Topics in Prediction and Learning Lecture 1: Optimal Universal Prediction—Quadratic Loss

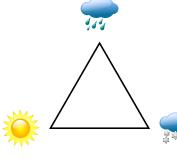
Peter Bartlett

Computer Science and Statistics University of California at Berkeley

Mathematical Sciences Queensland University of Technology

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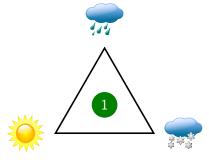
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At round *t*:

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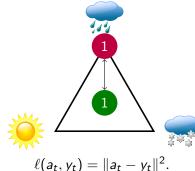
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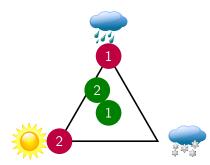
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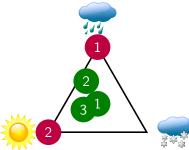
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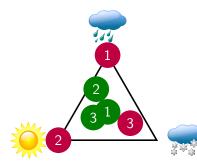
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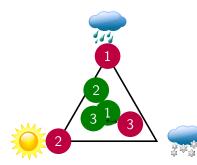
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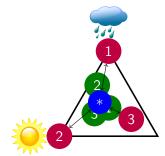
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Player's aim:

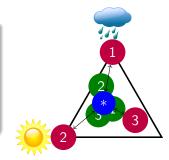
Minimize regret:

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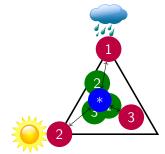
Minimize regret wrt comparison C:

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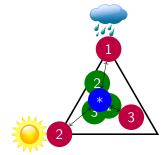
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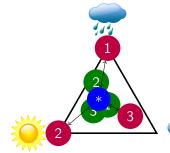
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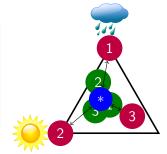
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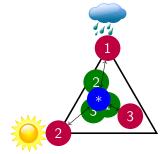
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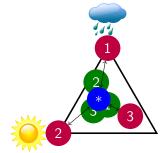
Examples:

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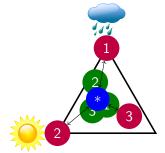
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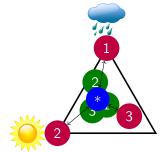
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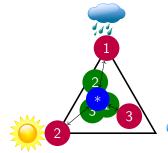
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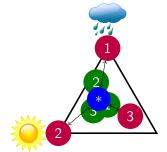
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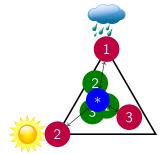
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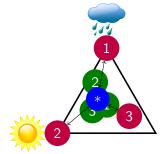
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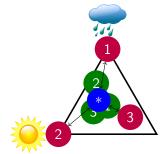
Examples:

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- Focus in this lecture: Minimax optimal strategies.

$$\sum_{t=1}^{T} \ell(a_t, y_t) - \min_{a \in \mathcal{A}} \sum_{t=1}^{T} \ell(a, y_t)$$

Minimax Regret

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 $\min_{a_1 \in \mathcal{A}}$

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The value of the game: Minimax Regret

$$V_T(\mathcal{Y}, \mathcal{A}) = \min_{a_1 \in \mathcal{A}} \max_{y_1 \in \mathcal{Y}} \cdots \min_{a_T \in \mathcal{A}} \max_{y_T \in \mathcal{Y}} \left(\sum_{t=1}^T \ell(a_t, y_t) - \min_{a \in \mathcal{A}} \sum_{t=1}^T \ell(a, y_t) \right)$$

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Strategy:

$$S: \bigcup_{t=0}^{T} \mathcal{Y}^t \to \mathcal{A}.$$

$$V_{T}(\mathcal{Y}, \mathcal{A}) = \min_{S} \max_{y_{1}^{T} \in \mathcal{Y}^{T}} \left(\sum_{t=1}^{T} \ell\left(S\left(y_{1}^{t-1}\right), y_{t}\right) - \min_{a \in \mathcal{A}} \sum_{t=1}^{T} \ell(a, y_{t}) \right)$$

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Minimax Optimal Strategy:

$$\begin{split} \mathcal{S}^* : & \bigcup_{t=0}^T \mathcal{Y}^t \to \mathcal{A}. \\ V_T(\mathcal{Y}, \mathcal{A}) &= \min_{\mathcal{S}} \max_{y_1^T \in \mathcal{Y}^T} \left(\sum_{t=1}^T \ell\left(\mathcal{S}\left(y_1^{t-1}\right), y_t\right) - \min_{a \in \mathcal{A}} \sum_{t=1}^T \ell(a, y_t) \right) \\ &= \max_{y_1^T \in \mathcal{Y}^T} \left(\sum_{t=1}^T \ell\left(\mathcal{S}^*\left(y_1^{t-1}\right), y_t\right) - \min_{a \in \mathcal{A}} \sum_{t=1}^T \ell(a, y_t) \right). \end{split}$$



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loss, $\ell(a, y)$:

1 $||a-y||_2^2$,

 $a, y \in \mathbb{R}^d$.

2 $(x^{\top}a - y)^2$.

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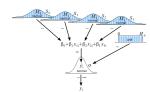
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Outline

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- Part 1: Euclidean loss.
- Part 2: Linear regression.

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$$V_T(\mathcal{Y}, \mathcal{A}) = \min_{a_1 \in \mathcal{A}} \max_{y_1 \in \mathcal{Y}} \cdots \min_{a_T \in \mathcal{A}} \max_{y_T \in \mathcal{Y}} \left(\sum_{t=1}^T \ell(a_t, y_t) - \min_{a \in \mathcal{A}} \sum_{t=1}^T \ell(a, y_t) \right).$$

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$$V(y_1, ..., y_T) := -\min_{a} \sum_{t=1}^{T} \ell(a, y_t),$$

$$V(y_1, ..., y_{t-1}) := \min_{a_t} \max_{y_t} (\ell(a_t, y_t) + V(y_1, ..., y_t)),$$

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$$V_{T}(\mathcal{Y}, \mathcal{A}) = \min_{a_{1} \in \mathcal{A}} \max_{y_{1} \in \mathcal{Y}} \cdots \min_{a_{T} \in \mathcal{A}} \max_{y_{T} \in \mathcal{Y}} \left(\sum_{t=1}^{I} \ell(a_{t}, y_{t}) - \min_{a \in \mathcal{A}} \sum_{t=1}^{I} \ell(a, y_{t}) \right).$$

$$egin{aligned} V(y_1, \dots, y_T) &:= -\min_{a} \sum_{t=1}^T \ell(a, y_t), \ V(y_1, \dots, y_{t-1}) &:= \min_{a_t} \max_{y_t} \left(\ell(a_t, y_t) + V(y_1, \dots, y_t) \right), \ V_T(\mathcal{Y}, \mathcal{A}) &= V(), \end{aligned}$$

The value of the game:

$$V_{\mathcal{T}}(\mathcal{Y}, \mathcal{A}) = \min_{a_1 \in \mathcal{A}} \max_{y_1 \in \mathcal{Y}} \cdots \min_{a_{\mathcal{T}} \in \mathcal{A}} \max_{y_{\mathcal{T}} \in \mathcal{Y}} \left(\sum_{t=1}^{I} \ell(a_t, y_t) - \min_{a \in \mathcal{A}} \sum_{t=1}^{I} \ell(a, y_t) \right).$$

$$\begin{split} V(y_1, \dots, y_T) &:= -\min_{a} \sum_{t=1}^T \ell(a, y_t), \\ V(y_1, \dots, y_{t-1}) &:= \min_{a_t} \max_{y_t} \left(\ell(a_t, y_t) + V(y_1, \dots, y_t) \right), \\ V_T(\mathcal{Y}, \mathcal{A}) &= V(), \\ S^*(y_1, \dots, y_{t-1}) &= \arg\min_{a_t} \max_{y_t} \left(\ell(a_t, y_t) + V(y_1, \dots, y_t) \right). \end{split}$$

To play the minimax strategy: after seeing y_1, \ldots, y_{t-1} ,

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Efficient minimax optimal strategies

When is V a simple function of (statistics of) the history y_1, \ldots, y_t ?

Outline

- Computing minimax optimal strategies.
- Part 1: Euclidean loss.
 - ullet $\mathcal{Y} = \mathsf{ball}$
 - \bullet $\mathcal{Y} = simplex$
 - ullet Closed, bounded ${\mathcal Y}$
 - Hilbert space
 - $\mathcal{Y} = \mathsf{ellipsoid}$
- Part 2: Linear regression.

Euclidean loss

$$\ell(\hat{y}, y) = \frac{1}{2} \|\hat{y} - y\|^2.$$

Constraints

Strategy chooses $\hat{\mathbf{y}}_n \in \mathbb{R}^d$.

Euclidean loss
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Regret
$$=\sum_{t=1}^{n}\ell(\hat{y}_t,y_t)$$
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Strategy chooses $\hat{y}_n \in \mathbb{R}^d$.

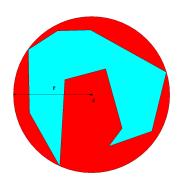
Adversary chooses $y_n \in \mathcal{Y}$, where $\mathcal{Y} \subseteq \mathbb{R}^d$.

$$\ell(\hat{y}, y) = \frac{1}{2} \|\hat{y} - y\|^2.$$

Regret
$$=\sum_{t=1}^n \ell(\hat{y}_t, y_t) - \inf_{a \in \mathbb{R}^d} \sum_{t=1}^n \ell(a, y_t).$$

The smallest ball containing \mathcal{Y} is $B_{\mathcal{Y}} = \{y \in \mathbb{R}^d : \|y - c\| \le r\}$, with center $c = \arg\min_{c} \max_{y \in \mathcal{Y}} \|y - c\|$, radius $r = \min_{c} \max_{y \in \mathcal{Y}} \|y - c\|$.

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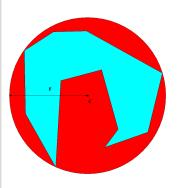


The smallest ball containing \mathcal{Y} is $B_{\mathcal{Y}} = \{ y \in \mathbb{R}^d : ||y - c|| \le r \}$, with center $c = \arg\min_c \max_{v \in \mathcal{Y}} ||y - c||$, radius $r = \min_c \max_{v \in \mathcal{Y}} ||y - c||$.

Theorem

• For closed, bounded $\mathcal{Y} \subset \mathbb{R}^d$, minimax strategy is empirical minimizer, shrunk towards c:

$$a_{t+1}^* = t\alpha_{t+1}\bar{y}_t + (1 - t\alpha_{t+1})c.$$



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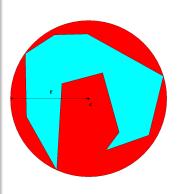
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• Optimal regret: $\frac{r^2}{2} \sum_{t=1}^{l} \alpha_t$;

$$\alpha_T = \frac{1}{T}, \quad \alpha_t = \alpha_{t+1}^2 + \alpha_{t+1}.$$



Outline

- Computing minimax optimal strategies.
- Part 1: Euclidean loss.
 - $\mathcal{Y} = \mathbf{ball}$
 - $\mathcal{Y} = \text{simplex}$
 - ullet Closed, bounded ${\cal Y}$
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The ball case: $\mathcal{Y} = \{y : ||y - c|| \le r\}$

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Maintain statistics:
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Value-to-go: quadratic

$$\frac{1}{2} \left(\alpha_n \|s_n\|^2 - \sigma_n^2 + r^2 \sum_{t=n+1}^T \alpha_t \right).$$

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Minimax strategy: affine

$$a_n^* = c + \alpha_n \sum_{t=1}^{n-1} (y_t - c).$$

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Minimax regret for ball

$$V(\mathcal{Y}) = \frac{r^2}{2} \sum_{t=0}^{T} \alpha_t.$$

The ball case: $\mathcal{Y} = \{y : ||y - c|| \le r\}$

Maintain statistics:
$$s_n = \sum_{t=1}^n (y_t - c), \qquad \sigma_n^2 = \sum_{t=1}^n \|y_t - c\|^2.$$

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$$\frac{1}{2}\left(\alpha_n\|s_n\|^2 - \sigma_n^2 + r^2\sum_{t=n+1}^T \alpha_t\right).$$

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Maximin distribution: same mean

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= $(n-1)\alpha_n \bar{y}_{n-1} + (1-(n-1)\alpha_n)$

Maximin distribution: same mean.

Minimax regret for ball

$$V(\mathcal{Y}) = \frac{r^2}{2} \sum_{t=0}^{T} \alpha_t.$$

$$V(y_1, ..., y_T) := -\min_{a} \sum_{t=1}^{I} \ell(a, y_t),$$

$$V(y_1, ..., y_{t-1}) = \min_{a_t} \max_{y_t} (\ell(a_t, y_t) + V(y_1, ..., y_t)).$$

Proof idea

$$V(y_1, ..., y_T) := -\min_{a} \sum_{t=1}^{T} \ell(a, y_t),$$

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The final $V(y_1, ..., y_T)$ is a (convex) quadratic in the state:

Proof idea

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$$V(y_1,\ldots,y_T) := -\min_{a} \frac{1}{2} \sum_{t=1}^{I} \|a - y_t\|^2$$

Proof idea

$$V(y_1, ..., y_T) := -\min_{a} \sum_{t=1}^{T} \ell(a, y_t),$$

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$$V(y_1, \dots, y_T) := -\min_{a} \frac{1}{2} \sum_{t=1}^{T} \|a - y_t\|^2 = -\frac{1}{2} \sum_{t=1}^{T} \|\bar{y} - y_t\|^2$$

Proof idea

$$V(y_1, ..., y_T) := -\min_{a} \sum_{t=1}^{T} \ell(a, y_t),$$

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$$= -\frac{1}{2} \sum_{t=1}^{T} \|\frac{s_T}{T} - y_t\|^2$$

Proof idea

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$$V(y_1, \dots, y_T) := -\min_{a} \frac{1}{2} \sum_{t=1}^{T} \|a - y_t\|^2 = -\frac{1}{2} \sum_{t=1}^{T} \|\bar{y} - y_t\|^2$$
$$= -\frac{1}{2} \sum_{t=1}^{T} \|\frac{s_T}{T} - y_t\|^2 = \frac{1}{2} \left(\frac{1}{T} \|s_T\|^2 - \sigma_T^2\right).$$

Proof idea

$$V(y_1, ..., y_T) := -\min_{a} \sum_{t=1}^{T} \ell(a, y_t),$$

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$$\begin{split} V(y_1,\ldots,y_T) &:= -\min_{a} \frac{1}{2} \sum_{t=1}^{T} \|a - y_t\|^2 = -\frac{1}{2} \sum_{t=1}^{T} \|\bar{y} - y_t\|^2 \\ &= -\frac{1}{2} \sum_{t=1}^{T} \|\frac{s_T}{T} - y_t\|^2 = \frac{1}{2} \left(\frac{1}{T} \|s_T\|^2 - \sigma_T^2\right). \end{split}$$
 i.e., $V(s_T,\sigma_T^2,T) = \frac{1}{2} \left(\alpha_T \|s_T\|^2 - \sigma_T^2\right).$

$$V(y_1,...,y_{t-1}) = \min_{a_t} \max_{y_t} (\ell(a_t, y_t) + V(y_1,...,y_t))$$

$$V(y_1, \dots, y_{t-1}) = \min_{\substack{a_t \\ y_t}} \max_{y_t} (\ell(a_t, y_t) + V(y_1, \dots, y_t))$$

$$= \frac{1}{2} \min_{\substack{a_t \\ y_t}} \max_{y_t} (\|a_t - y_t\|^2 + \alpha_t \|s_{t-1} + y_t\|^2$$

$$- \sigma_{t-1}^2 - \|y_t\|^2 + r^2 \sum_{n=t+1}^T \alpha_n$$

$$\begin{split} V(y_1, \dots, y_{t-1}) &= \min_{\substack{a_t \\ a_t \\ y_t}} \max_{y_t} \left(\|(a_t, y_t) + V(y_1, \dots, y_t) \right) \\ &= \frac{1}{2} \min_{\substack{a_t \\ y_t}} \max_{y_t} \left(\|(a_t - y_t)\|^2 + \alpha_t \|(s_{t-1} + y_t)\|^2 \right. \\ &\left. - \sigma_{t-1}^2 - \|(y_t)\|^2 + r^2 \sum_{n=t+1}^T \alpha_n \right) \\ &= \frac{1}{2} \min_{\substack{a_t \\ a_t \\ y_t}} \max_{y_t} \left(\|(a_t)\|^2 + 2 \left(\alpha_t s_{t-1} - a_t \right)^\top y_t + \alpha_t \|(y_t)\|^2 \right. \\ &\left. + \alpha_t \|(s_{t-1})\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right) \end{split}$$

Proof idea

$$\begin{split} V(y_1, \dots, y_{t-1}) &= \min_{\substack{a_t \\ a_t \\ y_t}} \max_{y_t} \left(\|(a_t, y_t) + V(y_1, \dots, y_t) \right) \\ &= \frac{1}{2} \min_{\substack{a_t \\ a_t \\ y_t}} \max_{y_t} \left(\|a_t - y_t\|^2 + \alpha_t \|s_{t-1} + y_t\|^2 \right. \\ &\left. - \sigma_{t-1}^2 - \|y_t\|^2 + r^2 \sum_{n=t+1}^T \alpha_n \right) \\ &= \frac{1}{2} \min_{\substack{a_t \\ a_t \\ y_t}} \max_{y_t} \left(\|a_t\|^2 + 2 \left(\alpha_t s_{t-1} - a_t \right)^\top y_t + \alpha_t \|y_t\|^2 \right. \\ &\left. + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right) \end{split}$$

Optimization of y_t : maximize a convex function over the ball.

Proof idea

$$V(y_1, \dots, y_{t-1}) = \min_{\substack{a_t \\ y_t}} \max_{y_t} (\ell(a_t, y_t) + V(y_1, \dots, y_t))$$

$$= \frac{1}{2} \min_{\substack{a_t \\ y_t}} \max_{y_t} (\|a_t - y_t\|^2 + \alpha_t \|s_{t-1} + y_t\|^2$$

$$- \sigma_{t-1}^2 - \|y_t\|^2 + r^2 \sum_{n=t+1}^T \alpha_n$$

$$= \frac{1}{2} \min_{\substack{a_t \\ a_t}} \max_{y_t} \left(\|a_t\|^2 + 2(\alpha_t s_{t-1} - a_t)^\top y_t + \alpha_t \|y_t\|^2 + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right)$$

Optimization of y_t : maximize a convex function over the ball.

But the solution is easy: choose y_t on the sphere, aligned with $\alpha_t s_{t-1} - a_t$.

$$V(y_1, \dots, y_{t-1}) = \frac{1}{2} \min_{a_t} \max_{y_t} \left(\|a_t\|^2 + 2(\alpha_t s_{t-1} - a_t)^\top y_t + \alpha_t \|y_t\|^2 + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right)$$

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$$= \frac{1}{2} \min_{a_t} \left(\|a_t\|^2 + 2r \|\alpha_t s_{t-1} - a_t\| + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right)$$

$$V(y_1, \dots, y_{t-1}) = \frac{1}{2} \min_{a_t} \max_{y_t} \left(\|a_t\|^2 + 2 \left(\alpha_t s_{t-1} - a_t \right)^\top y_t + \alpha_t \|y_t\|^2 + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right)$$

$$= \frac{1}{2} \min_{a_t} \left(\|a_t\|^2 + 2r \|\alpha_t s_{t-1} - a_t\| + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right)$$

$$(a_t^* := \alpha_t s_{t-1})$$

$$\begin{split} V(y_1, \dots, y_{t-1}) &= \frac{1}{2} \min_{a_t} \max_{y_t} \left(\|a_t\|^2 + 2 \left(\alpha_t s_{t-1} - a_t \right)^\top y_t + \alpha_t \|y_t\|^2 \right. \\ &\quad + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right) \\ &= \frac{1}{2} \min_{a_t} \left(\|a_t\|^2 + 2r \|\alpha_t s_{t-1} - a_t\| \right. \\ &\quad + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t}^T \alpha_n \right) \\ &\quad + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t}^T \alpha_n \right) \\ &\quad \cdot \left. \left(\frac{a_t^*}{t} := \alpha_t s_{t-1} \right) \right. \end{split}$$

Proof idea

$$V(y_1, \dots, y_{t-1}) = \frac{1}{2} \min_{a_t} \max_{y_t} \left(\|a_t\|^2 + 2 \left(\alpha_t s_{t-1} - a_t \right)^\top y_t + \alpha_t \|y_t\|^2 + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t+1}^T \alpha_n \right)$$

$$= \frac{1}{2} \min_{a_t} \left(\|a_t\|^2 + 2r \|\alpha_t s_{t-1} - a_t\| + \alpha_t \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{n=t}^T \alpha_n \right)$$

$$|a_t^*| = \alpha_t s_{t-1}| \qquad = \frac{1}{2} \left(\frac{\alpha_t^2}{\alpha_t^2} \|s_{t-1}\|^2 + \frac{\alpha_t}{\alpha_t} \|s_{t-1}\|^2 - \sigma_{t-1}^2 + r^2 \sum_{t=1}^T \alpha_t \alpha_t \right).$$

Principle of indifference: Any $||y_t|| = r$ is a best response.

The ball case: $\mathcal{Y} = \{y : ||y - c|| \le r\}$

Maintain statistics:
$$s_n = \sum_{t=1}^n (y_t - c), \qquad \sigma_n^2 = \sum_{t=1}^n \|y_t - c\|^2.$$

Value-to-go: quadratic in state

$$\frac{1}{2} \left(\alpha_n ||s_n||^2 - \sigma_n^2 + r^2 \sum_{t=n+1}^T \alpha_t \right).$$

$$\alpha_T = \frac{1}{T}, \quad \alpha_n = \alpha_{n+1}^2 + \alpha_{n+1} \le \frac{1}{n}.$$

Minimax strategy: affine in state

$$a_n^* = c + \alpha_n \sum_{t=1}^{n-1} (y_t - c).$$

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 Maximin distribution: same mean.

Minimax regret for ball

$$V(\mathcal{Y}) = \frac{r^2}{2} \sum_{t=1}^{T} \alpha_t.$$

Outline

- Computing minimax optimal strategies.
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$$V(y_1, \dots, y_T) := -\min_{a} \sum_{t=1}^{T} \ell(a, y_t),$$

$$V(y_1, \dots, y_{t-1}) := \min_{a_t} \max_{y_t} (\ell(a_t, y_t) + V(y_1, \dots, y_t)).$$

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$$\begin{split} V(y_1,\ldots,y_T) &:= -\min_{a} \frac{1}{2} \sum_{t=1}^{I} \|a - y_t\|^2 = -\frac{1}{2} \sum_{t=1}^{I} \|\bar{y} - y_t\|^2 \\ &= -\frac{1}{2} \sum_{t=1}^{T} \|\frac{s_T}{T} + c - y_t\|^2 \\ \text{i.e.,} \quad V(s_T, \sigma_T^2, T) &= \frac{1}{2} \left(\alpha_T \|s_T\|^2 - \sigma_T^2 \right). \end{split}$$

$$V(y_1, \dots, y_{t-1}) := \min_{a_t} \max_{y_t} (\ell(a_t, y_t) + V(y_1, \dots, y_t))$$

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$$V(y_1, \dots, y_{t-1}) = \frac{1}{2} \max_{p_t} \sum_{i=1}^{d+1} p_t(i) \left(\|Zp_t - z_i\|^2 + C_i \right)$$

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$$V(y_1, ..., y_{t-1}) = \frac{1}{2} \max_{p_t} \sum_{i=1}^{d+1} p_t(i) \left(\|Zp_t - z_i\|^2 + C_i \right)$$
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It's clear that, at each step, the unconstrained maximizer in $\{p \in \mathbb{R}^{d+1}: 1^\top p = 1\}$ keeps the value-to-go a quadratic function.

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It's clear that, at each step, the unconstrained maximizer in $\{p \in \mathbb{R}^{d+1}: 1^\top p = 1\}$ keeps the value-to-go a quadratic function. It turns out that when the simplex points z_i are on the surface of the smallest ball, the maximizer is a probability distribution.

Proof idea

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- It's clear from the proof that the maximin distribution is concentrated on the vertices of the simplex, and that a_t^* is its expectation.

The simplex case

Suppose \mathcal{Y} is a set of d+1 affinely independent points in \mathbb{R}^d , all lying on the surface of the smallest ball.

Maintain statistics: $s_n = \sum_{t=1}^n (y_t - c), \qquad \sigma_n^2 = \sum_{t=1}^n \|y_t - c\|^2.$

Value-to-go: quadratic in state

$$\frac{1}{2} \left(\alpha_n ||s_n||^2 - \sigma_n^2 + r^2 \sum_{t=n+1}^T \alpha_t \right).$$

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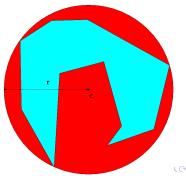
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- Computing minimax optimal strategies.
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The general case: closed, bounded $\mathcal{Y} \subset \mathbb{R}^d$

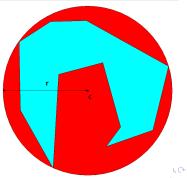
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Recall: the smallest ball containing \mathcal{Y} is $B_{\mathcal{Y}} = \{x \in \mathbb{R}^d : ||x - c|| \le r\}$. A Lagrange dual argument shows that the optimal center is in the convex hull of a set of *contact points* of \mathcal{Y} at radius r.

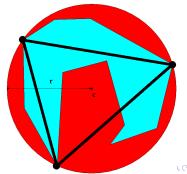


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From Carathéodory's Theorem, there is an affinely independent subset S of these contact points, with $|S| \le d + 1$.



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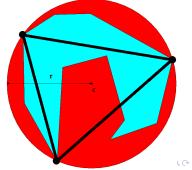
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From below

$$\mathcal{Y} \supseteq S$$
, so $V(\mathcal{Y}) \ge V(S) = \frac{r^2}{2} \sum_{i=1}^{T} \alpha_i$.



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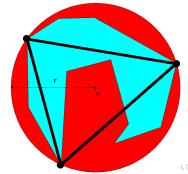
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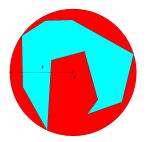
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$$\mathcal{Y} \subseteq B_{\mathcal{Y}}$$
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Main result: the role of the smallest ball



The smallest ball: B_y

The smallest ball containing \mathcal{Y} is $B_{\mathcal{Y}} = \{ y \in \mathbb{R}^d : \|y - c\| \le r \}$, with $c = \arg\min_c \max_{y \in \mathcal{Y}} \|y - c\|$, $r = \min_c \max_{y \in \mathcal{Y}} \|y - c\|$.

Main Theorem

For closed, bounded $\mathcal{Y} \subset \mathbb{R}^d$:

Minimax strategy is
$$a_{n+1}^* = n\alpha_{n+1} \frac{1}{n} \sum_{t=1}^n y_t + (1 - n\alpha_{n+1})c$$
.

Optimal regret is
$$V(\mathcal{Y}) = \frac{r^2}{2} \sum_{n=1}^{T} \alpha_n$$
.

Minimax regret

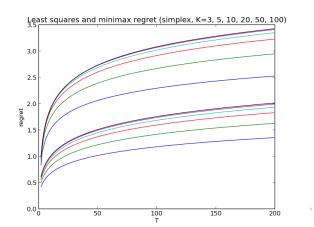
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$$V(\mathcal{Y}) = \frac{r^2}{2} \sum_{t=1}^{T} \alpha_t = \frac{r^2}{2} \left(\log T - \log \log T + O\left(\frac{\log \log T}{\log T}\right) \right).$$

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Loss $\ell(\hat{y}, y) = \frac{1}{2} \|\hat{y} - y\|^2$.

Constraints

Strategy chooses $\hat{y}_n \in \mathcal{H}$, a Hilbert space (separable, complete, inner product space).

$$\|\hat{y} - y\|^2 = \langle \hat{y} - y, \hat{y} - y \rangle.$$

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Regret
$$=\sum_{t=1}^{n}\ell(\hat{y}_t,y_t)-\inf_{a\in\mathcal{H}}\sum_{t=1}^{n}\ell(a,y_t).$$

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The smallest enclosing ball of a closed, bounded, convex ${\cal Y}$ is well-defined, with a unique center and radius.

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Recall

If $\mathcal Y$ lies in a finite-dimensional subset of $\mathcal H$, with smallest enclosing ball B(c,r), the minimax strategy is

$$a_n^* = c + \alpha_n \sum_{t=1}^{n-1} (y_t - c).$$

Theorem

For any $d \in \mathcal{H}$, the strategy

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Proof: a straightforward calculation.

Corollary

$$R_T \le \frac{r^2}{2} \sum_{t=1}^T \alpha_t$$

(By setting
$$d = c$$
.)

Theorem

For a closed, bounded, convex ${\cal Y}$

$$V_{\mathcal{T}}(\mathcal{Y}) = \frac{r^2}{2} \sum_{t=1}^{\mathcal{T}} \alpha_t.$$

Lower bound proof idea

We construct a sequence of finite sets $C_1, C_2, \ldots \subseteq \mathcal{Y}$ so that

$$\frac{r(C_i)}{r(\mathcal{Y})} \geq 1 - \sqrt{\frac{2}{i}},$$

where $r(C_i)$ is the radius of the smallest ball containing C_i .

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where $r(C_i)$ is the radius of the smallest ball containing C_i .

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To construct the C_i :

- **1** Start with $C_1 = \{y_1\}$.
- ② Set $C_{i+1} = C_i \cup \{y_{i+1}\}$, where $||c y_{i+1}|| \ge r(\mathcal{Y})$, for c the center of the smallest enclosing ball for \mathcal{Y} .

Lower bound proof idea

$$r^{2}(C_{i}) = \min_{c} \max_{y \in C_{i}} ||y - c||^{2}$$

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Evaluating $r(C_{i+1})$ and optimizing over λ gives the result.

Theorem

For a closed, bounded, convex ${\cal Y}$

$$V_{\mathcal{T}}(\mathcal{Y}) = \frac{r^2}{2} \sum_{t=1}^{\mathcal{T}} \alpha_t,$$

Theorem

For a closed, bounded, convex ${\cal Y}$

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and this is achieved by the minimax optimal strategy

$$a_n^* = c + \alpha_n \sum_{t=1}^{n-1} (y_t - c)$$

Outline

- Computing minimax optimal strategies.
- Part 1: Euclidean loss.
 - $\mathcal{Y} = \mathsf{ball}$
 - \bullet $\mathcal{Y} = simplex$
 - ullet Closed, bounded ${\cal Y}$
 - Hilbert space
 - $\mathcal{Y} = ellipsoid$
- Part 2: Linear regression.

Ellipsoid:

$$\mathcal{Y} = \{ y : (y - c)^{\top} W^{-1} (y - c) \le 1 \}$$

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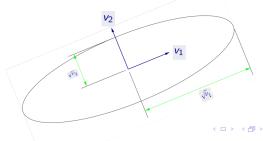
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, with $\nu_1 \ge \nu_2 \ge \cdots \ge \nu_d \ge 0$.

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Value-to-go: quadratic in state.

$$V(s_n, \sigma_n^2) = \frac{1}{2} \left(s_n^\top A_n s_n - \sigma_n^2 + \nu_1 \sum_{t=n+1}^I \alpha_n \right).$$

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Value-to-go: quadratic in state. Minimax strategy: affine in state

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Proof idea

$$V(y_1, ..., y_T) := -\min_{a} \sum_{t=1}^{r} \ell(a, y_t),$$

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The final $V(y_1, \ldots, y_T)$ is a (convex) quadratic in the state, as before:

$$V(y_1,\ldots,y_T) = \frac{1}{2} \left(s_T^\top A_T s_T - \sigma_T^2 \right).$$

Proof idea

$$V(y_1, \dots, y_{t-1}) = \min_{a_t} \max_{y_t} (\ell(a_t, y_t) + V(y_1, \dots, y_t))$$

$$= \frac{1}{2} \min_{a_t} \max_{y_t} \left(\|a_t - y_t\|^2 + (s_{t-1} + y_t)^\top A_n (s_{t-1} + y_t) - \sigma_{t-1}^2 - \|y_t\|^2 + \nu_1 \sum_{t=n+1}^T \alpha_n \right).$$

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 At each step, the inner maximum is of a (convex) quadratic criterion with a single quadratic constraint.

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- At each step, the inner maximum is of a (convex) quadratic criterion with a single quadratic constraint.
- This is a rare example of a nonconvex problem where strong duality holds.
- Evaluating the dual gives the recurrence for the value-to-go.

Minimax strategy:

$$a_n^* - c = (\alpha_{n+1}\nu_1 W^{-1} + I - A_{n+1})^{-1} A_{n+1} \sum_{t=1} (y_t - c).$$

$$\nu_1 = \lambda_{\max}(W), \quad \alpha_{n+1} = \lambda_{\max}(A_{n+1}),$$

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How does shrinkage behave?

Write
$$W = \sum_{i=1}^{d} \nu_i v_i v_i^{\top}$$
, with $\nu_1 \geq \nu_2 \geq \cdots \geq \nu_d$.

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How do the eigenvalues evolve?

$$\lambda_i^{(T)} = \frac{1}{T}, \quad \lambda_i^{(n)} = \frac{1}{1 + \lambda_1^{(n+1)} \nu_1 / \nu_i - \lambda_i^{(n+1)}} \left(\lambda_i^{(n+1)}\right)^2 + \lambda_i^{(n+1)}.$$

• $\alpha_n = \lambda_1^{(n)} \ge \lambda_2^{(n)} \ge \cdots \ge \lambda_d^{(n)}$; the gap increases with n for smaller ν_i .

Eigenvalues of A_n

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$$\bullet \ a_n^* - c = \sum_{i=1}^d \frac{\lambda_i^{(n)}}{1 + \lambda_1^{(n)} \nu_1 / \nu_i - \lambda_i^{(n)}} v_i^\top s_{n-1} v_i.$$

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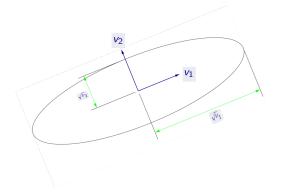
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Online prediction with quadratic loss



Subgame optimality

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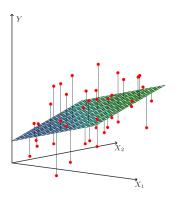
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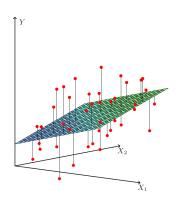
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- ullet Consider, for example, playing as if ${\mathcal Y}$ is the smallest ball when it is an ellipsoid. If the adversary plays only on the major axis, the optimal strategies are identical. If the adversary is suboptimal, the smallest ball strategy will under-regularize.

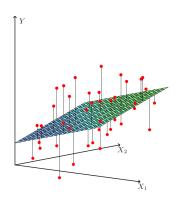
Outline

- Computing minimax optimal strategies.
- Part 1: Euclidean loss.
 - $\mathcal{Y} = \mathsf{ball}$
 - $\mathcal{Y} = simplex$
 - ullet Closed, bounded ${\cal Y}$
 - Hilbert space
 - $\mathcal{Y} = \mathsf{ellipsoid}$
- Part 2: Linear regression.
 - Fixed design.
 - Minimax strategy is regularized least squares.
 - Box and ellipsoid constraints.
 - Adversarial covariates.



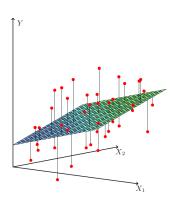


Protocol



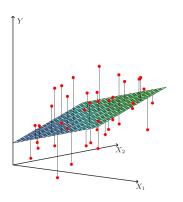
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Given: T;



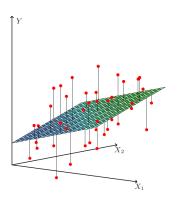
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Given: T; $x_1, \ldots, x_T \in \mathbb{R}^p$;



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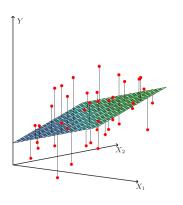
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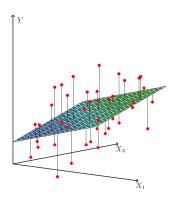


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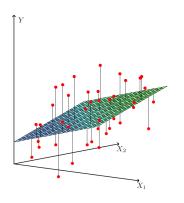
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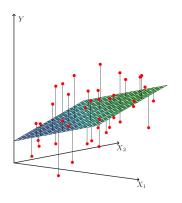
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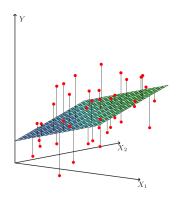
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- Learner predicts $\hat{y}_t \in \mathbb{R}$
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$$\mathsf{Regret} = \sum_{t=1}^T (\hat{y}_t - y_t)^2 - \min_{\beta \in \mathbb{R}^p} \sum_{t=1}^T \left(\beta^\top x_t - y_t \right)^2.$$

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Sufficient statistics

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Maintain statistics:
$$s_n = \sum_{t=1}^n y_t x_t$$

 $\mathcal{Y}^T = \{(y_1, \dots, y_T) : |y_t| < B_t\}.$

Value-to-go: quadratic

$$s_n^{\top} P_n s_n - \sigma_n^2 + \sum_{t=n+1}^I B_t^2 x_t^{\top} P_t x_t.$$

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Linear regression

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c.f. ridge regression:

$$\sum_{t=0}^{n} x_{t} x_{t}^{\top} + \lambda I.$$

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We'll show by induction that

$$V(s_t, \sigma_t^2, t) = s_t^{\top} P_t s_t - \sigma_t^2 + \gamma_t.$$

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Provided the problem is not too constrained (i.e., $B \ge |x_{t+1}^\top P_{t+1} s_t|$), the solution is $\hat{y}_{t+1} = x_{t+1}^\top P_{t+1} s_t$.

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$$V(s_t, \sigma_t^2, t) = s_t^{\top} \left(P_{t+1} x_{t+1} x_{t+1}^{\top} P_{t+1} + P_{t+1} \right) s_t$$
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Optimal predictions

$$\hat{y}_{n+1} = x_{n+1}^{\top} P_{n+1} s_n, \qquad s_n = \sum_{t=1}^n y_t x_t,$$

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Example

Suppose we wish to estimate a mean, and the covariate is a fixed scalar, say $x_n = 1$.

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Example

Suppose we wish to estimate a mean, and the covariate is a fixed scalar, say $x_n = 1$.

Then $P_T = 1/T$ and P_n evolves as α_n in the ball game.

An alternative recurrence

$$P_n^{-1} = \sum_{t=1}^n x_t x_t^{\top} + \sum_{t=n+1}^T \frac{x_t^{\top} P_t x_t}{1 + x_t^{\top} P_t x_t} x_t x_t^{\top}.$$

Proof

• The result is true for n = T:

$$P_T^{-1} = \sum_{t=1}^T x_t x_t^{\top}.$$

Proof

• If it's true for n, we apply the Sherman-Morrison formula:

$$(A + uv^{\top})^{-1} = A^{-1} - \frac{A^{-1}uv^{\top}A^{-1}}{1 + v^{\top}A^{-1}u},$$

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$$= \sum_{t=1}^{n} x_t x_t^{\top} + \sum_{t=n+1}^{T} \frac{x_t^{\top} P_t x_t}{1 + x_t^{\top} P_t x_t} x_t x_t^{\top} - \frac{x_n x_n^{\top}}{1 + x_n^{\top} P_n x_n}$$

$$= \sum_{t=1}^{n-1} x_t x_t^{\top} + \sum_{t=n}^{T} \frac{x_t^{\top} P_t x_t}{1 + x_t^{\top} P_t x_t} x_t x_t^{\top}.$$

Linear regression: Regret

$$\mathsf{Regret} = \sum_{t=1}^{T} B_t^2 x_t^{\top} P_t x_t.$$

Theorem

$$\max_{x_1, \dots x_T} \sum_{t=1}^T x_t^\top P_t x_t \leq p \left(1 + 2 \ln \left(1 + \frac{T}{2}\right)\right).$$

Outline

- Computing minimax optimal strategies.
- Part 1: Euclidean loss.
- Part 2: Linear regression.
 - Fixed design.
 - Minimax strategy is regularized least squares.
 - Box and ellipsoid constraints.
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Linear regression: Alternative constraints

Ellipsoid constraints (weighted 2-norm)

$$\mathcal{Y}_R^T = \left\{ (y_1, \dots, y_T) : \sum_{t=1}^T y_t^2 x_t^\top P_t x_t \leq R \right\}.$$

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$$\hat{y}_n^* = x_n^\top P_n s_{n-1}.$$

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Equalizer property

For all y_1, \ldots, y_T ,

Regret of (MM) :=
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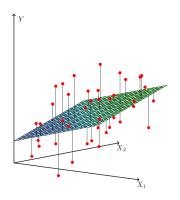
Corollary

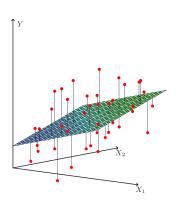
For every R, (MM) is minimax optimal on

$$\mathcal{Y}_R^T = \left\{ (y_1, \dots, y_T) : \sum_{t=1}^T y_t^2 x_t^\top P_t x_t \leq R \right\}.$$

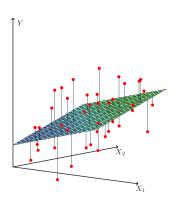
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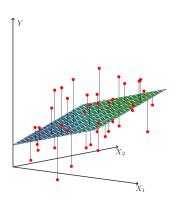


Protocol



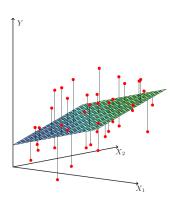
Protocol

Given: *T*;



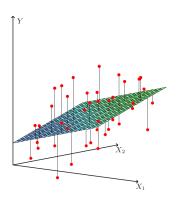
Protocol

Given: T; $\mathcal{X}^T \subset (\mathbb{R}^p)^T$;

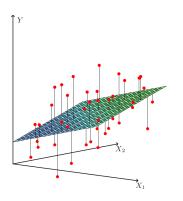


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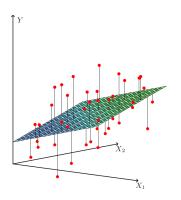


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For t = 1, 2, ..., T:

• Adversary reveals $x_t \in \mathbb{R}^p$

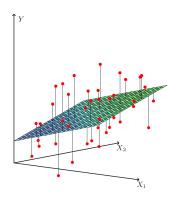


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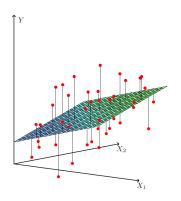
For t = 1, 2, ..., T:

• Adversary reveals $x_t \in \mathbb{R}^p \ (x_1^T \in \mathcal{X}^T)$



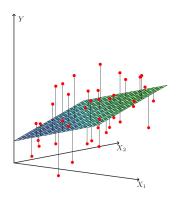
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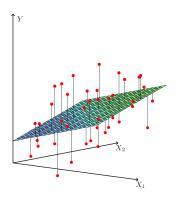
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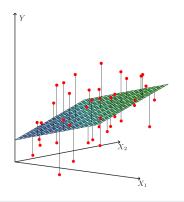
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$$\mathsf{Regret} = \sum_{t=1}^T (\hat{y}_t - y_t)^2 - \min_{\beta \in \mathbb{R}^p} \sum_{t=1}^T \left(\beta^\top x_t - y_t \right)^2.$$

A covariance budget

Recall:

$$P_T^{-1} = \sum_{t=1}^T x_t x_t^\top,$$

$$P_n = P_{n+1} x_{n+1} x_{n+1}^\top P_{n+1} + P_{n+1}.$$

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Theorem

$$P_{t+1} = P_t - \frac{a_t}{b_t^2} P_t x_{t+1} x_{t+1}^{\top} P_t,$$

where
$$a_t = rac{\sqrt{4b_t^2 + 1} - 1}{\sqrt{4b_t^2 + 1} + 1}, \ b_t^2 = x_{t+1}^ op P_t x_{t+1}.$$

Proof

Fix \tilde{P}_0 , define \tilde{P}_t by the forward iteration,

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We'll show that $P_t = \tilde{P}_t$ for t = T - 1, ..., 1.

Proof

$$P_t = P_{t+1} + P_{t+1} x_{t+1} x_{t+1}^\top P_{t+1}$$

Proof

$$P_t = \tilde{P}_{t+1} + \tilde{P}_{t+1} \boldsymbol{x}_{t+1} \boldsymbol{x}_{t+1}^\top \tilde{P}_{t+1}$$

Proof

$$\begin{split} P_{t} &= \tilde{P}_{t+1} + \tilde{P}_{t+1} x_{t+1} x_{t+1}^{\top} \tilde{P}_{t+1} \\ &= \tilde{P}_{t} - \frac{a_{t}}{b_{t}^{2}} \tilde{P}_{t} x_{t+1} x_{t+1}^{\top} \tilde{P}_{t} \\ &+ \left(\tilde{P}_{t} - \frac{a_{t}}{b_{t}^{2}} \tilde{P}_{t} x_{t+1} x_{t+1}^{\top} \tilde{P}_{t} \right) x_{t+1} x_{t+1}^{\top} \left(\tilde{P}_{t} - \frac{a_{t}}{b_{t}^{2}} \tilde{P}_{t} x_{t+1} x_{t+1}^{\top} \tilde{P}_{t} \right) \end{split}$$

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Suppose $P_{t+1} = \tilde{P}_{t+1}$. Then

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where we have used $(1 - a_t)^2 = \frac{a_t}{b_r^2}$.

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because

$$\left(\sqrt{4b_t^2+1}-1
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with orthonormal v_1, \ldots, v_m and $\lambda_1 \ge \cdots \ge \lambda_m \ge 0$.

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$$P_{t+1}^{-1} - \sum_{q=1}^{t+1} x_q x_q^{\top} = \sum_{i=1}^{m-1} \lambda_i v_i v_i^{\top}.$$

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$$P_{t+1}^{-1} - \sum_{q=1}^{t+1} x_q x_q^{\top} = \sum_{i=1}^{m-1} \lambda_i v_i v_i^{\top}.$$

Actually we can choose β as any smaller value to ensure that the rank does not drop in one step, so we can "complete the sequence" in any number of steps, provided that it is at least m.

Proof idea

To see the reverse implication, once we have computed the P_t s by the forward iteration, we can write the equivalent expression

$$P_n^{-1} = \sum_{t=1}^{n} x_t x_t^{\top} + \sum_{t=n+1}^{T} \frac{x_t^{\top} P_t x_t}{1 + x_t^{\top} P_t x_t} x_t x_t^{\top}$$

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$$\succeq \sum_{t=1}^{n} x_{t} x_{t}^{\top}.$$

Legal covariate sequences

For any $t \ge 0$, any x_1, \ldots, x_t and any P_t , the following two conditions are equivalent.

• There is a $T \ge t$ and a sequence x_{t+1}, \dots, x_T such that, under the forward iteration,

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Adversarial covariates

Thus, each $P_0 \succeq 0$ (a 'covariance budget') defines a set of sequences x_1, \ldots, x_T .

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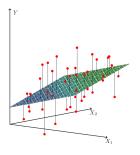
$$P_T^{-1} = \sum_{q=1}^{I} x_q x_q^{\top}.$$

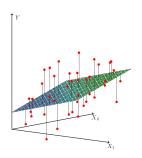
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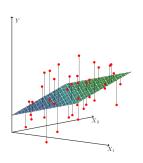
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The same strategy is optimal for each of these sequences.

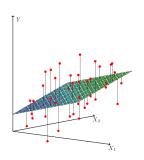






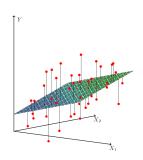


Protocol Given:



Protocol

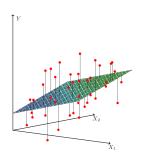
Given: $\mathcal{Z} \subset \bigcup_{T \geq 1} (\mathbb{R}^p \times \mathbb{R})^T$.



Protocol

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For t = 1, 2, 3, ...

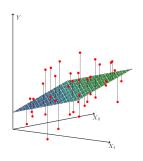


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Given: $\mathcal{Z} \subset \bigcup_{T \geq 1} (\mathbb{R}^p \times \mathbb{R})^T$.

For t = 1, 2, 3, ...

• Adversary reveals $x_t \in \mathbb{R}^p$

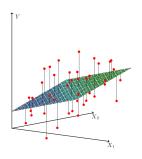


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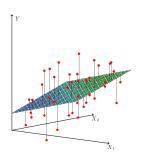
- Adversary reveals $x_t \in \mathbb{R}^p$
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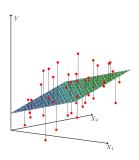


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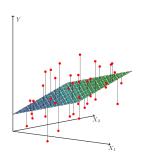
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- Learner incurs loss $(\hat{y}_t y_t)^2$.

 $(x_1^T, y_1^T \in \mathcal{Z})$



Protocol

Given: $\mathcal{Z} \subset \bigcup_{T > 1} (\mathbb{R}^p \times \mathbb{R})^T$.

For t = 1, 2, 3, ...

- Adversary reveals $x_t \in \mathbb{R}^p$
- Learner predicts $\hat{y}_t \in \mathbb{R}$
- Adversary reveals $y_t \in \mathbb{R}$ $(x_1^T, y_1^T \in \mathcal{Z})$
- Learner incurs loss $(\hat{y}_t y_t)^2$.

$$\mathsf{Regret} = \sum_{t=1}^{T} (\hat{y}_t - y_t)^2 - \min_{\beta \in \mathbb{R}^p} \sum_{t=1}^{T} \left(\beta^\top x_t - y_t \right)^2.$$

Constraints on y_t s

1 Box constraints: $\mathcal{B}(B) := \{y_1^T : |y_t| \le B_t\}$, for $B_1, B_2, \ldots > 0$.

Constraints on y_t s

- **1** Box constraints: $\mathcal{B}(B) := \{y_1^T : |y_t| \le B_t\}$, for $B_1, B_2, ... > 0$.
- 2 Ellipsoidal constraints:

$$\mathcal{E}(x_1^T, R) := \left\{ y_1^T : \sum_{t=1}^T y_t^2 x_t^\top P_t x_t \le R \right\}.$$

Constraints on x_t s

Compatibility constraints:

$$\mathcal{X}(B) = \left\{ x_1^T : B_t \ge \sum_{s=1}^{t-1} \left| x_t^\top P_t x_s \right| B_s \text{ for } 2 \le t \le T \right\}.$$

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2 Covariance constraints:

$$\overline{\mathcal{X}}(\Sigma) = \left\{ x_1^{\mathcal{T}} : \text{for } P_0, \dots, P_{\mathcal{T}} \text{ defined by } x_1^{\mathcal{T}}, \ P_0^{-1} = \Sigma \right\}.$$

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For all positive semidefinite Σ ; $B_1, B_2, \ldots > 0$ the forward strategy s^* is horizon-independent minimax optimal,

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$$\sup_{T} \sup_{x_1^T \in \mathcal{X}} \left(\sup_{y_1^T \in \mathcal{Y}(x_1^T)} R_T(s^*, x_1^T, y_1^T) - \min_{s} \sup_{y_1^T \in \mathcal{Y}(x_1^T)} R_T(s, x_1^T, y_1^T) \right) = 0.$$

with respect to the following $(\mathcal{X},\mathcal{Y}(x_1^t))$:

$$(\mathcal{X}(B_1^T) \cap \overline{\mathcal{X}}(\Sigma), \mathcal{B}(B_1^T)),$$

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That is, s^* performs as well as the best strategy that sees the covariate sequence.

The minimax strategy predicts $\hat{y}_n = \hat{\theta}_n^{\top} x_n$, where $\hat{\theta}_n$ is the solution to

$$\min_{\theta} \sum_{t=1}^{n-1} (\theta^{\top} x_t - y_t)^2 + \theta^{\top} R_n \theta,$$

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• Minimax optimal for two families of label constraints: box constraints and problem-weighted ℓ_2 norm constraints.

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- Minimax optimal for two families of label constraints: box constraints and problem-weighted ℓ_2 norm constraints.
- Strategy does not need to know the constraints.
- Regret is $O(p \log T)$.
- Same strategy is optimal for covariate sequences consistent with some 'covariance budget' P_0 .

Outline

- Computing minimax optimal strategies.
- Part 1: Euclidean loss.
- Part 2: Linear regression.