

On the coalescence time of reversible random walks

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Abstract

Consider a system of coalescing random walks where each individual performs random walk over a finite graph \mathbf{G} , or (more generally) evolves according to some reversible Markov chain generator Q . Let C be the first time at which all walkers have coalesced into a single cluster. C is closely related to the consensus time of the voter model for this \mathbf{G} or Q .

We prove that the expected value of C is at most a constant multiple of the largest hitting time of an element in the state space. This solves a problem posed by Aldous and Fill and gives sharp bounds in many examples, including all vertex-transitive graphs. We also obtain results on the expected time until only $k \geq 2$ clusters remain. Our proof tools include a new exponential inequality for the meeting time of a reversible Markov chain and a deterministic trajectory, which we believe to be of independent interest.

Keywords: coalescing random walks, voter model, hitting time.

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1 Introduction

Consider a system of continuous-time random walks on a finite connected graph \mathbf{G} , with a walker starting from each vertex of \mathbf{G} . Let the walkers evolve independently, except that any two that occupy the same vertex of \mathbf{G} at a given time *coalesce* into one (this is made precise in Section 3.2).

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As time goes by, larger and larger coalesced clusters emerge, until at a certain random time C only one cluster remains. The question we address here is: how large can C be in terms of other parameters of \mathbf{G} ? This is a natural question which has implications for the so-called *voter model on \mathbf{G}* , discussed in Section 1.1 below.

It is instructive to consider what happens in the simple case of $\mathbf{G} = K_n$, the complete graph on n vertices. An explicit calculation [2, Chapter 14, Sec. 3.3] shows that:

$$\frac{C}{n} \approx \sum_{i=1}^{+\infty} \frac{Z_i}{i(i+1)}, \text{ with } \{Z_i\}_{i \geq 1} \text{ i.i.d. exponential random variables with mean 1.} \quad (1)$$

In particular, $\mathbb{E}[C] \sim n$ as $n \rightarrow +\infty$. What is remarkable about this is that any two of the walkers will take an expected time $\sim n/2$ to meet and coalesce; the fact that we are dealing with an *unbounded* number of particles only increases the expected time by a constant factor.

It is natural to ask what happens in more general graphs. This is closely related to the following problem, which was posed by Aldous and Fill in the mid-nineties.

Problem 1 (Open problem 13, Chapter 14 of [2]) *Prove that there exists a universal constant $K > 0$ such that the expected value of C satisfies*

$$\mathbb{E}[C] \leq K T_{\text{hit}}^{\mathbf{G}}$$

irrespective of initial conditions, where $T_{\text{hit}}^{\mathbf{G}}$ is the maximum expected hitting time of a vertex in \mathbf{G} .

To see how this relates to our previous discussion, consider a vertex-transitive graph \mathbf{G} . Proposition 5 in [2, Chapter 14] implies that the maximum expected meeting time of two walkers on \mathbf{G} , denoted by $T_{\text{meet}}^{\mathbf{G}}$, actually equals $T_{\text{hit}}^{\mathbf{G}}/2$. This implies that, if Problem 1 has a positive solution, all vertex-transitive graphs are like K_n in that $\mathbb{E}[C]$ is at most a universal constant factor away from $T_{\text{meet}}^{\mathbf{G}}$. A similar conclusion holds for the many other families of graphs where $T_{\text{meet}}^{\mathbf{G}} = \Theta(T_{\text{hit}}^{\mathbf{G}})$ (eg. all regular graphs with $T_{\text{hit}}^{\mathbf{G}} = O(n)$). For more general graphs it is still true that $T_{\text{meet}}^{\mathbf{G}} \leq T_{\text{hit}}^{\mathbf{G}}$, as proven in the aforementioned Proposition (see also [1]), and the Problem may be viewed as an strengthening of this fact¹.

¹There are graphs such as stars where $T_{\text{hit}}^{\mathbf{G}}$ is much larger than $T_{\text{meet}}^{\mathbf{G}}$ or $\mathbb{E}[C]$.

To the best of our knowledge, Problem 1 has remained open up to now. The best known bound of this sort has an extra $\ln |\mathbf{V}|$ factor; see Proposition 11 in [2, Chapter 14]. Our main goal in this paper is to give a solution of Problem 1 in the more general setting of reversible Markov chains.

Assume that Q is the generator of a reversible, irreducible, continuous-time Markov chain over a finite set \mathbf{V} . Given $v \in \mathbf{V}$, let H_v be the hitting time of v , ie. the first time at which a trajectory of Q hits v . We define the following parameter of the chain:

$$T_{\text{hit}}^Q \equiv \max_{v,w \in \mathbf{V}} \mathbb{E}_w [H_v] = \text{largest expected hitting time for } Q. \quad (2)$$

Define a system of coalescing random walks as in the case of graphs, with the difference that each walker now evolves over \mathbf{V} according to Q . The following Theorem solves Problem 1.

Theorem 1.1 *There exists a universal constant $K > 0$ such that, with Q as above, for any $n \in \mathbb{N} \setminus \{0\}$ and for any $x^{(n)} = (x(1), \dots, x(n)) \in \mathbf{V}^n$:*

$$\mathbb{E}_{x^{(n)}} [C] \leq K T_{\text{hit}}^Q.$$

Remark 1 *Here $x^{(n)}$ is an initial condition, with n arbitrary. In particular, there may be more or less than one walker at each site $v \in \mathbf{V}$ in the beginning of the process. Allowing for arbitrary initial conditions is convenient for our proofs, but does not really change the results.*

We also prove a stronger result. Let C_k denote the first time at which there are at most k clusters of coalesced walkers ($k \geq 1$). Notice that $C_1 = C$ with this definition.

Theorem 1.2 *There exists a universal constant $K_1 > 0$ such that, in the same setting of Theorem 1.1:*

$$\forall k \in \mathbb{N} \setminus \{0\}, \mathbb{E}_{x^{(n)}} [C_k] \leq K_1 \left(\frac{T_{\text{hit}}^Q}{k} + T_{\text{mix}}^Q \right),$$

where T_{mix}^Q is the mixing time of Q (see Section 2.2 for a definition).

The dependence on k in this Theorem is essentially best possible, as $\mathbb{E}_{x^{(n)}} [C_k] \sim T_{\text{hit}}^Q/k$ for large complete graphs. The case $k = 1$ gives back Theorem 1.1, as $T_{\text{mix}}^Q \leq c T_{\text{hit}}^Q$ for

some universal $c > 0$ [2, Chapter 4]. We will nevertheless prove Theorem 1.1 first and then show how its proof can be modified to obtain Theorem 1.2.

One justification for proving this second result is that it is helpful in approximating the *distribution* of C . We are in the process of writing a paper where we show that, if Q is transitive and $T_{\text{mix}}^Q \ll T_{\text{hit}}^Q$, then

$$\text{Law of } \frac{C}{T_{\text{hit}}^Q} \approx \sum_{i=1}^{+\infty} \frac{Z_i}{i(i+1)}, \text{ as for the complete graph (cf. (1)).}$$

In particular, $\mathbb{E}[C] \sim T_{\text{hit}}^Q$. This was previously known only for discrete tori \mathbb{Z}_L^d with $L \gg 1$ in $d \geq 2$ dimensions, due to Cox’s paper [4]². An important step in both our proof and Cox’s argument is that $\mathbb{E}[C_k] \ll T_{\text{hit}}^Q$ if $k \gg 1$. Cox proves this in [4, Section 4] via a simple renormalization argument which is very specific for discrete tori, whereas we use Theorem 1.2 for the same purpose.

1.1 Application to the voter model

We now sketch the connection between our results and the voter model [6, 2] on a graph \mathbf{G} (this could be generalized to an arbitrary generator Q , but we will not do this here). The state of the process at a given time t is a function:

$$\eta_t : V(\mathbf{G}) \rightarrow \mathcal{O}$$

where $V(\mathbf{G})$ is the vertex set of \mathbf{G} and \mathcal{O} is a fixed set of possible *opinions*. The evolution of the process is as follows. Each vertex $v \in V(\mathbf{G})$ “wakes up” at rate 1; when that happens at a time $t > 0$, v chooses one of its neighbors w uniformly at random and updates its value of $\eta_t(v)$ to w ’s opinion $\eta_{t-}(w)$; all other opinions stay the same.

A classical duality result (see eg. [6, Chapter 5] or [2, Chapter 14]) relates the state of the process at a given time to a system of coalescing random walks on \mathbf{G} moving backwards in time. In particular, the consensus time for the voter model – ie. the least time at which all vertices of \mathbf{G} have the same opinion – is dominated by the coalescence time C from the initial state with all vertices occupied. This implies the following Corollary of Theorem 1.1.

²Transitivity can be dropped at the cost of making stronger assumptions on Q and using a different normalization factor.

Corollary 1.1 *There exists a universal constant $K > 0$ such that, for any graph \mathbf{G} and any set \mathcal{O} , the expected value of the consensus time of the voter model defined in terms of \mathbf{G} and \mathcal{O} , started from an arbitrary initial state, is bounded by $K T_{\text{hit}}^{\mathbf{G}}$.*

Proposition 5 in [2, Chapter 14] shows that the Corollary is tight up to the value of K for vertex-transitive \mathbf{G} , at least when the initial conditions are iid uniform over $\{-1, +1\}$ (say); we omit the details.

1.2 Main proof ideas

Let us give an outline of the (elementary) proof of Theorem 1.1; the proof of Theorem 1.2 is quite similar. For clarity, we first present an oversimplified account, and then explain how one can avoid the oversimplifications.

We label the n walkers $(X_t(a))_{t \geq 0}$ with numbers $a = 1, \dots, n$. Instead of having walkers coalesce, we will assume that a walker $\#b$ will kill any walker $\#a$ with $a > b$ that happens to be in the same state as itself (this is made precise in Section 3.3). The number of walkers that are alive at time t in this process is precisely the number of clusters in the coalescing random walks process, and C is the first time at which only walker $\#1$ is still alive. This implies that:

$$\mathbb{P}(C > t) \leq \sum_{a=2}^n \mathbb{P}(\text{walker } \# a \text{ alive at time } t).$$

We now make the following oversimplification:

Oversimplification #1: walker $\#a$ dies at the first time when $X_t(a) = X_t(b)$ for some $b < a$.

The reason why this is an oversimplification is that a walker $\#b$ may have died before meeting walker $\#a$. For the moment, we ignore this and write:

$$\mathbb{P}(\text{walker } \# a \text{ alive at time } t) \leq \mathbb{P}\left(\bigcap_{b=1}^{a-1} \{\forall 0 \leq s \leq t, X_s(a) \neq X_s(b)\}\right).$$

In order to simplify the RHS, we notice that the trajectories $(X_t(u))_{t \geq 0}$ of walkers $\#u$, $1 \leq u \leq a$, are independent realizations of Q . Conditioning on $X_s(a) = h_s$, $s \geq 0$, makes

the events in the RHS independent, and we deduce:

$$\mathbb{P}(\text{walker \# } a \text{ alive at time } t \mid (X_s(a))_{s \geq 0} = (h_s)_{s \geq 0}) \leq \prod_{b=1}^{a-1} \mathbb{P}(\forall 0 \leq s \leq t, X_s(b) \neq h_s).$$

We now make another oversimplification.

Oversimplification #2: $(X_t(b))_{t \geq 0}$ is started from the stationary distribution for all b .

This allows us to use the following Lemma, which we believe to be new (and of independent interest).

Lemma 1.1 (Meeting time Lemma; proven in Section 5.2) *Let $(X_t)_{t \geq 0}$ be a realization of Q starting from the stationary distribution π . Then there exist $v \in \mathbf{V}$ and a quasistationary distribution \mathbf{q}_v for $\mathbf{V} \setminus \{v\}$ such that for any deterministic path $h \in \mathbb{D}([0, +\infty), \mathbf{V})$, we have:*

$$\forall t \geq 0, \mathbb{P}_\pi(\forall 0 \leq s \leq t, X_s \neq h_s) \leq \mathbb{P}_{\mathbf{q}_v}(H_v > t) = \exp\left(-\frac{t}{\mathbb{E}_{\mathbf{q}_v}[H_v]}\right).$$

Remark 2 *The proof of Lemma 1.1 shows that we may take $v \in h([0, +\infty))$. This is a well-known result if $h \equiv v$ [2, Chapter 3, Section 6.5]. An application of this Lemma to so-called cat-and-mouse games is sketched in the final section.*

Notice that $\mathbb{E}_{\mathbf{q}_v}[H_v] \leq T_{\text{hit}}^Q$, so:

$$\mathbb{P}(\text{walker \# } a \text{ alive at time } t \mid (X_s(i))_{s \geq 0} = (h_s)_{s \geq 0}) \leq e^{-\frac{(a-1)t}{T_{\text{hit}}^Q}}.$$

This shows that:

$$\mathbb{P}(C > t) \leq \sum_{a=2}^n e^{-\frac{(a-1)t}{T_{\text{hit}}^Q}}.$$

If one takes $t = (\ln 2 + c)T_{\text{hit}}^Q$, the RHS becomes:

$$\mathbb{P}\left(C > (\ln 2 + c)T_{\text{hit}}^Q\right) \leq \sum_{a=2}^n 2^{-a+1} e^{-(a-1)c} \leq e^{-c},$$

and this gives $\mathbb{E}[C] \leq (\ln 2 + 2)T_{\text{hit}}^Q$.

Of course, this is *not* a proof of Theorem 1.1 because of the oversimplifications. Our way out of this is to introduce a process where at any given time there is a list of *allowed killings*. At any time t there will be a set \mathcal{A}_t , so that walker $\#b$ may kill walker $\#a$ at time t only if $b < a$ and $(b, a) \in \mathcal{A}_t$ (cf. Section 3.4). The salient characteristics of this process are:

1. For any choice of $\mathcal{A} = (\mathcal{A}_t)_{t \geq 0}$, the set of alive walkers in the process defined via \mathcal{A} dominates the corresponding set in the process without \mathcal{A} (see Proposition 3.2).
2. A judicious choice of \mathcal{A} will ensure that for each a , there will be a large enough time interval where a large number of walkers will be available to kill walker $\#a$. Moreover, many of these will be stationary.

Item 1 allows us to consider the process with a list of allowed killings instead of the original process in order to obtain upper bounds. Item 2 will mean that we may apply the Meeting Time Lemma to at least some of the walkers with indices $b < a$, in some time intervals. These two ingredients will allow us to “fix” the oversimplified proof just presented.

1.3 Organization

The remainder of the paper is organized as follows. Section 2 introduces our notation and recalls some basic concepts. Section 3 defines the main processes we consider in the paper. Section 4 presents the proofs of the two Theorems, and Section 5 presents the proof of Lemma 1.1. Some final comments are presented in the last Section.

2 Preliminaries

In what follows we recall some basic material while also fixing notation.

2.1 Basic notation

\mathbb{N} is the set of non-negative integers. Given $n \in \mathbb{N} \setminus \{0\}$, we set $[n] \equiv \{1, 2, \dots, n\}$.

We will often speak of *universal constants*. These are numbers that are independent of any other object or parameter under consideration, be it a Markov chain, the initial state of a process under consideration or anything else.

The cardinality of a finite set S is denoted by $|S|$, and 2^S represents the power set of S (ie. the set whose elements are the subsets of S). The set of all probability measures over S will be denoted by $M_1(S)$. \mathbb{R}^S denotes the space of all functions $f : S \rightarrow \mathbb{R}$, or equivalently of all (column) vectors with entries indexed by S . Linear operators acting on \mathbb{R}^S correspond to matrices with rows and columns indexed by the elements of S . If A is some matrix of this sort, $\|A\|_{\text{op}}$ is the operator norm of A . If A is symmetric, we let $\lambda_{\min}(A)$, $\lambda_{\max}(A)$ denote the minimum and maximum eigenvalues of A (respectively).

Given a finite set $F \neq \emptyset$, a function $\omega : [0, +\infty) \rightarrow F$ is said to be càdlàg if there exist $t_0 = 0 < t_1 < t_2 < \dots < t_n < \dots \nearrow +\infty$ with ω constant over each interval $[t_{i-1}, t_i)$. $\mathbb{D}([0, +\infty), F)$ is the set of all such càdlàg functions, with the σ -field generated by the projections “ $\omega \mapsto \omega(t)$ ” ($t \geq 0$).

2.2 Markov chain basics

Let \mathbf{V} be a finite, non-empty set. A matrix Q (with rows and columns labelled by \mathbf{V}) which acts on $\mathbb{R}^{\mathbf{V}}$ in the following way:

$$Q : f(\cdot) \in \mathbb{R}^{\mathbf{V}} \mapsto \sum_{x \in \mathbf{V}, x \neq \cdot} q(\cdot, x)(f(\cdot) - f(x)), \text{ with } q(\cdot, \cdot) \geq 0$$

defines a unique *continuous-time Markov chain* on \mathbf{V} . More precisely, there exists a unique family of measures $\{\mathbb{P}_x\}_{x \in \mathbf{V}}$ over $\mathbb{D}([0, +\infty), \mathbf{V})$ (with the σ -field generated by finite-dimensional projections) such that, letting

$$X_t : \omega \in \mathbb{D}([0, +\infty), \mathbf{V}) \mapsto X_t(\omega) \equiv \omega(t) \in \mathbf{V}$$

and $\mathcal{F}_t \equiv \sigma(X_s : s \leq t)$, we have $\mathbb{P}_x(X_0 = x) = 1$ and

$$\mathbb{P}_x(X_{t+s} = y \mid \mathcal{F}_s) = \text{the entry of } e^{-tQ} \text{ labelled by } (X_t, y) \text{ } (t, s \geq 0, y \in \mathbf{V}). \quad (3)$$

Q is said to be the generator of the Markov chain and the numbers $q(x, y)$ ($x, y \in \mathbf{V}$, $x \neq y$) are the transition rates. We let $\mathbb{E}_x[\cdot]$ denote expectation with respect to \mathbb{P}_x .

We also define

$$\mathbb{P}_\mu = \sum_{x \in \mathbf{V}} \mu(x) \mathbb{P}_x, \quad (\mu \in M_1(\mathbf{V}))$$

which we interpret in the customary way, as describing the law of the chain given by Q from a random initial state with law μ . $\mathbb{E}_\mu[\cdot]$ is the corresponding expectation symbol.

We will always assume that Q is *irreducible*, meaning that for all $A \subset \mathbf{V}$ with $A, \mathbf{V} \setminus A \neq \emptyset$ there exist $a \in A$ and $b \in \mathbf{V} \setminus A$ with $q(a, b) \neq 0$. In this case there exists a unique probability measure $\pi \in M_1(\mathbf{V})$ which is *stationary* in the sense that $\mathbb{P}_\pi(X_t = \cdot) = \pi(\cdot)$ for all $t \geq 0$. Moreover, we have that:

$$\forall x, y \in \mathbf{V}, \quad \lim_{t \rightarrow +\infty} \mathbb{P}_x(X_t = y) = \pi(y).$$

The *mixing time* of Q measures the speed of this convergence:

$$T_{\text{mix}}^Q \equiv \inf \left\{ t \geq 0 : \forall x \in \mathbf{V}, \max_{S \subset \mathbf{V}} |\mathbb{P}_x(X_t \in S) - \pi(S)| \leq 1/4 \right\}.$$

Finally, we will also assume that Q is *reversible* with respect to π , which means that $\pi(x)q(x, y) = \pi(y)q(y, x)$ for all distinct $x, y \in \mathbf{V}$. This is the same as requiring that the matrix $\Pi^{1/2} Q \Pi^{-1/2}$ is symmetric, where Π is diagonal and has the values $\pi(v)$, $v \in \mathbf{V}$ on the diagonal.

3 Processes with multiple random walks

We define here the main processes that we will be concerned with, all of which involve n random walkers for some integer $n \in \mathbb{N} \setminus \{0, 1\}$. We will assume that Q and $\{\mathbb{P}_x\}_{x \in \mathbf{V}}$ are as defined in Section 2.2.

3.1 Independent random walks

We first define a processes made out of n independent realizations of the Markov chain with generator Q . More specifically, given

$$x^{(n)} = (x(1), x(2), \dots, x(n)) \in \mathbf{V}^n,$$

we let $\mathbb{P}_{x^{(n)}}$ denote the distribution on $\mathbb{D}([0, +\infty), \mathbf{V}^n)$ corresponding to n independent trajectories of Q ,

$$(X_t^{(n)})_{t \geq 0} \equiv (X_t(a) : a \in [n])_{t \geq 0}, \quad (4)$$

with each $(X_t(a))_{t \geq 0}$ started from $x(a)$. That is, the joint law of $\{(X_t(a))_{t \geq 0}\}_{a \in [n]}$ is the product measure:

$$\mathbb{P}_{x(1)} \times \mathbb{P}_{x(2)} \times \cdots \times \mathbb{P}_{x(n)}.$$

Notice that our notation $\mathbb{P}_{x^{(n)}}$ does not refer explicitly to the fact that this is a process on \mathbf{V}^n , as opposed to the process over \mathbf{V} defined in the previous subsection. This distinction should be clear from context and from the fact that we write all $x^{(n)} \in \mathbf{V}^n$ with a “ (n) ” superscript. The independent random walks process is also a Markov chain: for $x^{(n)} = (x(1), \dots, x(n))$ and $y^{(n)} = (y(1), \dots, y(n))$ distinct, the transition rate from $x^{(n)}$ to $y^{(n)}$ is:

$$q^{(n)}(x^{(n)}, y^{(n)}) \equiv \begin{cases} q(x(i), y(i)) & \text{if } x(i) \neq y(i) \text{ and } x(j) = y(j) \text{ for all } j \in [n] \setminus \{i\}; \\ 0 & \text{otherwise.} \end{cases} \quad (5)$$

3.2 Coalescing random walks

For our purposes, it is convenient to define this process, denoted by

$$(\text{Co}_t^{(n)})_{t \geq 0} \equiv (\text{Co}_t(a) : a \in [n])_{t \geq 0}$$

as a deterministic function of the independent random walks process. The idea is that, once a walker meets another walker with smaller index, it starts following the trajectory of the latter. That is, consider a realization of $\mathbb{P}_{x^{(n)}}$ as in (4). First define:

$$\text{Co}_t(1) \equiv X_t(1) \quad (t \geq 0).$$

Given $a \in [n] \setminus \{1\}$, assume inductively that $(\text{Co}_t(b))_{t \geq 0}$ has been defined for $1 \leq b < a$. Since Q is irreducible, there a.s. is a *first* time τ_a at which $X_t(a) = \text{Co}_t(b)$ for some $b \in [a-1]$. More precisely, define:

$$\tau_a \equiv \inf\{t \geq 0 : \exists 1 \leq b < a, X_t(a) = \text{Co}_t(b)\}$$

and

$$B_a \equiv \min\{b \in [a-1] : X_{\tau_a}(a) = X_{\tau_a}(b)\}$$

and then set:

$$\text{Co}_t(a) \equiv \begin{cases} X_t(a), & 0 \leq t < \tau_a; \\ \text{Co}_t(B_a), & t \geq \tau_a; \end{cases} \quad \text{for each } t \geq 0.$$

One can show that the law of $(\text{Co}_t^{(n)})_{t \geq 0}$ is invariant under permutations of the $x(i)$. We also *define* the set:

$$S_t \equiv \{v \in \mathbf{V} : \exists a \in [n], \text{Co}_t(a) = v\} \quad (6)$$

as the set of occupied sites in this process. Our definition of $(S_t)_{t \geq 0}$ coincides with the more traditional coalescing random walks process defined in eg. [4]. We also set:

$$C_k \equiv \inf\{t \geq 0 : |S_t| \leq k\} \quad (k \in \mathbb{N} \setminus \{0\})$$

and $C \equiv C_k$.

Remark 3 *We note that this process makes sense even if $x^{(n)}$ contains repeats, ie. if there exist $i \neq j$ with $x(i) = x(j)$.*

3.3 Random walks with killings

Let $\partial \notin V$ be a “coffin state”. We define a new process

$$(Y_t^{(n)})_{t \geq 0} \equiv (Y_t(a) : a \in [n])_{t \geq 0}.$$

The new idea is that a walker with index a will be killed by a walker of index $b < a$ occupying the same site. More precisely, we first define:

$$Y_t(1) \equiv X_t(1) \quad (t \geq 0).$$

Given $a \in [n] \setminus \{1\}$, assume inductively that $(Y_t(b))_{t \geq 0}$ has been defined for $1 \leq b < a$. Define:

$$\tau_a \equiv \inf\{t \geq 0 : \exists 1 \leq b < a, X_t(a) = Y_t(b)\}$$

and set:

$$Y_t(a) \equiv \begin{cases} X_t(a), & 0 \leq t < \tau_a; \\ \partial, & t \geq \tau_a; \end{cases} \quad \text{for each } t \geq 0.$$

Although our new definition of τ_a different from the previous one, it is easy to show that the two definitions coincide, and that in fact:

Proposition 3.1 (Proof omitted) *Let S_t be as (6). Then for all $t \geq 0$,*

$$S_t = \{v \in \mathbf{V} : \exists a \in [n], Y_t(v) = a\}$$

and

$$|S_t| = |\{a \in [n] : Y_t(a) \neq \partial\}|.$$

Therefore, for all $k \in \mathbb{N} \setminus \{0\}$,

$$\mathbb{P}_{x^{(n)}}(C_k > t) = \mathbb{P}_{x^{(n)}}(|S_t| \geq k + 1) = \mathbb{P}_{x^{(n)}}(|\{a \in [n] : Y_t(a) \neq \partial\}| \geq k + 1).$$

Remark 4 *As in Remark 3, we may allow $x^{(n)}$ where $x(i) = x(j)$ for some pair $i \neq j$. Notice, however, that $Y_0^{(n)} \neq x^{(n)}$ in this case.*

3.4 Random walks with a list of allowed killings

Now assume that we have a deterministic càdlàg trajectory:

$$\mathcal{A} : t \geq 0 \mapsto 2^{[n]^2}.$$

We define yet another process:

$$((Y_t^{\mathcal{A}})^{(n)})_{t \geq 0} \equiv (Y_t^{\mathcal{A}}(a) : a \in [n])_{t \geq 0}$$

where a walker with index a may be killed by a walker with index b only if they occupy the same site at some time t and $(b, a) \in \mathcal{A}_t$. Intuitively, this means that b is allowed to kill a only at times t with $(b, a) \in \mathcal{A}_t$.

For a formal definition, we first set:

$$Y_t^{\mathcal{A}}(1) \equiv X_t(1) \quad (t \geq 0).$$

Given $a \in [n] \setminus \{1\}$, assume inductively that $(Y_t^{\mathcal{A}}(b))_{t \geq 0}$ has been defined for $1 \leq b < a$.

Define:

$$\tau_a^{\mathcal{A}} \equiv \inf\{t \geq 0 : \exists 1 \leq b < a, (b, a) \in \mathcal{A}_t \text{ and } X_t(a) = Y_t^{\mathcal{A}}(b)\}$$

and set:

$$Y_t^{\mathcal{A}}(a) \equiv \begin{cases} X_t(a), & 0 \leq t < \tau_a^{\mathcal{A}}; \\ \partial, & t \geq \tau_a^{\mathcal{A}}; \end{cases} \quad \text{for each } t \geq 0.$$

The following Proposition shows that the process with a list of allowed killings can be used to upper bound $\mathbb{E}_{x^{(n)}}[C_k]$.

Proposition 3.2 *Define:*

$$S_t^{\mathcal{A}} \equiv \{Y_t^{\mathcal{A}}(a) : a \in [n]\}.$$

For any choice of \mathcal{A} as above and of initial state $x^{(n)}$, one can couple $(S_t)_{t \geq 0}$ and $(S_t^{\mathcal{A}})_{t \geq 0}$ such that (almost surely) $S_t \subset S_t^{\mathcal{A}}$ for all $t \geq 0$. In particular, for all $k \in \mathbb{N} \setminus \{0\}$,

$$\mathbb{P}_{x^{(n)}}(C_k > t) = \mathbb{P}_{x^{(n)}}(|S_t| \geq k + 1) \leq \mathbb{P}_{x^{(n)}}(|S_t^{\mathcal{A}}| \geq k + 1).$$

We omit the proof of this rather intuitive Proposition. The key idea here is this: suppose we do *not* kill a walker a at a given time t_0 . The only way this could make S_t “smaller” is if $X_t(a)$ were to meet a walker $X_t(c)$ with $c > a$ at some later time $t \geq t_0$. But if this happens, we may pretend that $X_s(a)$ follows the trajectory of $X_s(c)$ for $s \geq t$; this follows from the Markov property coupled with the fact that $X_t(a) = X_t(c)$. This shows that in fact S_t does *not* become smaller.

Remark 5 *Similarly to Remark 4, we note that we may allow $x^{(n)}$ with $x(i) = x(j)$ for some pair $i \neq j$, but then $(Y_0^{\mathcal{A}})^{(n)} \neq x^{(n)}$.*

4 Proofs of the main Theorems

We prove Theorems 1.1 and 1.2 in this Section. For simplicity, we first focus on the proof of Theorem 1.1, and then show how it can be modified to prove the second Theorem. We will take the notation and definitions in Sections 2.2 and 3 for granted.

4.1 Preliminaries for the proof of Theorem 1.1

We first note that Theorem 1.1 follows from a seemingly different statement.

Proposition 4.1 *Let $c, \gamma > 0$ be given universal constants. Suppose we can show that there exists some choice of $\mathcal{A} = (\mathcal{A}_t)_{t \geq 0}$ as in Section 3.4 and of $0 \leq t_0 \leq c(\mathbb{T}_{\text{mix}}^Q + \mathbb{T}_{\text{hit}}^Q)$ with*

$$\forall n \in \mathbb{N} \setminus \{0\}, \forall x^{(n)} \in \mathbf{V}^n, \mathbb{P}_{x^{(n)}}(|S_{t_0}^{\mathcal{A}}| \geq 2) \leq 1 - \gamma.$$

Then

$$\forall n \in \mathbb{N} \setminus \{0\}, \forall x^{(n)} \in \mathbf{V}^n, \mathbb{E}_{x^{(n)}}[C] \leq K \mathbb{T}_{\text{hit}}^Q$$

where $K > 0$ is universal.

Proof: Given $s \geq 0$, denote:

$$E(s) \equiv \{C > s\} = \{|S_s| \geq 2\} = \bigcup_{a \in [n] \setminus \{1\}} \{\tau_a > s\}. \quad (7)$$

Combining the assumption of the Proposition with Proposition 3.2 gives:

$$\forall n \in \mathbb{N} \setminus \{1\}, \forall x^{(n)} \in \mathbf{V}^n, \mathbb{P}_{x^{(n)}}(E(t_0)) \leq \mathbb{P}_{x^{(n)}}(|S_{t_0}^{\mathcal{A}}| \geq 2) \leq 1 - \gamma \quad (8)$$

We now consider $E(\ell t_0)$ where $\ell > 1$ is an integer. Let $(\Theta_s)_{s \geq 0}$ denote the time-shift operators for the independent random walks process and let $(\mathcal{F}_s^{(n)})_{s \geq 0}$ denote the filtration generated by this process.

$$\begin{aligned} \mathbb{P}_{x^{(n)}}(E(k t_0)) &\leq \mathbb{P}_{x^{(n)}}(E((\ell - 1)t_0) \cap (\cup_{a=2}^n \{\tau_a \circ \Theta_{(\ell-1)t_0} > t_0\})) \\ &= \mathbb{P}_{x^{(n)}}(E((\ell - 1)t_0) \cap \Theta_{(\ell-1)t_0}^{-1}(E(t_0))) \\ (E((\ell - 1)t_0) \in \mathcal{F}_{(\ell-1)t_0}^{(n)}) &\leq \mathbb{E}_{x^{(n)}} \left[\mathbb{I}_{E((\ell-1)t_0)} \mathbb{P}_{x^{(n)}} \left(\Theta_{(\ell-1)t_0}^{-1}(E(t_0)) \mid \mathcal{F}_{(\ell-1)t_0}^{(n)} \right) \right] \\ (\text{Markov property}) &= \mathbb{E}_{x^{(n)}} \left[\mathbb{I}_{E((\ell-1)t_0)} \mathbb{P}_{X_{(\ell-1)t_0}^{(n)}}(E(t_0)) \right] \\ (\text{inequality (8)}) &\leq \mathbb{E}_{x^{(n)}} \left[\mathbb{I}_{E((\ell-1)t_0)} (1 - \gamma) \right] \\ &= \mathbb{P}_{x^{(n)}}(E((\ell - 1)t_0)) (1 - \gamma) \\ (\text{induction on } k) &\leq (1 - \gamma)^\ell. \end{aligned}$$

Recalling the definition of $E(s)$, we deduce that:

$$\forall n \in \mathbb{N} \setminus \{1\}, \forall x^{(n)} \in \mathbf{V}^n, \mathbb{E}_{x^{(n)}} [C] \leq \sum_{\ell \in \mathbb{N}} (1 - \gamma)^\ell t_0 = \frac{t_0}{\gamma} \leq \frac{c(\mathbb{T}_{\text{mix}}^Q + \mathbb{T}_{\text{hit}}^Q)}{\gamma}.$$

The Proposition follows from this because $\mathbb{T}_{\text{mix}}^Q \leq c_0 \mathbb{T}_{\text{hit}}^Q$ for some universal $c_0 > 0$ [2, Chapter 3] and both c and γ are universal. \square

4.2 Construction of \mathcal{A}

Notational convention 1 From now on, we fix some $x^{(n)}$ and write \mathbb{P} instead of $\mathbb{P}_{x^{(n)}}$.

We will now design a specific trajectory $\mathcal{A} = (\mathcal{A}_t)_{t \geq 0}$ which will allow for a simple analysis of $S_t^{\mathcal{A}}$. Let $m \in \mathbb{N}$ be the smallest non-negative number with $n \leq \sum_{i=0}^m 2^i$. Define sets

$$\begin{aligned} A_0 &= \{1\}; \\ A_j &\equiv \left[\sum_{i=0}^j 2^i \right] \setminus \left[\sum_{i=0}^{j-1} 2^i \right] \quad (1 \leq j \leq m-1); \\ \text{and } A_m &\equiv [n] \setminus \left[\sum_{i=0}^{m-1} 2^i \right]. \end{aligned}$$

We will consider different *epochs*, numbered backwards in time. It is convenient to have the following notation.

$$\begin{aligned} t_m &\equiv 2\mathbb{T}_{\text{mix}}^Q; \\ t_j &= t_{j+1} + (\ln 5) 2^{4-j} \mathbb{T}_{\text{hit}}^Q, \quad j = m-1, m-2, \dots, 0. \end{aligned} \tag{9}$$

1. *Epoch # ∞* is the time interval $[0, t_m)$. We set $\mathcal{A}_t \equiv \emptyset$ for all t in this interval, ie. no killings are allowed up to time $2\mathbb{T}_{\text{mix}}^Q$.
2. *Epochs # m through #1* correspond to time intervals $I_j = [t_j, t_{j-1})$ as j decreases from m to 1. For each such j we set:

$$\mathcal{A}_t \equiv A_{j-1} \times \cup_{p=j}^m A_p, \quad t \in I_j.$$

That is, the only killings allowed are between walkers with labels in A_{j-1} and A_p with $p \geq j$.

3. *Epoch #0* corresponds to the time interval,

$$I_0 \equiv [t_0, +\infty),$$

(the remaining time), where we set $\mathcal{A}_t \equiv [n]^2$.

We note for later convenience that:

$$t_0 \leq 2T_{\text{mix}}^Q + (\ln 5) \sum_{j \geq 0} 2^{4-j} T_{\text{hit}}^Q \leq c(T_{\text{mix}}^Q + T_{\text{hit}}^Q) \quad (10)$$

with $c > 0$ universal, since $\sum_j 2^{-j} < +\infty$. We will use this in our application of Proposition 4.1.

4.3 Abundance of good walkers

We have the following simple proposition about the epoch $\#\infty$. Intuitively, it says that, at the end of this epoch, a positive proportion of the random walkers are “good”, in that they have converged to stationarity.

Proposition 4.2 *One can construct a (random) subset $R \subset [n]$ such that:*

1. R is $\mathcal{H}_{t_m}^{(n)}$ -measurable, where $\mathcal{H}_{t_m}^{(n)}$ is the sigma-field generated by $(X_s^{(n)})_{s \leq t_m}$ and by some additional independent random variable U .
2. Each $r \in [n]$ belongs to R with probability $1/4$, independently of all other $r' \in [n]$.
3. Conditionally on R and on $(X_{t_m}(i))_{i \in [n] \setminus R}$, the vector $(X_{t_m}(r))_{r \in R}$ has iid coordinates, each with distribution π .

Proof: Consider a single $a \in [n]$. Since Q is reversible, Lemma 7 in [2, Chapter 4] shows that:

$$\forall a \in [n], \forall v \in \mathbf{V} : \mathbb{P} \left(X_{2T_{\text{mix}}^Q}(a) = v \right) \geq \frac{\pi(v)}{4};$$

in other words, for each a there exists some $\nu_a \in M_1(\mathbf{V})$ such that:

$$\mathbb{P}\left(X_{2T_{\text{mix}}^Q}(a) = \cdot\right) = \frac{1}{4}\pi(\cdot) + \frac{3}{4}\nu_a(\cdot).$$

Since the random variables $(X_{t_m}(a))_{a \in [n]}$ are independent, we may assume that they sampled as follows:

1. Let $(I(a))_{a \in [n] \setminus A_m}$ be iid with $\mathbb{P}(I(a) = 1) = 1 - \mathbb{P}(I(a) = 0) = 1/4$.
2. For each a with $I(a) = 1$, let $X_{t_m}(a)$ be a sample from π , independent of everything else.
3. For each b with $I(b) = 0$, let $X_{t_m}(b)$ be a sample from ν_b , independent of everything else.

One may check that $R \equiv \{a \in [n] \setminus A_m : I(a) = 1\}$ has the desired properties. \square

The next proposition means that, with positive probability, there is a constant proportion of good walkers within each A_i with $i \leq m - 1$.

Proposition 4.3 *Let \mathcal{G} be the event:*

$$\mathcal{G} \equiv \bigcap_{i=0}^{m-1} \{|R \cap A_i| \geq 2^{i-3}\}.$$

Then $\mathbb{P}(\mathcal{G}) \geq \alpha > 0$, where

$$\alpha \equiv \prod_{i=0}^{+\infty} (1 - e^{-2^{i-7}}) > 0$$

is universal.

Proof: Let $\text{Bin}(m, x)$ denote a binomial random variable with parameters m and x , so that:

$$\mathbb{P}(\text{Bin}(m, x) = k) = \binom{m}{k} x^k (1-x)^{m-k} \quad (k \in [m] \cup \{0\}).$$

The random variables $N_i = |R \cap A_i|$, $0 \leq i \leq m - 1$ are independent, and each N_i has the law of $\text{Bin}(2^i, 1/4)$. Chernoff bounds [3, Appendix A.1] imply:

$$\mathbb{P}(|R \cap A_i| < 2^{i-3}) = \mathbb{P}(\text{Bin}(2^i, 1/4) < \mathbb{E}[\text{Bin}(2^i, 1/4)] - 2^{i-3}) \leq e^{-\frac{(2^{i-3})^2}{2 \cdot 2^i}} = e^{-2^{i-7}}.$$

We deduce:

$$\mathbb{P}(\mathcal{G}) \geq \prod_{i=0}^{m-1} \mathbb{P}(\text{Bin}(2^i, 1/4) < 2^{i-3}) \geq \alpha.$$

The positivity of α follows from $0 < e^{-2^{i-7}} < 1$ for all i and $\sum_i e^{-2^{i-7}} < +\infty$. \square

4.4 The probability of being alive

Let $E(a, t)$ denote the event:

$$E(a, t) \equiv \{Y_t^A(a) \neq \partial\} = \{\tau_a^A > t\}.$$

Notice that:

$$\{|S_t^A| \geq 2\} \subset \bigcup_{a=2}^n E(a, t). \quad (11)$$

We will now compute estimate the conditional probability of $E(a, t)$ given \mathcal{G} .

Proposition 4.4 *Let $a \in A_i$ for some $1 \leq i \leq m$. Then for all $1 \leq j < i$:*

$$\mathbb{P}(E(a, t_j) \mid \mathcal{G}) \leq 5^{j-i}.$$

Proof: We will prove a stronger statement: that for almost all $R_0 \subset [n]$ and $(h_t)_{t \in [t_i, t_j]}$:

$$\mathbb{P}\left(E(a, t_j) \mid R = R_0, (X_t(a))_{t \in [t_i, t_j]} = (h_t)_{t \in [t_i, t_j]}\right) \leq 5^{-\sum_{r=j}^{i-1} \frac{|R_0 \cap A_r|}{2^{r-4}}}. \quad (12)$$

This implies the proposition because the occurrence of \mathcal{G} implies $|R \cap A_r| \geq 2^{r-4}$ for all $1 \leq r < m - 1$.

To prove (12) we first observe that the event $E(a, t_j)$ satisfies:

Claim 4.1 *Suppose $b \in A_r$ with $j \leq r < i$. Then:*

$$E(a, t_j) \subset \{\forall t \in [t_{r+1}, t_r), X_t(a) \neq X_t(b)\}.$$

Proof: [of the Claim] If the event in the RHS does not hold, there exists a $t \in [t_r, t_{r-1})$ with $X_t(a) = X_t(b)$. We now argue that $\tau_a^A \leq t$ in this case. Indeed, this follows from the definition of τ_a^A and the following observations:

1. $X_t(a) = Y_t^A(b)$: this follows from $X_t(b) = Y_t^A(b)$, which is a consequence of the fact that $(b, c) \notin \mathcal{A}_s$ for any $c > b$ and $s \leq t_r$ (ie. b cannot be killed before time t_r).
2. $(b, a) \in \mathcal{A}_t$: this follows from $t \in [t_{r+1}, t_r) = I_{r+1}$.

□

The Claim implies:

$$\begin{aligned}
& \mathbb{P} \left(E(a, t_j) \mid R = R_0, (X_t(a))_{t \in [t_i, t_{i-1}]} = (h_t)_{t \in [t_i, t_{i-1}]} \right) \\
& \leq \mathbb{P} \left(\bigcap_{r=j}^{i-1} \bigcap_{b \in A_r} \{ \forall t \in [t_r, t_{r-1}), X_t(a) \neq X_t(b) \} \mid R = R_0, (X_t(a))_{t \in [t_i, t_j]} = (h_t)_{t \in [t_i, t_j]} \right) \\
& \leq \mathbb{P} \left(\bigcap_{r=j}^{i-1} \bigcap_{b \in A_r \cap R_0} \{ \forall t \in [t_{r+1}, t_r), X_t(a) \neq X_t(b) \} \mid R = R_0, (X_t(a))_{t \in [t_i, t_j]} = (h_t)_{t \in [t_i, t_j]} \right).
\end{aligned}$$

Now observe that we are conditioning on $R = R_0$ and on the trajectory of $(X_s(a))_{s \in [t_i, t_{i-1}]}$. Since $a \notin A_{i-1}$, Proposition 4.2 implies that:

Under the conditioning, $(X_{t_m}(b) : b \in R_0 \cap (\cup_{r=j}^{i-1} A_r))$ are iid with common law π .

Since R is $\mathcal{H}_{t_m}^{(n)}$ -measurable, the Markov property for the independent random walks process implies that

Under the conditioning, $(X_{t+t_m}(b) : b \in R_0 \cap (\cup_{r=j}^{i-1} A_r))_{t \geq 0}$ are iid realizations of \mathbb{P}_π .

We deduce:

$$\begin{aligned}
& \mathbb{P} \left(E(a, t_j) \mid R = R_0, (X_t(a))_{t \in [t_i, t_j]} = (h_t)_{t \in [t_i, t_j]} \right) \\
& = \prod_{r=j}^{i-1} \prod_{b \in R_0 \cap A_r} \mathbb{P}_\pi (\forall t \in [t_{r+1}, t_r), X_t \neq h_t) \\
& (\mathbb{P}_\pi \text{ is stationary}) = \prod_{r=j}^{i-1} \mathbb{P}_\pi (\forall t \in [0, t_r - t_{r+1}), X_t(b) \neq h_{t+t_r})^{|R_0 \cap A_r|}.
\end{aligned}$$

We apply the Meeting Time Lemma (Lemma 1.1 above) to each term in the product and deduce that, for some choice of (v, \mathbf{q}_v) as in the Lemma,

$$\begin{aligned} \mathbb{P}(E(a, t_{i-1}) \mid R = R_0, (X_t(a))_{t \in [t_i, t_{i-1}]} = (h_t)_{t \in [t_i, t_{i-1}]}) \\ \leq \prod_{r=j}^{i-1} \left\{ \exp \left(- \frac{(t_r - t_{r+1}) |R_0 \cap A_r|}{\mathbb{E}_{\mathbf{q}_v} [H_v]} \right) \right\}. \end{aligned}$$

The proof of (12) finishes once we realize that $t_r - t_{r+1} = 2^{4-r} (\ln 5) T_{\text{hit}}^Q$ and $T_{\text{hit}}^Q \geq \mathbb{E}_{\mathbf{q}_v} [H_v]$.
□

4.5 End of proof of Theorem 1.1

We now complete the proof of Theorem 1.1. By Proposition 4.1, it suffices to show that:

$$\mathbb{P}_{x^{(n)}} (|S_{t_0}^A| \geq 2) \leq 1 - \gamma$$

for some universal $\gamma > 0$, with t_0 as in (10). To see this, we will use (11) and recall our convention of omitting $x^{(n)}$ from the notation (cf. Notational convention 1).

$$\begin{aligned} \mathbb{P}(|S_{t_0}^A| \geq 2) &= \mathbb{P} \left(\bigcup_{a=2}^n E(a, t_0) \right) \\ (\mathcal{G} \text{ as in Prop. 4.3}) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P} \left(\mathcal{G} \cap \bigcup_{a=2}^n E(a, t_0) \right) \\ (\text{union bound}) &\leq 1 - \mathbb{P}(\mathcal{G}) + \sum_{a=2}^n \mathbb{P}(\mathcal{G} \cap E(a, t_0)) \\ ([n] \setminus \{1\} = \cup_{i=1}^m A_i) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \sum_{i=1}^m \sum_{a \in A_i} \mathbb{P}(E(a, t_0) \mid \mathcal{G}) \\ (\text{Prop. 4.4} + |A_i| \leq 2^i) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \sum_{i=1}^{+\infty} \left(\frac{2}{5} \right)^i \\ &= 1 - \frac{\mathbb{P}(\mathcal{G})}{3}. \end{aligned}$$

Since $\mathbb{P}(\mathcal{G}) \geq \alpha$ for some universal $\alpha > 0$ (cf. Proposition 4.3), we deduce:

$$\mathbb{P}_{x^{(n)}}(|S_{t_0}^A| \geq 2) \leq 1 - \gamma \text{ with } \gamma \equiv \frac{\alpha}{3} \text{ universal.}$$

This finishes the proof.

4.6 Proof of Theorem 1.2

We now present the modifications of the previous proof that are necessary to prove Theorem 1.2. We keep the definitions from previous subsections. We will also assume that $k > 4$, so that there exists some $j \in [m]$ with:

$$h \equiv 2^{j+1} - 1 = 1 + 2 + \dots + 2^j < k/2;$$

in fact, we will assume that j is the *largest* number satisfying this, so that $2^{j+2} \geq k/2$. (The case of $k \leq 4$ follows from Theorem 1.1, with an increase in the universal constant if necessary. If m is too small to allow for this choice of j , we may increase n – and thus m – at the cost of having more walkers in the beginning of the process.)

We first need an analogue of Proposition 4.1.

Proposition 4.5 (Proof omitted) *Suppose that there exists a universal $\gamma > 0$ such that for all k as above, all $n \in \mathbb{N}$ and all $x^{(n)} \in \mathbf{V}^n$,*

$$\mathbb{P}_{x^{(n)}}(|S_{t_j}^A| \geq k + 1) \leq 1 - \gamma. \tag{13}$$

Then there exists a universal $K_1 > 0$ with:

$$\mathbb{E}_{x^{(n)}}[C_k] \leq K_1 \left(\frac{T_{\text{hit}}^Q}{k} + T_{\text{mix}}^Q \right).$$

We omit the proof of this, which follows that of Proposition 4.1 quite closely. The key point is to notice that:

$$t_j = 2T_{\text{mix}}^Q + \sum_{i=j}^m 2^{4-i} (\ln 5) T_{\text{hit}}^Q \leq 2T_{\text{mix}}^Q + c_1 2^{-j-1} T_{\text{hit}}^Q \leq 2T_{\text{mix}}^Q + c_2 \frac{T_{\text{hit}}^Q}{k}$$

with $c_1, c_2 > 1$ universal (here we used $2^{j+2} \geq k/2$).

We will now bound $\mathbb{P}_{x^{(n)}}(|S_{t_j}^A| \geq k+1)$ in terms of h and k . Using Notational convention 1, we first observe that, since $h < k$:

$$\mathbb{P}\left(|S_{t_j}^A| \geq k+1\right) = \mathbb{P}\left(\sum_{a=1}^m \mathbb{I}_{E(a, t_j)} \geq k+1\right) \leq \mathbb{P}\left(\sum_{a=h+1}^m \mathbb{I}_{E(a, t_j)} \geq k-h+1\right).$$

Now follow the long chain of inequalities in the previous subsection to deduce:

$$\begin{aligned} \mathbb{P}\left(|S_{t_j}^A| \geq k+1\right) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \mathbb{P}\left(\sum_{a=h+1}^m \mathbb{I}_{E(a, t_j)} \geq k-h+1 \mid \mathcal{G}\right) \\ \text{(Markov ineq.)} &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \frac{\mathbb{E}\left[\sum_{a=h+1}^m \mathbb{I}_{E(a, t_j)}\right]}{k-h+1} \\ ([n] \setminus [h] = \cup_{i=j+1}^m A_i) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \frac{\mathbb{E}\left[\sum_{i=j+1}^m \sum_{a \in A_i} \mathbb{I}_{E(a, t_j)}\right]}{k-h} \\ &= 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \sum_{i=j+1}^m \sum_{a \in A_i} \frac{\mathbb{P}(E(a, t_j))}{k-h+1} \\ \text{(Prop. 4.4 + } |A_i| \leq 2^i) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \sum_{i=j+1}^{+\infty} \frac{2^i \left(\frac{2}{5}\right)^{j-i}}{k-h+1} \\ &= 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \frac{2^{j+1}}{3(k-h+1)} \\ (2^{j+1} = h+1) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \frac{h+1}{3(k-h+1)} \\ (h < k/2) &\leq 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \frac{k+2}{6(k-k/2)} \\ &= 1 - \mathbb{P}(\mathcal{G}) + \mathbb{P}(\mathcal{G}) \frac{k+2}{3k} \\ (k > 4) &\leq 1 - \frac{8}{15} \mathbb{P}(\mathcal{G}). \end{aligned}$$

To finish, we note that $\mathbb{P}(\mathcal{G}) \geq \alpha > 0$ with α universal (Proposition 4.3), hence we may take $\gamma = 8\alpha/15$ in Proposition 4.5.

5 On the Meeting Time Lemma

5.1 Preliminaries on quasistationary distributions

In this section we review some facts about quasistationary distributions that will be needed in the proof of Lemma 1.1. We will use the definitions of Section 2.2 throughout the section.

Given any $v \in \mathbf{V}$, we let \mathbf{q}_v be a quasistationary distribution for $\mathbf{V} \setminus \{v\}$: that is, $\mathbf{q}_v \in M_1(\mathbf{V})$ satisfies

$$\forall b \in \mathbf{V}, \mathbf{q}_v(b) = \mathbb{P}_{\mathbf{q}_v}(X_t = b \mid H_v > t).$$

All quasistationary distributions \mathbf{q}_v correspond to eigenvalues of restriction of $\Pi^{1/2}Q\Pi^{-1/2}$ to a subspace $\mathbb{R}_{-v}^{\mathbf{V}}$ of $\mathbb{R}^{\mathbf{V}}$ defined below. Here is the recipe.

1. Consider the subspace:

$$\mathbb{R}_{-v}^{\mathbf{V}} \equiv \{u \in \mathbb{R}^{\mathbf{V}} : u(v) = 0\}$$

and let $\mathcal{P}_{-v} : \mathbb{R}^{\mathbf{V}} \rightarrow \mathbb{R}_{-v}^{\mathbf{V}}$ denote the standard projection onto $\mathbb{R}_{-v}^{\mathbf{V}}$.

$$Q_{-v} \equiv \mathcal{P}_{-v}\Pi^{1/2}Q\Pi^{-1/2}\mathcal{P}_{-v}$$

is a symmetric linear operator from $\mathbb{R}_{-v}^{\mathbf{V}}$ to itself with identical diagonal entries and non-positive off-diagonal entries in the ‘‘obvious’’ basis for that space, ie. the one given by the canonical basis vectors e_b , $b \in \mathbf{V} \setminus \{v\}$.

2. By Perron-Frobenius, each irreducible block of the matrix Q_{-v} has a unique eigenvector $w_v \in \mathbb{R}_{-v}^{\mathbf{V}} \setminus \{0\}$ with non-negative entries which achieves the smallest eigenvalue $\lambda(w_v)$ corresponding to that block.
3. A simple calculation shows that the vector:

$$\mathbf{q}_v \equiv \frac{\Pi^{1/2}w_v}{\sum_{b \in \mathbf{V}} \pi^{1/2}(b)w_v(b)}$$

defines a probability distribution over \mathbf{V} with:

$$\mathbb{P}_{\mathbf{q}_v}(X_t = b, H_v > t) = \mathbf{q}_v^\dagger e^{-t\mathcal{P}_{-v}Q\mathcal{P}_{-v}} e_b = e^{-\lambda(w_v)t} \mathbf{q}_v(b),$$

which in particular implies that \mathbf{q}_v is a quasistationary distribution associated with

v . In particular, $\mathbb{E}_{\mathbf{q}_v} [H_v] = 1/\lambda(w_v) > 0$. Notice moreover that $\mathbf{q}_v(v) = 0$.

The following proposition – an immediate consequence of the third item above – will be all we need.

Proposition 5.1 *Let Q_{-v} be defined as above and let $\lambda(v)$ denote the smallest eigenvalue of Q_{-v} . Then there exists a quasistationary distribution \mathbf{q}_v for $\mathbf{V} \setminus \{v\}$ such that $\lambda(v) = 1/\mathbb{E}_{\mathbf{q}_v} [H_v]$ and:*

$$\mathbb{P}_{\mathbf{q}_v} (H_v > t) = e^{-\lambda(v)t}.$$

Proof: This smallest eigenvalue is the smallest eigenvalue of some block of Q_{-v} , and thus equals some w_v . The rest follows from item 3. and from summing the formula for $\mathbb{P}_{\mathbf{q}_v} (X_t = b, H_v > t)$ over b . \square

5.2 Proof of the Meeting Time Lemma

Proof: [of Lemma 1.1] Fix $n \in \mathbb{N} \setminus \{0\}, 0 < \Delta < n$. We note that:

$$\begin{aligned} \mathbb{P}_\pi (\forall 0 \leq s \leq t, X_s \neq h_s) &\leq \mathbb{P}_\pi (\cap_{i=1}^n \{X(it/n) \neq h(it/n)\}) \\ &\leq \mathbb{E}_\pi \left[\prod_{i=1}^n \left(1 - \frac{\Delta}{n} \mathbb{I}_{\{X(it/n)=h(it/n)\}} \right) \right]. \end{aligned}$$

For a given $v \in \mathbf{V}$, let D_v be the matrix with a 1 in position (v, v) and 0s elsewhere. A calculation reveals that the RHS above can be rewritten as:

$$(\Pi \mathbf{1})^\dagger \left\{ e^{-\frac{tQ}{n}} \left(I - \frac{\Delta D_{h(t/n)}}{n} \right) \right\} \left\{ e^{-\frac{tQ}{n}} \left(I - \frac{\Delta D_{h(2t/n)}}{n} \right) \right\} \dots \left\{ e^{-\frac{tQ}{n}} \left(I - \frac{\Delta D_{h(t)}}{n} \right) \right\} \mathbf{1}$$

where $\mathbf{1}$ is the all-ones vector and $\Pi = \text{diag}(\pi(v))_{v \in \mathbf{V}}$ was introduced in Section 5.1. Since Π commutes with all D_v , we can rewrite the above expression as:

$$(\Pi^{1/2} \mathbf{1})^\dagger \left\{ \prod_{1 \leq i \leq n}^* e^{-\frac{t \Pi^{1/2} Q \Pi^{-1/2}}{n}} \left(I - \frac{\Delta D_{h(it/n)}}{n} \right) \right\} \Pi^{1/2} \mathbf{1}$$

where the \prod^* symbol means that the order of the terms in the product is from left to right.

The vector $\Pi^{1/2}\mathbf{1}$ has norm $|\Pi^{1/2}\mathbf{1}|^2 = \sum_v \pi(v) = 1$. This implies that the above expression is at most the operator norm of the product of matrices. It follows that:

$$\mathbb{P}_\pi(\forall 0 \leq s \leq t, X_s \neq h_s) \leq \left\| \left\{ \prod_{1 \leq i \leq n}^* e^{-\frac{t\Pi^{1/2}Q\Pi^{-1/2}}{n}} \left(I - \frac{\Delta D_{h(it/n)}}{n} \right) \right\} \right\|_{\text{op}}$$

Since the operator norm is submultiplicative, we obtain:

$$\begin{aligned} \mathbb{P}_\pi(\forall 0 \leq s \leq t, X_s \neq h_s) &\leq \prod_{i=1}^n \left\| \left\{ e^{-\frac{t\Pi^{1/2}Q\Pi^{-1/2}}{n}} \left(I - \frac{\Delta D_{h(it/n)}}{n} \right) \right\} \right\|_{\text{op}} \\ &\leq \left(\max_{v \in \mathbf{V}} \left\| e^{-\frac{t\Pi^{1/2}Q\Pi^{-1/2}}{n}} \left(I - \frac{\Delta D_v}{n} \right) \right\|_{\text{op}} \right)^n. \end{aligned} \quad (14)$$

We now consider the terms of which we take the maximum in the RHS, for large $n \in \mathbb{N}$. For a given $v \in \mathbf{V}$, we have:

$$\left\| e^{-\frac{t\Pi^{1/2}Q\Pi^{-1/2}}{n}} \left(I - \frac{\Delta D_v}{n} \right) - e^{-\frac{t\Pi^{1/2}Q\Pi^{-1/2} - \Delta D_v}{n}} \right\|_{\text{op}} = O(n^{-2}),$$

where the constant implicit in the $O(n^{-2})$ term depends only on Δ , t and Q (and not on a , say). Letting $n \rightarrow +\infty$ while keeping Δ fixed, we get:

$$\begin{aligned} \lim_{n \rightarrow +\infty} \left(\max_{v \in \mathbf{V}} \left\| e^{-\frac{t\Pi^{1/2}Q\Pi^{-1/2}}{n}} \left(I - \frac{\Delta D_v}{n} \right) \right\|_{\text{op}} \right)^n \\ = \lim_{n \rightarrow +\infty} \left(\max_{v \in \mathbf{V}} \left\| e^{-\frac{t\Pi^{1/2}Q\Pi^{-1/2} - \Delta D_v}{n}} \right\|_{\text{op}} \right)^n \\ = \max_{v \in \mathbf{V}} \left\| e^{-t\Pi^{1/2}Q\Pi^{-1/2} - \Delta D_v} \right\|_{\text{op}}. \end{aligned} \quad (15)$$

Indeed, last the line follows from the self-adjointness of the exponential and from the fact that $\|B^k\|_{\text{op}} = \|B\|_{\text{op}}^k$ for self-adjoint matrices B . We now use the positive-definiteness of matrix exponentials, together with the spectral mapping property, to deduce:

$$\forall v \in \mathbf{V}, \|e^{-t\Pi^{1/2}Q\Pi^{-1/2} - \Delta D_v}\|_{\text{op}} = \lambda_{\max}(e^{-t\Pi^{1/2}Q\Pi^{-1/2} - \Delta D_v}) = e^{-\lambda_{\min}(t\Pi^{1/2}Q\Pi^{-1/2} + \Delta D_v)}.$$

This implies:

$$\mathbb{P}_\pi(\forall 0 \leq s \leq t, X_s \neq h_s) \leq \exp\left(-\min_{v \in \mathbf{V}, \Delta > 0} \lambda_{\min}(t\Pi^{1/2}Q\Pi^{-1/2} + \Delta D_v)\right).$$

We now make the following Claim.

Claim 5.1 *As $\Delta \nearrow +\infty$,*

$$\lambda_{\min}(t\Pi^{1/2}Q\Pi^{-1/2} + \Delta D_v) \rightarrow t\lambda_{\min}(Q_{-v})$$

where Q_{-v} is defined as in Section 5.1.

This result is probably well-known; for instance, it is a weaker variant of Lemma 3.1 in [5].

We will prove it below for completeness, but first we deduce from it that:

$$\mathbb{P}_\pi(\forall 0 \leq s \leq t, X_s \neq h_s) \leq \exp\left(-t \min_{v \in \mathbf{V}} \lambda_{\min}(Q_{-v})\right) = e^{-\frac{t}{\mathbb{E}_{\mathbf{q}_v}[H_v]}} = \mathbb{P}_{\mathbf{q}_v}(H_v > t)$$

via Proposition 5.1, where \mathbf{q}_v is some quasistationary distribution associated with v .

We now prove the Claim. Recall the definition of \mathcal{P}_{-v} in Section 5.1 and notice that $D_v = I - \mathcal{P}_{-v}$. This shows that $D_a w = 0$ for all $w \in \mathbb{R}_{-v}^{\mathbf{V}}$ and therefore:

$$\begin{aligned} \lambda_{\min}(t\Pi^{1/2}Q\Pi^{-1/2} + \Delta D_v) &= \inf_{w \in \mathbb{R}^{\mathbf{V}}, |w|=1} w^\dagger(t\Pi^{1/2}Q\Pi^{-1/2} + \Delta D_v) \\ &\leq \inf_{w \in \mathbb{R}_{-v}^{\mathbf{V}}, |w|=1} w^\dagger(t\Pi^{1/2}Q\Pi^{-1/2} + \Delta D_v)w \\ &= \inf_{w \in \mathbb{R}_{-v}^{\mathbf{V}}, |w|=1} w^\dagger(t\Pi^{1/2}Q\Pi^{-1/2})w \\ \text{(use } \mathcal{P}_{-v}w = w) &= t \inf_{w \in \mathbb{R}_{-v}^{\mathbf{V}}, |w|=1} w^\dagger(\mathcal{P}_{-v}\Pi^{1/2}Q\Pi^{-1/2}\mathcal{P}_{-v})w \\ &= t\lambda_{\min}(Q_{-a}). \end{aligned} \tag{16}$$

To get an opposite inequality, we set $A = t\Pi^{1/2}Q\Pi^{-1/2}$ for convenience. We first show that there exists some $c > 0$ such that for all large enough $\Delta > 0$,

$$A + \Delta D_v \succeq \mathcal{P}_{-v}A\mathcal{P}_{-v} - \frac{c\mathcal{P}_{-v}}{\Delta} + \frac{\Delta}{2}D_v, \tag{17}$$

where for symmetric matrices B_1, B_2 with the same size, $B_1 \preceq B_2$ means that $B_2 - B_1$ is

positive semidefinite. To see this, we use $\mathcal{P}_{-v} + D_v = I$ several times and notice that for any $x \in \mathbb{R}^{\mathbf{V}}$,

$$\begin{aligned}
x^\dagger(A + \Delta D_v)x &= x^\dagger \mathcal{P}_{-v} A \mathcal{P}_{-v} x + x^\dagger D_v A D_v x + 2x^\dagger (\mathcal{P}_{-v} A D_v)x \\
&\quad + \Delta (x^\dagger D_v x) \\
\text{(Cauchy-Schwartz)} &\geq x^\dagger \mathcal{P}_{-v} A \mathcal{P}_{-v} x + x^\dagger \frac{\Delta D_v}{2} x + |D_v x|^2 \left(\frac{\Delta}{2} - \|A\|_{\text{op}} \right) \\
&\quad - 2\|A\|_{\text{op}} |\mathcal{P}_{-v} x| |D_v x| \\
\text{(assume } \Delta > 4\|A\|_{\text{op}}) &\geq x^\dagger \left(\mathcal{P}_{-v} A \mathcal{P}_{-v} + \frac{\Delta D_v}{2} \right) x \\
&\quad + \left(\frac{\sqrt{\Delta} |D_v x|}{2} - \frac{2\|A\|_{\text{op}} |\mathcal{P}_{-v} x|}{\sqrt{\Delta}} \right)^2 - \frac{4\|A\|_{\text{op}} |\mathcal{P}_{-v} x|^2}{\Delta} \\
\text{(set } c \equiv 4\|A\|_{\text{op}}) &\geq x^\dagger \left(\mathcal{P}_{-v} A \mathcal{P}_{-v} - \frac{c \mathcal{P}_{-v}}{\Delta} + \frac{\Delta D_v}{2} \right) x.
\end{aligned}$$

This proves (17), which implies:

$$\lambda_{\min}(A + \Delta D_v) \geq \lambda_{\min} \left(\mathcal{P}_{-v} A \mathcal{P}_{-v} - \frac{c \mathcal{P}_{-v}}{\Delta} + \frac{\Delta}{2} D_v, \right). \quad (18)$$

Notice that the matrix in the RHS has $\mathbb{R}_{-v}^{\mathbf{V}}$ as an invariant subspace, which implies that all of its eigenvectors lie in $\mathbb{R}_{-v}^{\mathbf{V}}$ or in its orthogonal complement. It is easy to see that the all vectors in the latter space are eigenvectors with eigenvalue $\Delta/2$; therefore, for all large enough Δ the *minimal* eigenvalue corresponds to a vector in $\mathbb{R}_{-v}^{\mathbf{V}}$. We deduce that for all large $\Delta > 0$,

$$\begin{aligned}
\lambda_{\min} \left(\mathcal{P}_{-v} A \mathcal{P}_{-v} - \frac{c \mathcal{P}_{-v}}{\Delta} + \frac{\Delta}{4} D_v, \right) \\
= \min_{w \in \mathbb{R}_{-v}^{\mathbf{V}}, |w|=1} w^\dagger \left(\mathcal{P}_{-v} A \mathcal{P}_{-v} - \frac{c \mathcal{P}_{-v}}{\Delta} \right) w = t \lambda_{\min}(Q_{-v}) - \frac{2c}{\Delta}
\end{aligned}$$

because $w^\dagger \mathcal{P}_{-v} A \mathcal{P}_{-v} w = t w^\dagger Q_{-v} w$ for all w as above. Together with (16) and (18), this shows that:

$$\text{For large enough } \Delta > 0, t \lambda_{\min}(Q_{-v}) - \frac{c}{\Delta} \leq \lambda_{\min}(t \Pi^{1/2} Q \Pi^{-1/2} + \Delta D_v) \leq t \lambda_{\min}(Q_{-v}),$$

and the Claim follows when we let $\Delta \nearrow +\infty$. \square

6 Final remarks

- Let T_{meet}^Q denote the maximum expected meeting time of two independent realizations of Q . In light of the discussion in the Introduction, it would be natural to expect that $\mathbb{E}[C] \leq K_2 T_{\text{meet}}^Q$ for some universal $K_2 > 0$ and all Q . Is this actually true? A more modest question is whether the constants in the two Theorems can be improved.
- The Meeting Time Lemma (Lemma 1.1) can be used in the study of a cat-and-mouse game proposed in [2, Chapter 4, page 17]. In this game a cat moves according to a reversible Markov chain Q . A mouse chooses a trajectory $(h_s)_{s \geq 0}$ for itself and an initial distribution for the cat. Aldous and Fill asked if staying put at some carefully chosen state gives an optimal strategy for the mouse in terms of maximizing $\mathbb{E}[M]$, where M is the meeting time of cat and mouse. One can use Lemma 1.1 to prove that if $T_{\text{mix}}^Q \ll \max_v \mathbb{E}_\pi[H_v]$ (a natural condition in many examples), then the strategy where the mouse stays at v and chooses \mathbf{q}_v as the initial distribution nearly maximizes $\mathbb{P}(M > t)$ simultaneously for all $t \geq 0$. We expect to comment on this and related results in a upcoming note.

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