

$$A = \begin{pmatrix} -1 & -3 \\ -2 & 2 \end{pmatrix} \begin{matrix} \text{row min} \\ -3 \\ -2 \end{matrix} \Rightarrow \text{No saddle point}$$

col max -1 2

Let the optimal strategy for player I be  $(p, 1-p)$

$$\Rightarrow -p - 2(1-p) = -3p + 2(1-p) \quad (\text{using equalising strategies})$$

$$\Rightarrow -p + 2p - 2 = -3p + 2 - 2p$$

$$6p = 4 \quad p = 4/6 = 2/3, \quad 1-p = 1/3$$

Similarly, let  $(q, 1-q)^T$  be optimal for II under pure strategies  $e_i$  from player I,

$$-q - 3(1-q) = -2q + 2(1-q)$$

$$-q - 3 + 3q = -2q + 2 - 2q$$

$$6q = 5, \quad q = 5/6 \quad (5/6, 1/6)$$

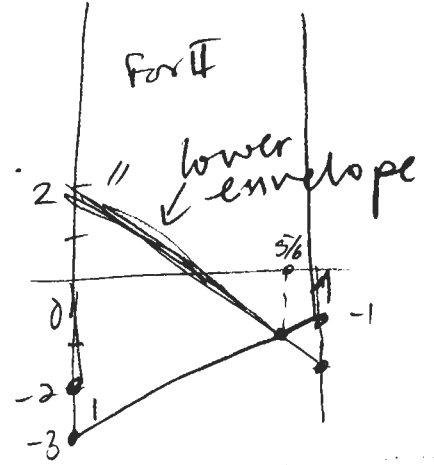
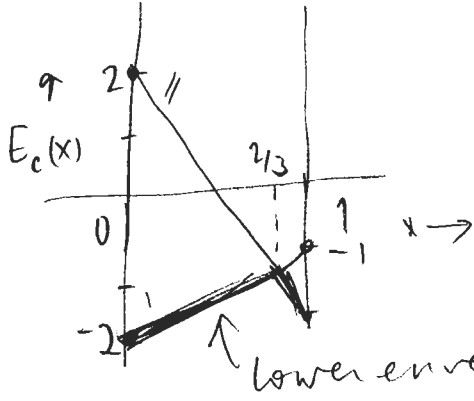
$$v = -\frac{2}{3} - 2\left(\frac{1}{3}\right) = -\frac{4}{3}$$

$$\text{or } v = -\frac{5}{6} - \frac{3}{6} = -\frac{8}{6} = -\frac{4}{3}$$

Ans.  $(2/3, 1/3)$  for player I

Alternatively : Graphical method

$$A = \begin{matrix} & y & 1-y \\ x & \begin{pmatrix} -1 & -3 \\ -2 & 2 \end{pmatrix} \\ 1-x & & \end{matrix}$$

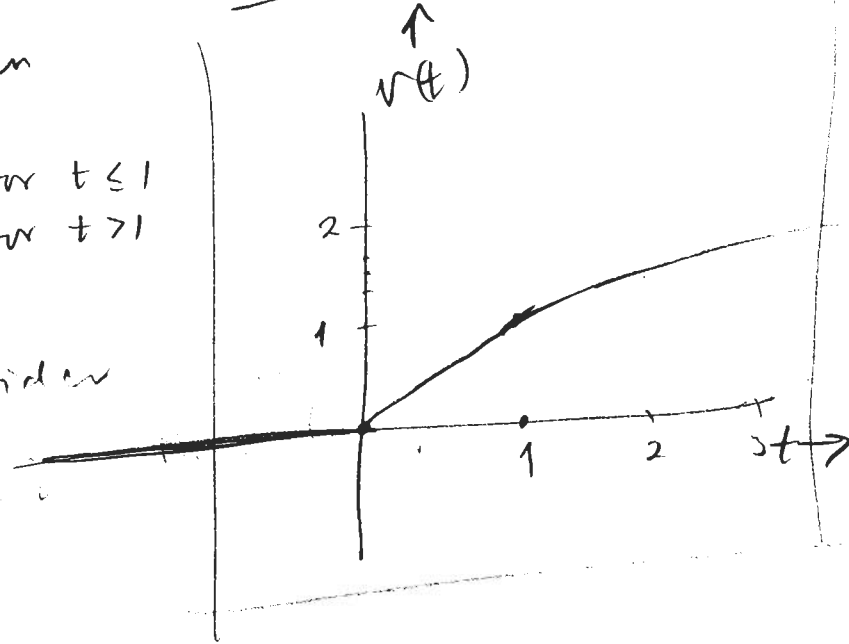


$\frac{2}{II-14}$

$$A = \begin{pmatrix} 0 & 2 \\ t & 1 \end{pmatrix} \begin{matrix} \text{row min} \\ 0 \\ \begin{cases} t & \text{for } t \leq 1 \\ 1 & \text{for } t > 1 \end{cases} \end{matrix}$$

$$\text{col max } \begin{cases} t, t \geq 0 \\ 0, t < 0 \end{cases} \begin{matrix} 2 \\ 0 \end{matrix}$$

- ⇒ We have to consider 3 cases:
- (i)  $t < 0$
  - (ii)  $0 \leq t \leq 1$
  - (iii)  $t > 1$



Case (i)

$$t < 0 \begin{matrix} \text{row min} \\ \begin{pmatrix} 0 & 2 \\ t & 1 \end{pmatrix} \\ \text{col max } \begin{matrix} 0 \\ 2 \end{matrix} \end{matrix} \begin{matrix} 0 \\ t \end{matrix}$$

⇒ 0 is a saddle point &  $v(t) = 0$

Case (ii)  $0 \leq t \leq 1$

$$\begin{matrix} \text{row min} \\ \begin{pmatrix} 0 & 2 \\ t & 1 \end{pmatrix} \\ \text{col max } \begin{matrix} 0 \\ t \end{matrix} \end{matrix} \begin{matrix} 0 \\ t \end{matrix}$$

⇒ t is a saddle point &  $v(t) = t$

Case (iii)  $t > 1$

$$\begin{matrix} \text{row min} \\ \begin{pmatrix} 0 & 2 \\ t & 1 \end{pmatrix} \\ \text{col max } \begin{matrix} 0 \\ 1 \end{matrix} \end{matrix} \begin{matrix} 0 \\ 1 \end{matrix}$$

no saddle point

$$t(1-p) = 2p + 1-p$$

$$\Rightarrow p = \frac{t-1}{t+1}$$

$$v(t) = \frac{2t}{t+1}$$

$$\lim_{t \rightarrow \infty} \frac{2t}{t+1} = \frac{2}{1+1/t} = 2$$

3

3.2  
P70

$$A = \begin{pmatrix} 0 & 9 & 1 & 1 \\ 5 & 0 & 6 & 7 \\ 2 & 4 & 3 & 3 \end{pmatrix} \rightarrow \begin{pmatrix} 0 & 9 \\ 5 & 0 \\ 2 & 4 \end{pmatrix} \rightarrow \begin{pmatrix} 0 & 9 \\ 5 & 0 \end{pmatrix} \quad \textcircled{c}$$

Note that col 3 & col 4 dominate col 1, so they can be discarded.

Now note that  $\frac{1}{2}(a_{11} + a_{21}) \geq 2$  &  $\frac{1}{2}(a_{12} + a_{22}) \geq 4$   
 $\Rightarrow$  a convex combination of rows 1 & 2 dominates row 3.

Now we can solve this  $2 \times 2$  matrix to get

$$p = \begin{pmatrix} 5/14 \\ 9/14 \end{pmatrix} = q, \quad v = 45/14$$

4

3.3  
P70

There are two possible ways to solve this problem.

① Use a payoff mix of probabilities  $\begin{pmatrix} 3/4 & 0 & 0 \\ 0 & 1/4 & 0 \\ 0 & 0 & 1/2 \end{pmatrix}$  where Player I wants to maximise his chances of destroying the item.

Solve to get  $p = (2/11, 6/11, 3/11)^T = q$  &  $v = \frac{6}{44}$ .

Note  $v < \frac{1}{2}$  & it is a probability, so game favours PI.

② Suppose PI wins \$1 for destroying item & loses \$1 for failing. Then payoff matrix =  $\begin{pmatrix} 1/2 & -1 & -1 \\ -1 & -1/2 & -1 \\ -1 & -1 & 0 \end{pmatrix}$  where the entries are his expected gains under the given probabilities

( $a_{11} = (\frac{3}{4})(1) + (\frac{1}{4})(-1) = \frac{1}{2}$   
 $a_{12} = (+1)(-1)$  etc.)  
 Solving this game gives the same optimal strategy &  $v = -8/11$ , so game favours PII.