### Asymptotic Learning on Social Networks

Elchanan Mossel Joint work with: Allan Sly (UC Berkeley) Omer Tamuz (Weizmann)

February 23, 2012

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1/25

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- Question: Do agents in this decentralized model aggregate their information effectively?

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3/25

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  - ▶ F<sup>-</sup> if S = −1
  - ► *F*<sup>+</sup> if *S* = 1
- Conditioned on S, private signals are independent.



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4/25

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  - When do the agents learn from each other efficiently?
    - Generally poorly understood.

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- ► Challenge: G = {0,1}<sup>3</sup> with signals 0.11, 0.22, 0.33, 0.44, 0.55, 0.66, 0.77, 0.88.
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- Does it converge in a finite number of iterations?

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► Note: Proof doesn't require independent signals.

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- Stronger version (M-Sly-Tamuz-12): Under the non-atomic beliefs three possible limiting actions are possible

1. For all 
$$i$$
,  $A_i(t) \rightarrow 1$  and  $X_i(\infty) > \frac{1}{2}$ .

- 2. For all i,  $A_i(t) \rightarrow -1$  and  $X_i(\infty) < \frac{1}{2}$ .
- 3. For all *i*,  $A_i(t)$  does not converge and  $X_i(\infty) = \frac{1}{2}$ .

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- Aumann's original motivation: Bayesian Economics doesn't make sense. Doesn't allow to "agree to disagree".

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  - "different motivation is simply technical expediency" (Ellison and Fundenberg)
  - "to keep the model mathematically tractable... this possibility [fully Bayesian agents] is precluded in our model... simplifying the belief revision process considerably" (Bala and Goyal)

Belief Learning Theorem (M., Sly and Tamuz (2012)) In the revealed beliefs model the limit  $X = \lim_{t\to\infty} X_i(t)$  satisfies:

$$X = \mathbb{P}\left[S = 1 \mid \omega_1, \ldots, \omega_n\right],$$

In particular there exist c > 0,

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- Independence of signals is needed. Example:  $S = S_1 \oplus S_2$ .

► Recall that in the revealed actions model:  $A_i(t) = \operatorname{argmax}_{a \in \{-1,1\}} \mathbb{P}\left[a = S | \mathcal{F}_i(t)\right] \text{ is announced to the neighbors at time } t.$ 

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- False: without non-atomic assumption, on directed graphs, w.o independence.

#### Prior work - some tractable models

- In the revealed actions model when the social network is the complete graph then Bayesian updates are tractable (Mossel and Tamuz (2010)).
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- When calculations are tractable then they are also easier to analyze.

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False in general without non-atomic assumption and on directed graphs, independence is also needed.

Dynamics are very complicated. Abstract approach needed: Assume by contradiction there is a sequence G<sub>n</sub> = (V<sub>n</sub>, E<sub>n</sub>) of graphs with |V<sub>n</sub>| → ∞ and lim sup P [Learning in G<sub>n</sub>] < 1.</p>

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- $\blacktriangleright \implies G_n$  are bounded degree.

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- Let

$$p^* = \inf_{G \text{ infinite}} p(G).$$

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- More work needed for general graphs.

The proof proceeds by using induction to find a vertex *i* and times t<sub>1</sub> < ... < t<sub>k</sub> such that ℙ[A<sub>i</sub>(t<sub>ℓ</sub>) = S] ≈ p\* and the A<sub>i</sub>(t<sub>ℓ</sub>) are almost independent.

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- Inductive Hypothesis: On any infinite graph G for any k, € > 0 there exists a vertex i such by some time t there exist F<sub>i</sub>(t) measureable random variables Y<sub>1</sub>,..., Y<sub>k</sub> taking values in {−1,1} such that
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- Note that the case k = 1 follows from the definition of  $p^*$ .
- ► The induction claim implies the theorem by taking the majority of the Y<sub>ℓ</sub>. This identifies S with probability better than p<sup>\*</sup> unless p<sup>\*</sup> = 1.

▶ Note for any *i* and any  $\epsilon' > 0$ , there exists t' and an  $\mathcal{F}_{t'}$ -measurable  $\tilde{A^*}$  such that  $\mathbb{P}\left[A^* = \tilde{A^*}\right] \ge 1 - \epsilon'$ .

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But eventually agent j will learn A\* too giving j another informative time. This completes the induction. ► Finite to infinite principle.

- Finite to infinite principle.
- Like many finite to infinite proofs, no rate.

Belief Learning Theorem (M. Sly and Tamuz (2012)) If there exists a random variable X such that  $X = X_i := \mathbb{E} [S \mid \mathcal{F}_i(\infty)]$  for all *i* then all agents learned optimally:

$$X = \mathbb{P}\left[S = 1 \mid \omega_1, \ldots, \omega_n\right].$$

# Proof Sketch

$$Z_i := \log \frac{\mathbb{P}\left[S = 1 \mid \omega_i\right]}{\mathbb{P}\left[S = 0 \mid \omega_i\right]} = \log \frac{\mathbb{P}\left[\omega_i \mid S = 1\right]}{\mathbb{P}\left[\omega_i \mid S = 0\right]}, \quad Z = \sum_i Z_i$$

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$$\mathbb{P}\left[S=1\mid\omega_{1},\ldots,\omega_{n}\right]=L(Z)$$
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• Hence since  $Z_i$  is  $\mathcal{F}_i$  measurable

$$\mathbb{E} \left[ Z_i \cdot L(Z) \mid X \right] = \mathbb{E} \left[ \mathbb{E} \left[ Z_i \cdot L(Z) \mid \mathcal{F}_i \right] \mid X \right]$$
$$= \mathbb{E} \left[ Z_i \cdot X \mid X \right]$$
$$= \mathbb{E} \left[ Z_i \mid X \right] \mathbb{E} \left[ L(Z) \mid X \right]$$

22 / 25

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23/25

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So the agreed value X equal to the optimal estimator L(Z) as needed.

Both models: Rate of convergence? Dependence on graph?

- ▶ Both models: Rate of convergence? Dependence on graph?
- Actions Learning: how does the probability of learning change with the graph?

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- Actions Learning: how does the probability of learning change with the graph?
- What if the agents aren't truthful but act strategically (in a game theoretic manner)?

Questions?