit formulas can be calculated (giving  $n \rightarrow \infty$  asymptotics for finite-n models).



#### Q: So what's the use of knowing LWC ...?

One goal is to prove that solution of CO problem on  $G_n$  converges to solution of CO problem on  $G_\infty$ . Not always true, of course!

**Example E:** Suppose edge-lengths are distinct. Then  $G_n$  has a unique MST (minimum spanning tree). Also we can define the (wired) minimum spanning forest (MSF) on an infinite network  $G_{\infty}$ .

**Lemma:** If  $G_n \to G_\infty$  (LWC), if (technical condition on  $G_\infty$ ), then

$$(G_n, \mathsf{MST}(G_n)) \to (G_\infty, \mathsf{MSF}(G_\infty))$$
 (LWC).

In particular, on the PWIT one can calculate

*E*(length of MSF per vertex)

 $=\frac{1}{2}E(\text{length of MSF-edges at root})$ 

$$=\zeta(3)=\sum_{j=1}^{\infty}j^{-3}$$

and re-derive result of Frieze (1985) that in complete graph with random edge lengths model

$$n^{-1}E(\text{length of MST}) \rightarrow \zeta(3).$$

xxx java picture

**Example F:** (in Aldous-Steele survey). Uniform random tree on n vertices. Put i.i.d. positive weights on edges.

 $M_n :=$  weight of max-weight partial matching Then  $n^{-1}EM_n \to EM_\infty$  where

 $M_{\infty} = \frac{1}{2}$  (weight of edge at root)

in max-weight matching on limit infinite tree



Explicitly, for Exponential(1) edge-weights we get  $EM_{\infty} \approx 0.2396$  where the limit equals

 $\int_0^\infty se^{-s}ds \int_0^s c(e^{-y}-be^{-s}) \exp(-ce^{-y}-ce^{-(s-y)}) dy$ where  $c \approx 0.7146$  is the strictly positive solution of  $c^2 + e^{-c} = 1$  and  $b = \frac{c^2}{c^2+2c-1} \approx 0.5433$ .

# LWC and the cavity method

Recall model

 $G_n$ : complete graph on *n* vertices; edge-lengths random, independent Exponential(mean *n*) distribution.

Write  $L_n$  for length of TSP tour under this model. Mezard-Parisi (1980s) used replica/cavity methods to argue

 $n^{-1}EL_n \rightarrow c \approx 2.04.$ 

We have explicit program to make rigorous – but can't carry through two of the technical steps. Here are two relaxations of TSP for n-vertex network.

(M2F): minimum 2-factor. Minimize total length of a 2-factor, that is an edge-set in which each vertex has degree 2. That is, a union of cycles which spans all vertices.

(MA2F): minimum almost 2-factor. Minimize total length over edge-sets  $\mathcal{E}_n$  such that

$$n^{-1}|\{v: \text{ degree}(v) \neq 2\}| \rightarrow 0.$$

Proposition (Frieze 2004) In this model,

$$n^{-1}(EL_n - EL'_n) \to 0$$

where  $L_n$  is TSP length and  $L'_n$  is M2F length.

### Missing Proposition Want to know

$$n^{-1}(EL_n - EL_n'^{prime}) \to 0$$

where  $L_n$  is TSP length and  $L_n^{'prime}$  is MA2F length.

xxx java slide; hand write Zs

Central part of method – which I'll explain only superficially – is to do analysis of TSP on the (infinite) PWIT. each edge e of PWIT splits it into two subtrees. There are random variables  $Z^1(e), Z^2(e)$ , measurable functions of the subtrees, such that

 $e \in \mathsf{TSP}\text{-path iff length}(e) < Z^1(e) + Z^2(e)$ 

(another "missing proposition" in proving this) from which one can calculate mean length of TSP-path edges.

**Q:** How do we go back from the PWIT to the finite-n model?

### More math infrastructure

A measurable function  $f(Y_1, Y_2, ...)$  of some infinite collection of r.v.'s can be approximated arbitrary closely by continuous functions  $f_k(Y_1, ..., Y_k)$ of finitely many of the r.v.'s. So on the PWIT we can define an edge-set  $\mathcal{E}_{r}$  such that

(i) the edges of  $\mathcal{E}_r$  at a vertex v are determined by the restriction of the PWIT to the window of radius r around v; (ii)  $\delta_r := P($  some edge at  $v \in \mathcal{E}_r \Delta\{\} \to 0$  as  $r \to \infty$ .

Using LWC of the finite-n model to the PWIT, we can apply the same rule to a window of radius r around a vertex v, and define edgesets  $\mathcal{E}_{r,n}$  such that

 $\limsup_{n} P(\operatorname{degree}(v) \text{ in } \mathcal{E}_{r,n} \neq 2) \leq 2\delta_r$ xxx and similarly the edge-lengths xxx.

This constructs an almost-2-factor of  $G_n$  whose cost-per-vertex converges to the c given in the PWIT analysis.